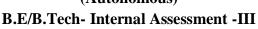


# Reg.No SNS College of Technology, Coimbatore-35.

### (Autonomous)

Academic Year 2023-2024(ODD)



**Fifth Semester** 

**Computer Science and Engineering** 

19CSB301 – Automata Theory and Compiler Design



Time: 1.5 Hours **Maximum Marks: 50** 

Part-A  $(5 \times 2 = 10 \text{Marks})$ 

CO **Blooms** 

**UND** 

CO<sub>4</sub>

1. Infer about back patching

> Backpatching is basically a process of fulfilling unspecified information. This information is of labels. It basically uses the appropriate semantic actions during the process of code generation. It may indicate the address of the Label in goto statements while producing TACs for the given expressions.

List three kind of intermediate representation 2.

CO<sub>4</sub> **REM** 

Syntax Tree

Postfix Notation

Three address Statement

3. Define Dead Code CO<sub>5</sub> **REM** 

Dead code includes code that can never be executed (unreachable code), and code that only affects dead variables (written to, but never read again), that is, irrelevant to the program.

4. What is Peephole optimization? CO<sub>5</sub> **REM** 

Peephole optimization is a simple and effective technique for locally improving target code.

Redundant Instruction Elimination

Unreachable code

5. Outline about global data flow analysis CO<sub>5</sub> **UND** 

To efficiently optimize the code compiler collects all the information about the program and distribute this information to each block of the flow graph. This process is known as data-flow graph analysis. Certain optimization can only be achieved by examining the entire program.

### Part-B (2x13+14=40Marks)

C a. Construct the CLR parsing table for the following grammar. check whether 04 P the string (a) is accepted or not.

 $S \rightarrow (L)|a$ 

L->L,S|S

1. Construct Augumented Grammar

P

- 2. Construct Canonical LR(1) items
- 3. Construct CLR Parsing table
- 4.Parsing i/p string using CLR Parse table
  After the elimination of left recursion: (2)

S->(L) | a L->SLí Lí->,SLí | Â

Calculation of First: (2)

First(S) =  $\{(, a\}$ First(L) =  $\{(, a\}$ First(Li) =  $\{,, \hat{A}\}$ 

Calculation of Follow: (2)

Follow(S) =  $\{\$, , ,)\}$ Follow (L) =  $\{\}$ Follow (Li) =  $\{\}$ 

Predictive parsing table: (4)

Non	Input symbol							
terminals	a	(	)	,	\$			
S	S->a	S->(L)						
L	L->SLí	L->SLí						
Lí			Lí->Â	Lí->,SLí				

or

- b. Illustrate the Storage allocation strategies in perspective of compiler with neat diagram.
  - The different storage allocation strategies are :

- .
- Static **allocation** lays out **storage** for all data objects at compile time.
- Stack **allocation** manages the run-time **storage** as a stack.
- Heap **allocation** allocates and deallocates **storage** as needed at run time from a data area known as heap.
- 7. a. Demonstrate about the translation scheme to generate three address code for the Boolean Expression

nn 13 C O4

13

C

04

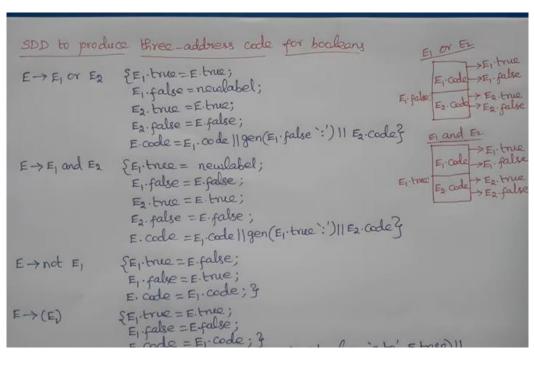
U

N

D

Α

P P



b.		Explain the principle sources of code optimization in detail with example  1. Common-Sub expression Elimination:  2. Copy Propagation			U N D
		3. Dead Code Elimination			
		4. Constant folding			
		5. Loop Optimizations			
8.	a.	(i) What is an activation record? Explain how it is related with run time storage	10	C	A
	organization			O4	N
	An activation record is a contiguous block of storage that manages information requ by a single execution of a procedure		4		
				C	
				O4	U
		Return Values			N
					D
		Parameter List			
		Control Links			

# (ii) Summarize on back patching

Access Links

Saved Machine Status

Local Data

Temporaries

6. t6:=Prod+t5 7. Prod:=t6 8. t7:=i+1 9. i:= t7

10. if  $i \le 20$  goto (1).

Backpatching is basically **a process of fulfilling unspecified information**. This information is of labels. It basically uses the appropriate semantic actions during the process of code generation. It may indicate the address of the Label in goto statements while producing TACs for the given expressions

or

b. Construct the DAG for the following Basic block & explain it.

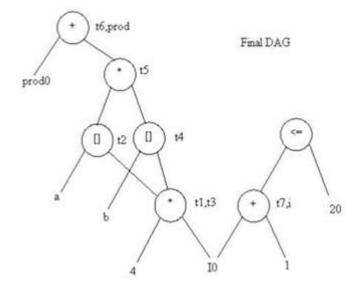
1. t1: = 4 \* i

2. t2:= a [t1]

3. t3: = 4 \* i

4. t4:= b [t3]

5. t5:=t2\*t4



Und-Understanding Rem-Remembering App-Applying Ana-AnalyzeCre-Creating Eva-Evaluating

Prepared By Verified By HoD