



# SNS COLLEGE OF TECHNOLOGY

(Autonomous)

COIMBATORE-35



## *Basic blocks and Flow graphs & Next Use Information*

(1) prod := 0  
(2) i := 1  
(3) t<sub>1</sub> := 4\* i  
(4) t<sub>2</sub> := a[t<sub>1</sub>]  
(5) t<sub>3</sub> := 4\* i  
(6) t<sub>4</sub> := b[t<sub>3</sub>]  
(7) t<sub>5</sub> := t<sub>2</sub> \* t<sub>4</sub>  
(8) t<sub>6</sub> := prod + t<sub>5</sub>  
(9) prod := t<sub>6</sub>  
(10) t<sub>7</sub> := i + 1  
(11) i := t<sub>7</sub>  
(12) if i <= 20 goto (3)

(1) prod := 0  
(2) i := 1

B<sub>1</sub>

(3) t<sub>1</sub> := 4\* i  
(4) t<sub>2</sub> := a[t<sub>1</sub>]  
(5) t<sub>3</sub> := 4\* i  
(6) t<sub>4</sub> := b[t<sub>3</sub>]  
(7) t<sub>5</sub> := t<sub>2</sub> \* t<sub>4</sub>  
(8) t<sub>6</sub> := prod + t<sub>5</sub>  
(9) prod := t<sub>6</sub>  
(10) t<sub>7</sub> := i + 1  
(11) i := t<sub>7</sub>  
(12) if i <= 20 goto (3)

B<sub>2</sub>



# *Basic blocks and Flow graphs*

## **Basic Block:**

A **basic block** is a **sequence** of consecutive statements in which **flow of control enters** at the beginning and leaves at the end without halt or possibly of the branching except at the end.

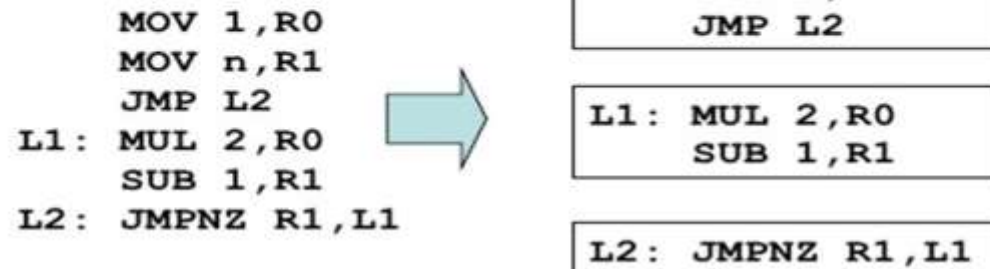
- **Flow Graph:** A **graph representation** of three address statements, called flow graph.
- **Nodes** in the flow graph represent **computations**.
- **Edges** represent the **flow of control**.
- Used to do **better job** of **register allocation** and **instruction selection**.



# Basic blocks and Flow graphs

## Basic Blocks

- A *basic block* is a sequence of consecutive instructions with exactly one entry point and one exit point (with natural flow or a branch instruction)



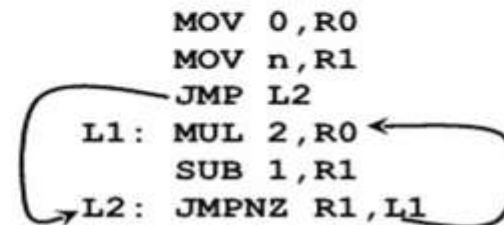


# *Basic blocks and Flow graphs*

## Flow Graphs

- A *flow graph* is a graphical depiction of a sequence of instructions with control flow edges
- A flow graph can be defined at the intermediate code level or target code level

```
MOV 1, R0
MOV n, R1
JMP L2
L1: MUL 2, R0
SUB 1, R1
L2: JMPNZ R1, L1
```



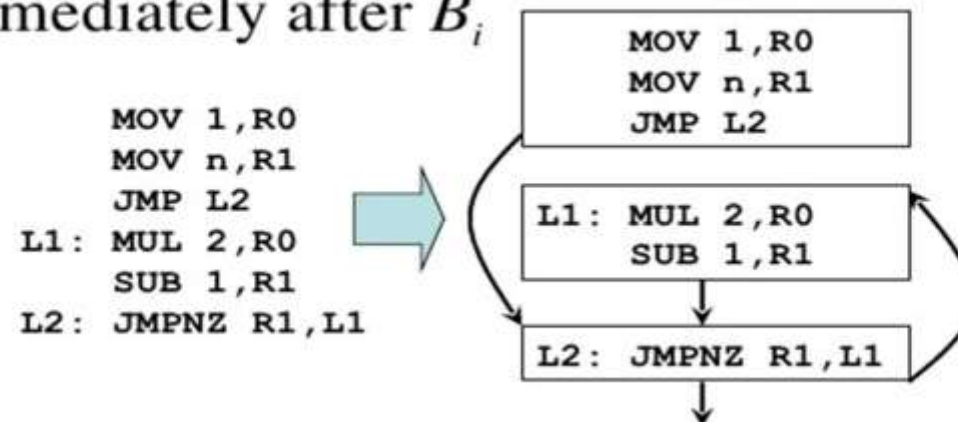


# Basic blocks and Flow graphs

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## Basic Blocks and Control Flow Graphs

- A *control flow graph* (CFG) is a directed graph with basic blocks  $B_i$  as vertices and with edges  $B_i \rightarrow B_j$  iff  $B_j$  can be executed immediately after  $B_i$



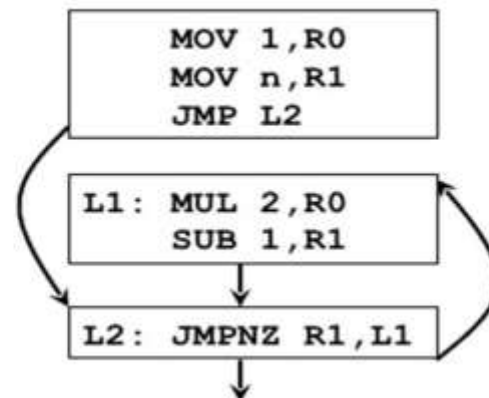


# Basic blocks and Flow graphs

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## Successor and Predecessor Blocks

- Suppose the CFG has an edge  $B_1 \rightarrow B_2$ 
  - Basic block  $B_1$  is a *predecessor* of  $B_2$
  - Basic block  $B_2$  is a *successor* of  $B_1$





# *Basic blocks and Flow graphs*

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## Loops

- A *loop* is a collection of basic blocks, such that
  - All blocks in the collection are *strongly connected*
  - The collection has a unique *entry*, and the only way to reach a block in the loop is through the entry

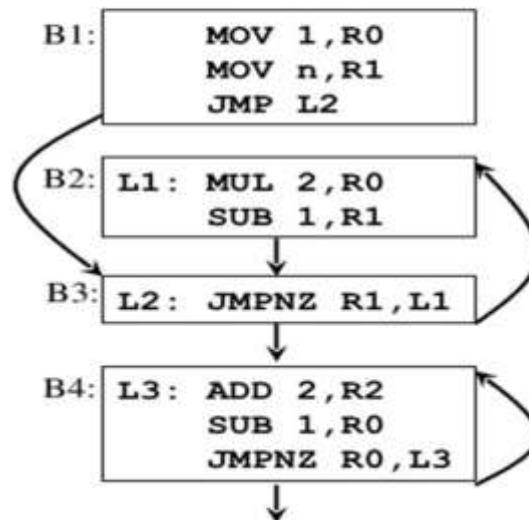




# *Basic blocks and Flow graphs*

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## Loops (Example)



Strongly connected components:

SCC = { {B2, B3},  
{B4} }

Entries:  
B3, B4





# Basic blocks and Flow graphs

## Construction of the Representation

1. Partition the intermediate code into basic blocks
2. The basic blocks become the nodes of a flow graph and the edges indicate the *flow* (which blocks follow which)

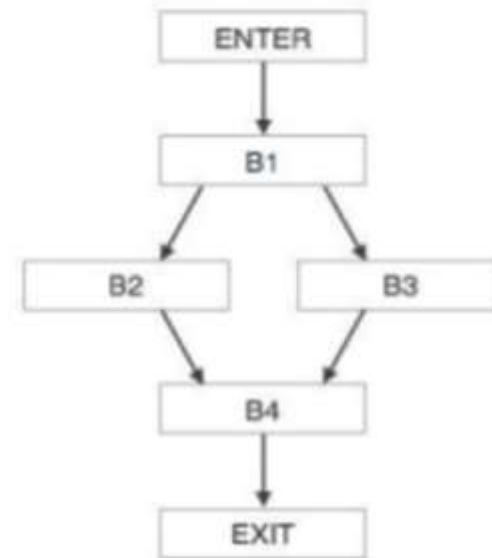
**B1**  
w = 0;  
x = x + y;  
y = 0;  
if( x > z)

**B2**  
y = x;  
x++;

**B3**  
y = z;  
z++;

**B4**  
w = x + z;

Basic Blocks



Flow Graph



# Basic blocks and Flow graphs

## Constructing Basic Blocks

Determine a set of leaders

- 1) The first instruction is a leader
- 2) Instruction L is a leader if there is an instruction `if ... goto L`  
or `goto L`
- 3) Instruction L is a leader if it immediately follows an instruction `if ... goto B`  
or `goto B`

```
1) i = 1
2) j = 1
3) t1 = 10 * i
4) t2 = t1 + j
5) t3 = 8 * t2
6) t4 = t3 - 88
7) a[t4] = 0.0
8) j = j + 1
9) if j <= 10 goto (3)
10) i = i + 1
11) if i <= 10 goto (2)
12) i = 1
13) t5 = i - 1
14) t6 = 88 * t5
15) a[t6] = 1.0
16) i = i + 1
17) if i <= 10 goto (13)
```

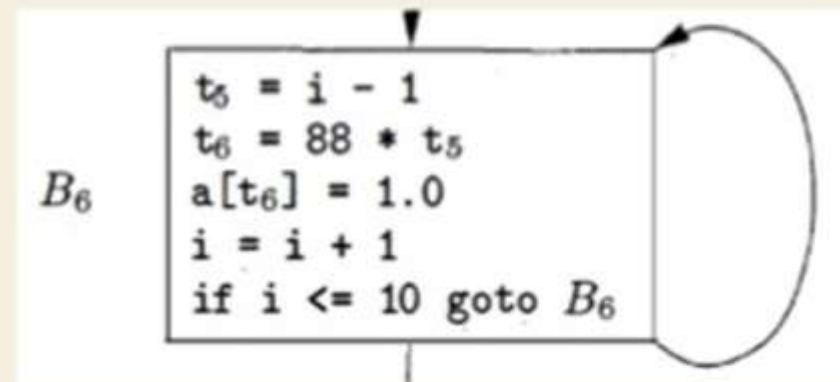


# *Basic blocks and Flow graphs*

## *Constructing the Flow Graph*

There is an edge from block *B* to block *C* iff it is possible for the first instruction in block *C* to immediately follow the last instruction in block *B*

1. There is a conditional or unconditional jump from the end of *B* to the beginning of *C*



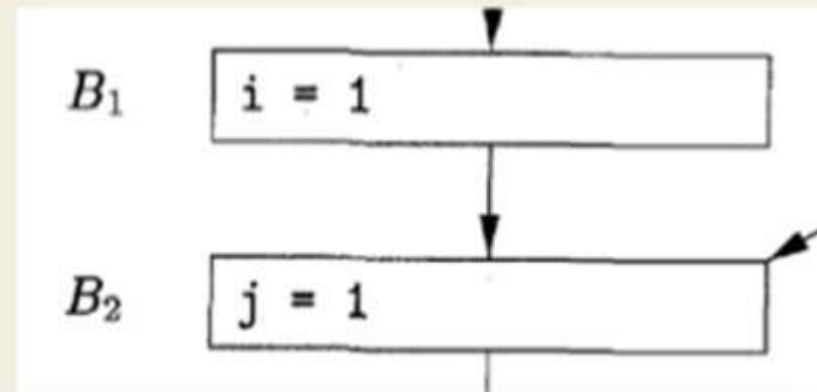


# Basic blocks and Flow graphs

## Constructing the Flow Graph

There is an edge from block  $B$  to block  $C$  iff it is possible for the first instruction in block  $C$  to immediately follow the last instruction in block  $B$

2.  $C$  immediately follows  $B$  in the original order of the three-address instructions, and  $B$  does not end in an unconditional jump



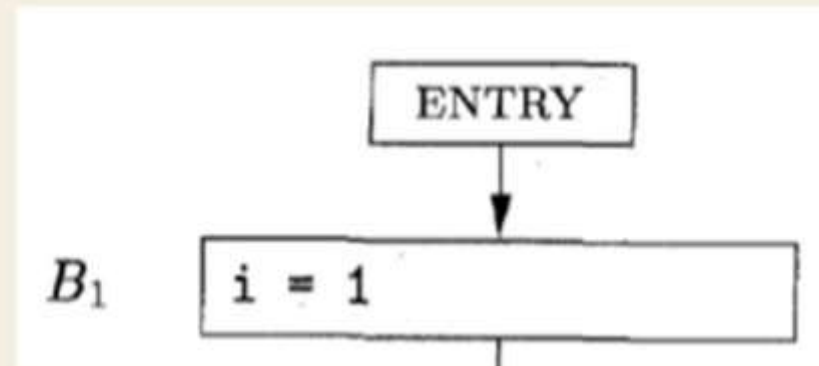


# *Basic blocks and Flow graphs*

## *Constructing the Flow Graph*

Often we add two nodes, called  
the entry and exit.

1. There is an edge from the entry to the first executable node.



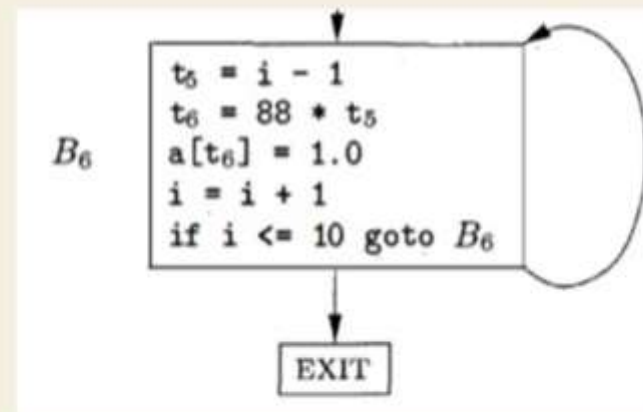


# Basic blocks and Flow graphs

## Constructing the Flow Graph

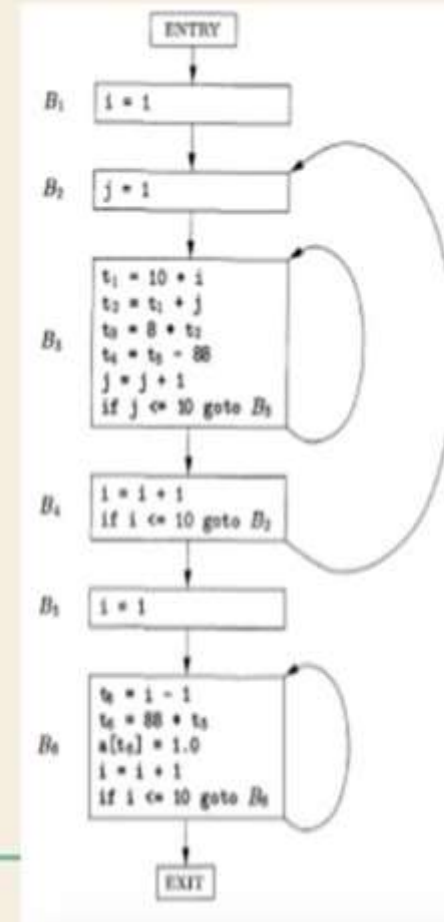
Often we add two nodes, called the entry and exit.

2. There is an edge to the exit from any basic block that contains an instruction that could be the last executed instruction of the program





```
1) i = 1
2) j = 1
3) t1 = 10 * i
4) t2 = t1 + j
5) t3 = 8 * t2
6) t4 = t3 - 88
7) a[t4] = 0.0
8) j = j + 1
9) if j <= 10 goto (3)
10) i = i + 1
11) if i <= 10 goto (2)
12) i = 1
13) t5 = i - 1
14) t6 = 88 * t5
15) a[t6] = 1.0
16) i = i + 1
17) if i <= 10 goto (13)
```







# *Summarization*