



Intermediate Code Optimization

Unit V

- Program transformation technique
- Improves code consume less resources
- Transforms the code to make it more efficient
- Output is not changed
- Intermediate code optimization generation is made easier

Intermediate Code Generator		
•		
<pre>t1 = inttofloat(60)</pre>		
t2 = id3 * t1		
t3 = id2 + t2		
id1 = t3		
+		
Code Optimizer		
•		
t1 = id3 * 60.0		
id1 = id2 + t1		



Code Optimization



- Optimization
 - Machine Independent Optimization

do

- takes in the intermediate code and transforms a part of the code that does not involve any CPU registers
- Example:

```
{
    item = 10;
    value = value + item;
} while(value<100);</pre>
```

This code involves repeated assignment of the identifier item, which if we put this way:

```
Item = 10;
do
{
  value = value + item;
} while(value<100);</pre>
```

- Machine Dependent optimization

- Target code
 - Rearrangement of machine instructions to improve the efficiency of the code
 - Divide the code into basic blocks





Peephole Optimization

- Optimization eliminates the redundant instruction from a small area of code
- Set of code peephole / window
- Goals :
 - Improves performance
 - Reduce memory footprint
 - Reduce code size





Principle sources of optimization





Compile Time Evaluation



• Constant Folding

- Folding the constants
- The expressions that contain the operands having constant values at compile time are evaluated.
- <u>Example:</u>
- return (3+5); 2 return 8;
- Cir=(22/7)*diameter 2cir = 3.14*diameter
- Constant Propagation
 - If some variable has been assigned some constant value, then it replaces that variable with its constant value in the further program during compilation.
 - <u>Example:</u>
 - radius =10,pi=3.14
 - area=pi*radius*radius; 2area=3.14*10*10;





Common Sub Expression

Code before Optimization	Code after Optimization
S1 = 4 x i	S1 = 4xi
S2 = a[S1]	$S_{2} = a[S_{1}]$
S3 = 4 x j	$S_3 = 4 x i$
S4 = 4 x i // Redund ant Expression	00 - TAJ
\$5 = n	50 = n
S6 = b[S4] + S5	S6 = b[S1] + S5





Code Movement

Code before Optimization	Code after Optimization
for $(int j = 0; j < n; j ++)$	$\mathbf{x} = \mathbf{y} + \mathbf{z};$
{	for (int $j = 0; j < n; j ++$)
$\mathbf{x} = \mathbf{y} + \mathbf{z};$	{
$\mathbf{a[j]} = 6 \mathbf{x} \mathbf{j};$	$\mathbf{a[j]} = 6 \mathbf{x} \mathbf{j};$
}	}
1	1





Dead Code Elimination

• Eliminates the dead code

Code before Optimization

Code after Optimization







Strength Reduction

- Reduces the strength of expressions
- Replaces expensive operators with cheaper one
- <u>Example</u>
 - − B=A*2 ?B=A+A
 - Cost of multiplication is higher than the addition