



Intermediate Code generation for *Boolean Expressions*



- Boolean Expression
 - Logical values
 - Conditional Expression – change the flow of program (if-else, do-while)
- Boolean operator
 - And
 - Or (lowest precedence)
 - Not
- Example
 - $E \rightarrow E \text{ or } E$
 - $E \rightarrow E \text{ and } E$
 - $E \rightarrow \text{not } E$
 - $E \rightarrow (E)$
 - $E \rightarrow \text{id relop id}$
 - $E \rightarrow \text{TRUE}$ $E \rightarrow \text{id}$
 - $E \rightarrow \text{FALSE}$



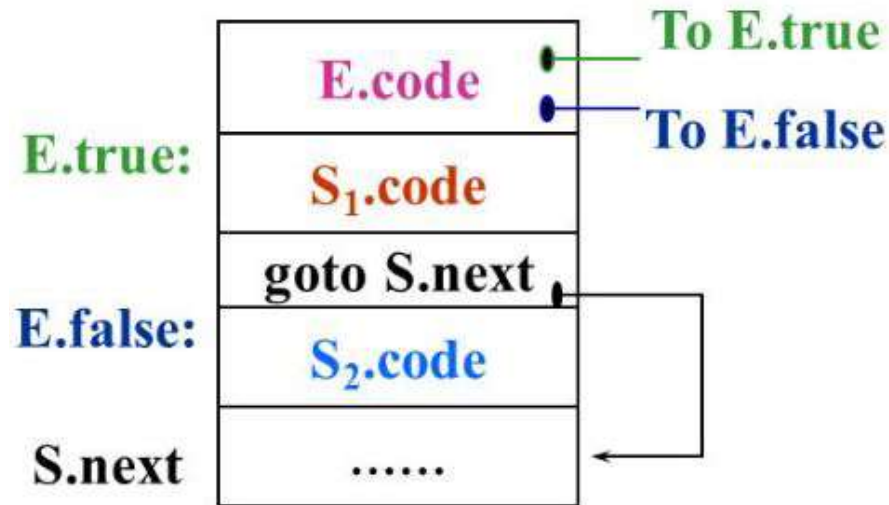
Intermediate Code generation for *Boolean Expressions*



- Numerical representation of Boolean Expression
 - Example1: A or B and C
 - Three Address Sequence:
 - T1=B and C
 - T2=A or T1
 - Example2: $A < B \rightarrow$ if A<B then 1 else 0
 - Three Address Sequence:
 - 1. If A<B goto (4)
 - 2. T=0
 - 3. goto (5)
 - 4. t=1
 - 5. ---



Attributes used for “if E then S1 else S2”



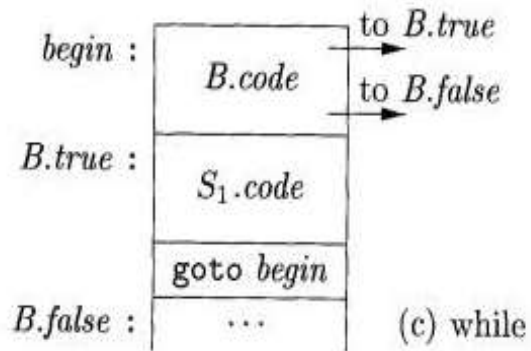
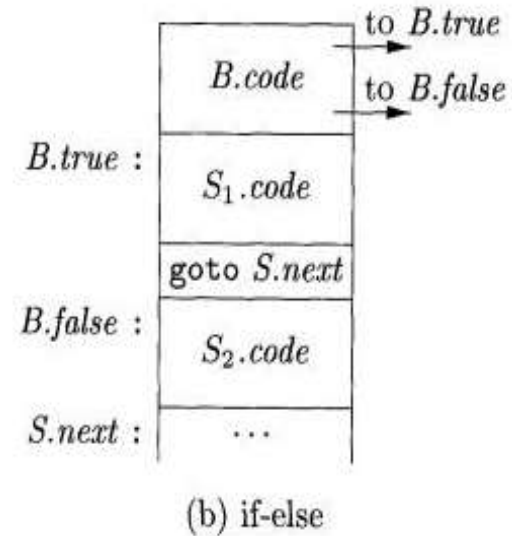
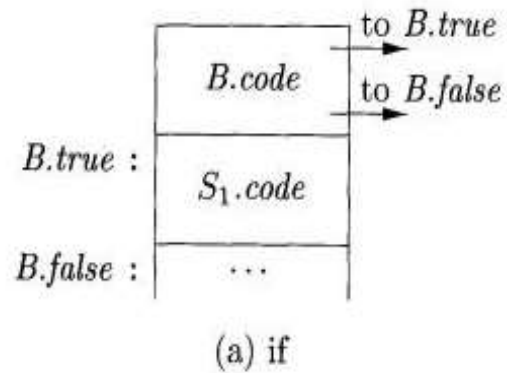


Flow of Control Statements

If E then S1

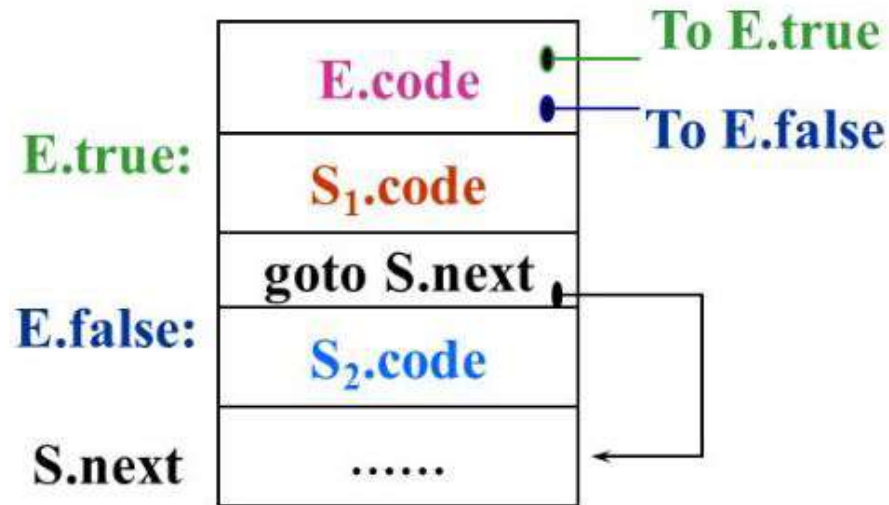
IF E then S1 else S2

While E do S1





Attributes used for “if E then S1 else S2”



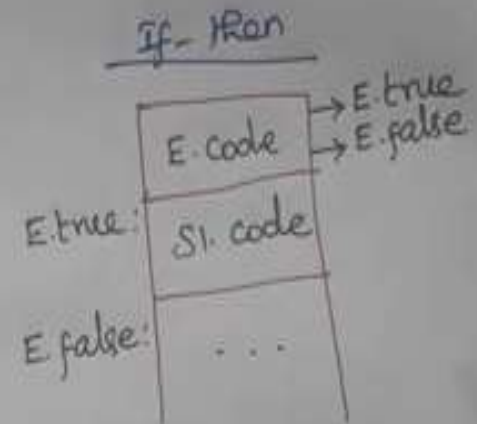


Flow of Control Statements

Flow of control statements

$S \rightarrow \text{if } E \text{ then } S_1 \quad \left\{ \begin{array}{l} E.\text{true} = \text{newlabel}; \\ E.\text{false} = S.\text{next}; \\ S_1.\text{next} = S.\text{next}; \\ S.\text{code} = E.\text{code} \parallel \text{gen}(E.\text{true} ':') \\ \parallel S_1.\text{code} \end{array} \right\}$

$S \rightarrow \text{if } E \text{ then } S_1 \text{ else } S_2 \quad \left\{ \begin{array}{l} E.\text{true} = \text{newlabel}; \\ E.\text{false} = \text{newlabel}; \\ S_1.\text{next} = S.\text{next}; \\ S_2.\text{next} = S.\text{next}; \\ S.\text{code} = E.\text{code} \parallel \text{gen}(E.\text{true} ':') \\ \parallel S_1.\text{code} \parallel \text{gen}(\text{goto}' S.\text{next}) \\ \parallel \text{gen}(E.\text{false} ':') \parallel S_2.\text{code} \end{array} \right\}$





```
do while E do S, S.begin : 2 new label ;  
E.true : 2 new label ;  
E.false : 2 S.next ;  
S1.next : 2 S.begin ;  
S.code : 2 gen (S.begin : 1) || E.code ||  
gen (E.true : 1) || S1.code ||  
gen ('goto' S.begin)
```



Control Flow Translation of Boolean Expressions

Short-circuit code (or) Jumping code

- without generating code for the boolean operators
- without evaluating the entire expression

E_1 or E_2 E_1 and E_2

$a < b$

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if a < b goto E.true  
goto E.false
```



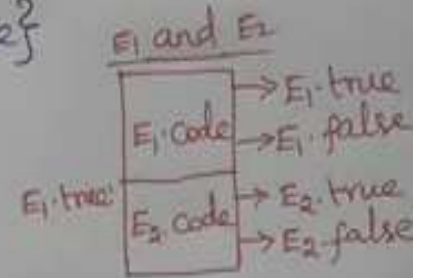

Boolean Expressions

SDD to produce three-address code for booleans

$E \rightarrow E_1 \text{ OR } E_2$ $\{ E_1.\text{true} = E.\text{true};$
 $E_1.\text{false} = \text{newlabel};$
 $E_2.\text{true} = E.\text{true};$
 $E_2.\text{false} = E.\text{false};$
 $E.\text{code} = E_1.\text{code} \parallel \text{gen}(E_1.\text{false} ':') \parallel E_2.\text{code} \}$



$E \rightarrow E_1 \text{ AND } E_2$ $\{ E_1.\text{true} = \text{newlabel};$
 $E_1.\text{false} = E.\text{false};$
 $E_2.\text{true} = E.\text{true};$
 $E_2.\text{false} = E.\text{false};$
 $E.\text{code} = E_1.\text{code} \parallel \text{gen}(E_1.\text{true} ':') \parallel E_2.\text{code} \}$



$E \rightarrow \text{not } E_1$ $\{ E_1.\text{true} = E.\text{false};$
 $E_1.\text{false} = E.\text{true};$
 $E.\text{code} = E_1.\text{code}; \}$

$E \rightarrow (E_1)$ $\{ E_1.\text{true} = E.\text{true};$
 $E_1.\text{false} = E.\text{false};$
 $E.\text{code} = E_1.\text{code}; \}$

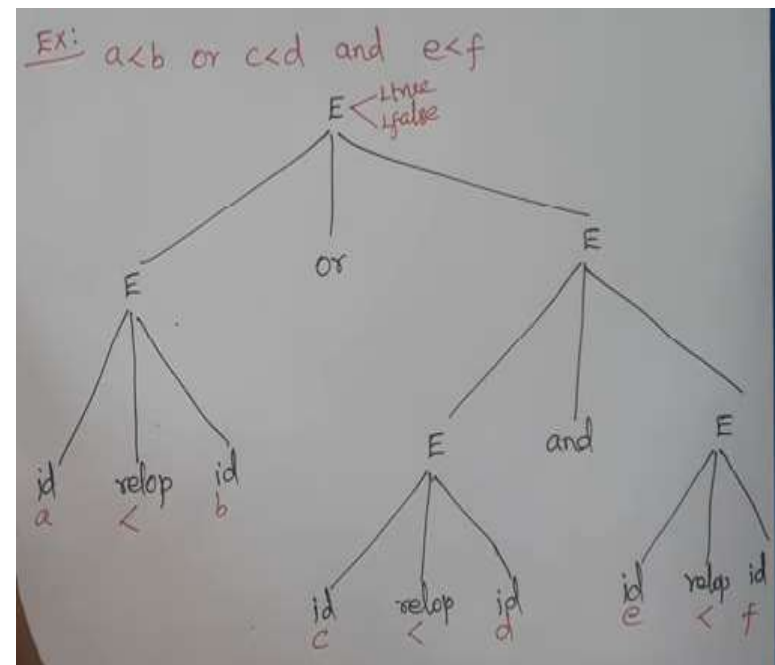


$F \rightarrow (E_1)$ $\{E_1.true = E.true;$
 $E_1.false = E.false;$
 $E.code = E_1.code;\}$

$E \rightarrow id_1, relop, id_2$ $\{E.code = gen(\text{if } id_1 \text{ place } relop \text{ place } id_2 \text{ place 'goto' } E.true) ||$
 $gen(\text{'goto' } E.false)\}$

$E \rightarrow true$ $\{E.code = gen(\text{'goto' } E.true)\}$

$E \rightarrow false$ $\{E.code = gen(\text{'goto' } E.false)\}$





Three Address Code

$a < b$ or $c < d$ and $e < f$

if $a < b$ goto Ltrue
goto L1

L1: if $c < d$ goto L2
goto Lfalse

L2: if $e < f$ goto Ltrue
goto Lfalse



BACKPATCHING

- Process of backpatching
 - A marker Non-terminal M – next instruction to be executed
 - Example
 - $E \rightarrow E1$ and M E2
 - Incomplete jumps with unfilled labels \rightarrow E.truelist and E.falselist
 - E1 – false , E is also false \rightarrow E1.falselist becomes a part of E.flaselist
 - E1 – true \rightarrow E2 test \rightarrow E1.truelist becomes the beginning code for E2 \leftarrow marker non-terminal M