SNS COLLEGE OF TECHNOLOGY

Coimbatore-35
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DEPARTMENT OF INFORMATION TECHNOLOGY

19ITB201 – DESIGN AND ANALYSIS OF ALGORITHMS

II YEAR IV SEM

UNIT-II-BRUTE FORCE AND DIVIDE AND CONQUER

TOPIC: Divide and Conquer –Quick Sort

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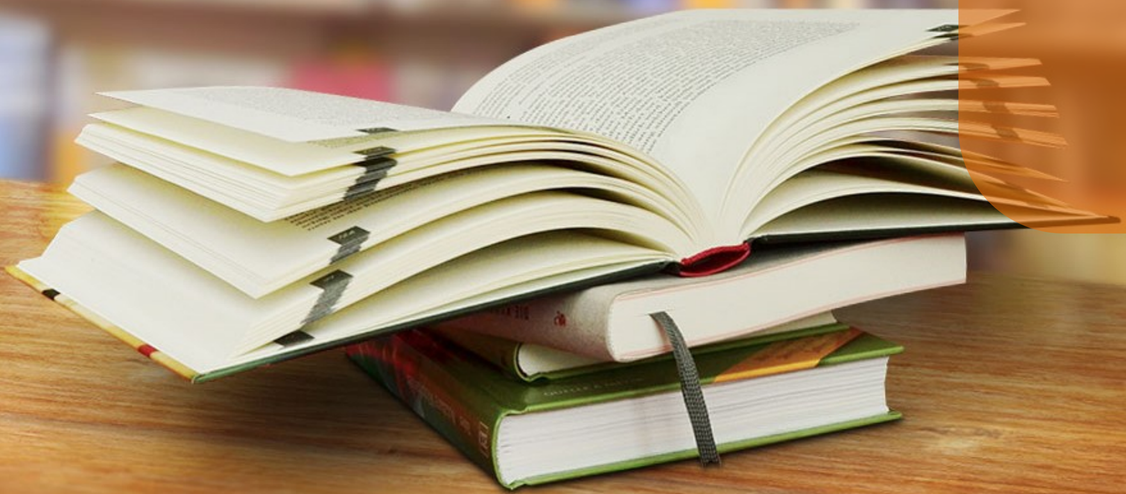
DIVIDE AND CONQUER-QUICK SORT

Subject :Design and Analysis of Algorithm
Unit :II

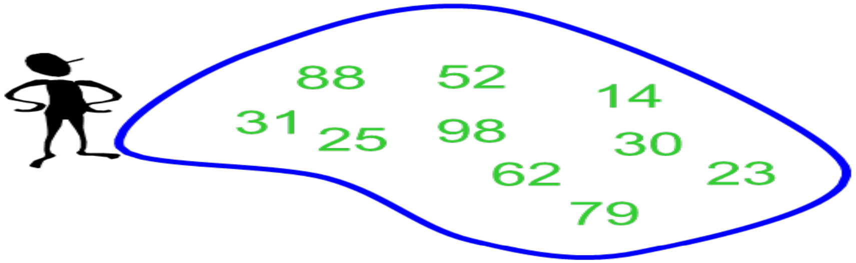


Identify the problem

Sorting

Problem Example

Quick Sort

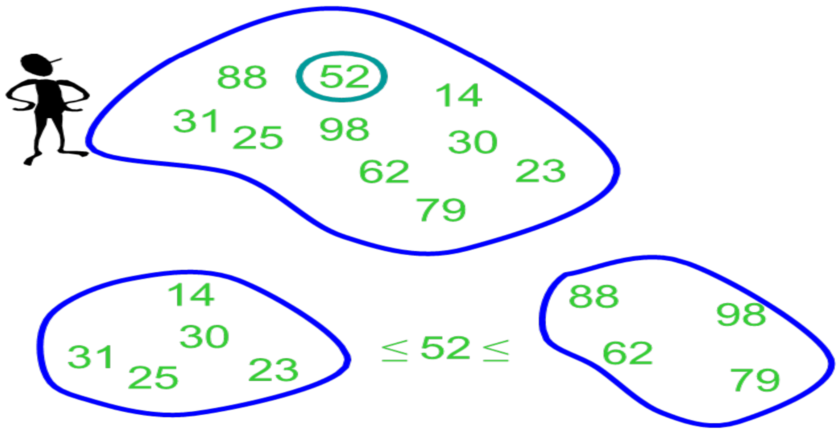


Divide and Conquer

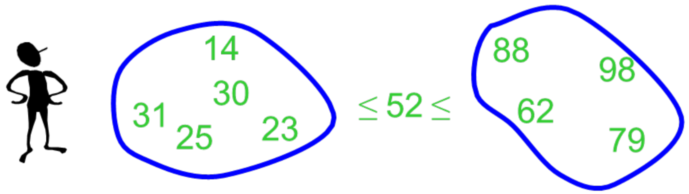
Problem Example

Quick Sort

Partition set into two using randomly chosen pivot



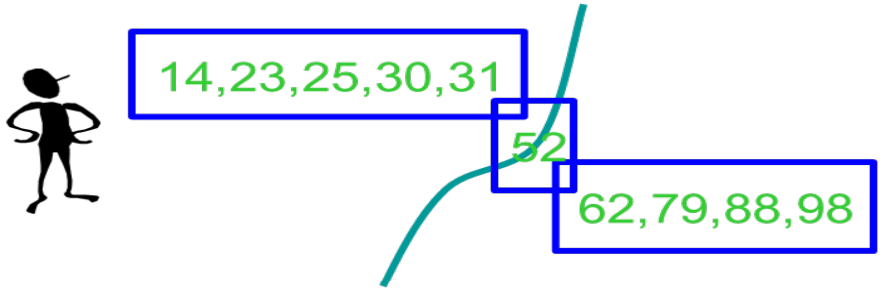
Quick Sort



sort the first half.

sort the second half.

Quick Sort



Glue pieces together.

14,23,25,30,31,52,62,79,88,98

Quicksort

- Quicksort pros [advantage]:
  - Sorts in place
  - Sorts $O(n \lg n)$ in the average case
  - Very efficient in practice , it's quick

Quicksort cons [disadvantage]:

- – Sorts $O(n^2)$ in the worst case
  - And the worst case doesn't happen often … sorted

Quicksort

Another divide-and-conquer algorithm:

Divide: **A[p…r]** is partitioned (rearranged) into two nonempty subarrays **A[p…q-1]** and **A[q+1…r]** s.t. each element of $A[p…q-1]$ is less than or equal to each element of $A[q+1…r]$. Index q is computed here, called **pivot**.

Conquer: two subarrays are **sorted by recursive calls** to quicksort.

Combine: unlike merge sort, no work needed since the subarrays are sorted in place already.

Algorithm

ALGORITHM *Quicksort(A*[*l..r*]*)*

//Sorts a subarray by quicksort

//Input: Subarray of array *A*[0..*n* 1]*,* defined by its left and right

//indices *l and r*

//Output: Subarray *A*[*l..r*] sorted in nondecreasing order

if *l* < *r*

s ←*Partition(A*[*l..r*]*)* //*s* is a split position

Quicksort(A[*l....s* − 1]*)*

Quicksort(A[*s* + 1*.....r*]*)*

Algorithm

ALGORITHM *HoarePartition(A*[*l..r*]*)*

//Partitions a subarray by Hoare's algorithm, using the first element

//as a pivot

//Input: Subarray of array *A*[0..*n*-1], defined by its left and right

//indices *l* and *r (l < r)*

//Output: Partition of *A*[*l..r*], with the split position returned as

//this function's value

P=A[l]

i=l;j=r+1;

repeat

    repeat i=i+1 until A[i]>=p

    repeat j=j-1 until A[j]<=p

swap*(A*[*i*]*, A*[*j*]*)*

Until i>=j

swap*(A*[*i*]*, A*[*j*]*)* //undo last swap when *i>= j*

Swap (A[*l*]*, A*[*j*]*)*

return *j*

Complexity Analysis of Quick Sort

Worst Case Time Complexity [Big-O]: **$O(n^2)$**

Best Case Time Complexity [Big-omega]: **O(n*log n)**

Average Time Complexity [Big-theta]: **O(n*log n)**

Space Complexity: **O(n*log n)**

Assessment

1. Which of the following sorting algorithms is the fastest?
   a) Merge sort
   b) Quick sort
   c) Insertion sort
   d) Shell sort
2. Quick sort follows Divide-and-Conquer strategy.
   a) True
   b) False
3. What is the worst case time complexity of a quick sort algorithm?
   a) O(N)
   b) O(N log N)
   c) $O(N^2)$
   d) O(log N)


Thank you!