



# SNS COLLEGE OF TECHNOLOGY

Coimbatore-35  
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## DEPARTMENT OF INFORMATION TECHNOLOGY

19ITB201– Design and Analysis of Algorithms

II YEAR IV SEM

UNIT 4 – FLOW NETWORKS AND STRING MATCHING

TOPIC – String Matching-Navie String Matching Algorithm



# NAIVE STRING MATCHING



# What is a String?



In C, String is a sequence of characters or more specifically can be regarded as an array of characters.



# String Matching



One of the most commonest operation done with strings is String Matching.



# WHAT IS STRING MATCHING?

If we are given a string “p” which is nothing but stream of characters  $p[0], p[1], p[2], \dots, p[n-1]$  (where n is the length of the string), then searching for a specific pattern in the string is known as string matching



# Applications of String Matching

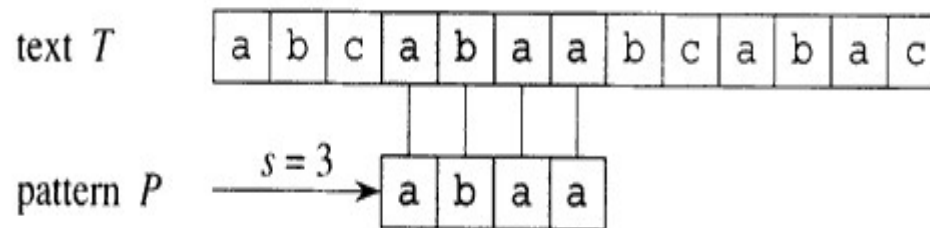
- Intrusion Detection
- String matching in bioinformatics
- String matching in Digital Forensics



# An example of string matching



Here we discuss an example on string matching where the given text consists of a pattern which is to be searched for. The following figure shows the location of pattern  $P$  in a given text  $T$ .





# Algorithm For String Matching



There are many ways to do the String Matching but here we discuss the concept of Naïve String Matching.





# Naive string matching



The naive algorithm finds all valid shifts using a loop that checks the condition  $P[1..m] = T[s + 1..s + m]$  for each of the  $n - m + 1$  possible values of  $s$ .

NAIVE-STRING-MATCHER( $T, P$ )

```
1   $n \leftarrow \text{length}[T]$ 
2   $m \leftarrow \text{length}[P]$ 
3  for  $s \leftarrow 0$  to  $n - m$ 
4      do if  $P[1..m] = T[s + 1..s + m]$ 
5          then print "Pattern occurs with shift"  $s$ 
```

*Running time:  $O((n-m+1)m)$ .*



## Problem with naive algorithm



- Suppose  $p=ababc$ ,  $T=cabababcd$
- T: c a b a b a b c d
- P: a ...
- P: a b a b c
- P: a ...
- P: a b a b c
- Whenever a character mismatch occurs after matching of several characters, the comparison begins by going back in T from the character which follows the last beginning character. This makes the process very slow.



# Solution to the Problem

- There are many more algorithms which work in a more efficient way than the naive string matching and are listed below:-
- Rabin–Karp string search algorithm  
Finite-state automaton  
Knuth–Morris–Pratt algorithm  
Boyer–Moore string search algorithm



**THANK YOU**