



EUCLID'S ALGORITHM



Subject :Design and Analysis of Algorithm
Unit :I





GCD(Greatest common divisor)



- The greatest common divisor of two nonnegative, not-both-zero integers m and n , denoted $\gcd(m, n)$, is defined as the largest integer that divides both m and n evenly, i.e., with a remainder of zero
- Euclid of Alexandria (third century B.c.) outlined an algorithm for solving this problem in one of the volumes of his *Elements* most famous for its systematic exposition of geometry





Euclid's algorithm

Euclid's algorithm is based on applying repeatedly the equality

$$\gcd(m, n) = \gcd(n, m \bmod n),$$

Where $m \bmod n$ is the remainder of the division of m by n , until $m \bmod n$ is equal to 0. Since $\gcd(m, 0) = m$ (why?), the last value of m is also the greatest common divisor of the initial m and n .

For example, $\gcd(60, 24)$ can be computed as follows:

$$\gcd(60, 24) = \gcd(24, 12) = \gcd(12, 0) = 12.$$





Euclid's algorithm for computing $\gcd(m, n)$

Step 1 If $n=0$, return the value of m as the answer and stop;
otherwise, proceed to Step 2.

Step 2 Divide m by n and assign the value of the remainder to r .

Step 3 Assign the value of n to m and the value of r to n . Go to Step





Alternatively, we can express the same algorithm in pseudocode:



ALGORITHM *Euclid*(m, n)

//Computes $\text{gcd}(m, n)$ by Euclid's algorithm

//Input: Two nonnegative, not-both-zero integers m and n

//Output: Greatest common divisor of m and n

while $n \neq 0$ **do**

$r \leftarrow m \bmod n$

$m \leftarrow n$

$n \leftarrow r$

return m





Consecutive integer checking algorithm



Consecutive integer checking algorithm for computing $\text{gcd}(m, n)$

Step 1 Assign the value of $\min\{m, n\}$ to t .

Step 2 Divide m by t . If the remainder of this division is 0, go to Step 3; otherwise, go to Step 4.

Step 3 Divide n by t . If the remainder of this division is 0, return the value of t as the answer and stop; otherwise, proceed to Step 4.

Step 4 Decrease the value of t by 1. Go to Step 2.





Middle-school procedure



Middle-school procedure for computing $\text{gcd}(m, n)$

Step 1 Find the prime factors of m .

Step 2 Find the prime factors of n .

Step 3 Identify all the common factors in the two prime expansions found in Step 1 and Step 2. (If p is a common factor occurring pm and pn times in m and n , respectively, it should be repeated $\min\{pm, pn\}$ times.)

Step 4 Compute the product of all the common factors and return it as the greatest common divisor of the numbers given.

Thus, for the numbers 60 and 24, we get

$$60 = 2 \cdot 2 \cdot 3 \cdot 5$$

$$24 = 2 \cdot 2 \cdot 2 \cdot 3$$

$$\text{gcd}(60, 24) = 2 \cdot 2 \cdot 3 = 12.$$





sieve of Eratosthenes



As an example, consider the application of the algorithm to finding the list of primes not exceeding $n = 25$:

2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
2	3		5		7		9		11		13		15		17		19		21		23		25
2	3		5		7				11		13				17		19				23		25
2	3		5		7				11		13				17		19				23		



ALGORITHM *Sieve(n)*

//Implements the sieve of Eratosthenes

//Input: A positive integer $n > 1$

//Output: Array L of all prime numbers less than or equal to n

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for  $p \leftarrow 2$  to  $n$  do  $A[p] \leftarrow p$ 
for  $p \leftarrow 2$  to  $\lfloor \sqrt{n} \rfloor$  do //see note before pseudocode
    if  $A[p] \neq 0$  //p hasn't been eliminated on previous passes
         $j \leftarrow p * p$ 
        while  $j \leq n$  do
             $A[j] \leftarrow 0$  //mark element as eliminated
             $j \leftarrow j + p$ 
//copy the remaining elements of  $A$  to array  $L$  of the primes
 $i \leftarrow 0$ 
for  $p \leftarrow 2$  to  $n$  do
    if  $A[p] \neq 0$ 
         $L[i] \leftarrow A[p]$ 
         $i \leftarrow i + 1$ 
return  $L$ 
```



Thank you!

