



# SNS COLLEGE OF TECHNOLOGY

Coimbatore-37.

An Autonomous Institution



**COURSE NAME : 19ITB201 & DESIGN AND ANALYSIS OF ALGORITHMS**

**II YEAR/ IV SEMESTER**

**UNIT – 1 INTRODUCTION**

**Topic: Fundamentals of the Analysis of Algorithm Efficiency**

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# *Fundamentals of the Analysis of Algorithm Efficiency*

- Analysis Framework
- Asymptotic Notations and its properties
- Mathematical analysis for Recursive algorithms.
- Mathematical analysis for Non recursive algorithms.



## *Analysis Framework*

There are two kinds of efficiencies to analyze the efficiency of any algorithm. They are:

- *Time efficiency*, indicating how fast the algorithm runs, and
- *Space efficiency*, indicating how much extra memory it uses

**The algorithm analysis framework consists of the following:**

- Measuring an Input's Size
- Units for Measuring Running Time
- Orders of Growth
- Worst-Case, Best-Case, and Average-Case Efficiencies



# *Fundamentals of the Analysis of Algorithm Efficiency*

## Orders of Growth

- A **difference in running times on small inputs** is not what really distinguishes efficient algorithms from inefficient ones.
- For example, the greatest common divisor of two small numbers, it is not immediately clear how much more efficient Euclid's algorithm is compared to the other algorithms, the difference in **algorithm efficiencies becomes clear** for larger numbers only.



## *Worst-Case, Best-Case, and Average-Case Efficiencies*

Consider Sequential Search algorithm some search key  $K$  **ALGORITHM** *SequentialSearch*( $A[0..n - 1], K$ )

//Searches for a given value in a given array by sequential search

//Input: An array  $A[0..n - 1]$  and a search key  $K$

//Output: The index of the first element in  $A$  that matches  $K$  or -1 if there are no

// matching elements

$i \leftarrow 0$

**while**  $i < n$  **and**  $A[i] \neq K$  **do**

$i \leftarrow i + 1$

**if**  $i < n$  **return**  $i$

**else return** -1

Clearly, the running time of this algorithm can be quite different for the same list size  $n$ .



## *Worst-Case, Best-Case, and Average-Case Efficiencies*

### *Worst-case efficiency:-*

- The *worst-case efficiency* of an algorithm is its efficiency for the worst case input of size  $n$ .
- The algorithm runs the longest among all possible inputs of that size.
- For the input of size  $n$ , the running time is  $C_{worst}(n) = n$ .

### *Best case efficiency*

- The *best-case efficiency* of an algorithm is its efficiency for the best case input of size  $n$ .
- The algorithm runs the fastest among all possible inputs of that size  $n$ .
- In sequential search, If we search a first element in list of size  $n$ . (*i.e.* first element equal to a search key), then the running time is
- $C_{best}(n) = 1$ .



## *Fundamentals of the Analysis of Algorithm Efficiency*

### *Average case efficiency*

- The Average case efficiency lies **between best case and worst case**.
- To analyze the algorithm's average case efficiency, we must make some
- assumptions about possible inputs of size  $n$ .

The standard assumptions are that

- The probability of a successful search is equal to  $p$  ( $0 \leq p \leq 1$ ) and
- The probability of the first match occurring in the  $i$ th position of the list is
- the same for every  $i$ .



## *References*

1. Anany Levitin, “Introduction to the Design and Analysis of Algorithms”, Pearson Education, 3rd Edition, 2012
2. Ellis Horowitz, Sartaj Sahni and Sanguthevar Rajasekaran, “Fundamentals of Computer Algorithms”, Galgotia Publications, 2<sup>nd</sup> edition, 2003