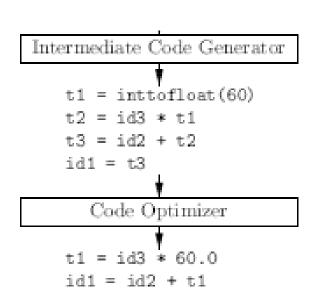




Unit V Intermediate Code Optimization

- Program transformation technique
- Improves code consume less resources
- Transforms the code to make it more efficient
- Output is not changed
- Intermediate code →
 optimization → code
 generation is made easier





Code Optimization



- Optimization
 - Machine Independent Optimization
 - takes in the intermediate code and transforms a part of the code that does not involve any CPU registers
 - Example:

```
do
{
   item = 10;
   value = value + item;
} while(value<100);</pre>
```

This code involves repeated assignment of the identifier item, which if we put this way:

```
Item = 10;
do
{
    value = value + item;
} while(value<100);</pre>
```

- Machine Dependent optimization
 - Target code
 - Rearrangement of machine instructions to improve the efficiency of the code
 - Divide the code into basic blocks





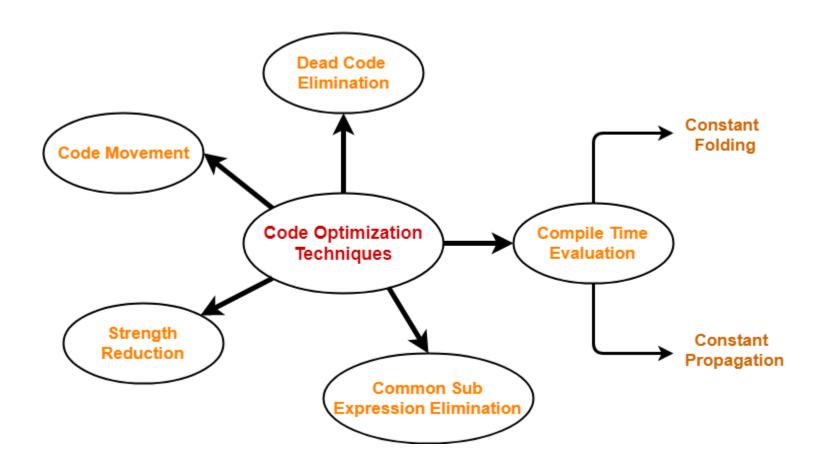
Peephole Optimization

- Optimization eliminates the redundant instruction from a small area of code
- Set of code peephole / window
- Goals:
 - Improves performance
 - Reduce memory footprint
 - Reduce code size





Principle sources of optimization





Compile Time Evaluation



Constant Folding

- Folding the constants
- The expressions that contain the operands having constant values at compile time are evaluated.
- Example:
- return (3+5); \rightarrow return 8;
- $Cir=(22/7)*diameter \rightarrow cir = 3.14*diameter$

Constant Propagation

 If some variable has been assigned some constant value, then it replaces that variable with its constant value in the further program during compilation.

- Example:

- radius =10,pi=3.14
- area=pi*radius*radius; → area=3.14*10*10;





Common Sub Expression

Code before Optimization	Code after Optimization
$S1 = 4 \times i$	$S1 = 4 \times i$
S2 = a[S1]	S2 = a[S1]
$S3 = 4 \times j$ $S4 = 4 \times i // \text{ Red und ant Expression}$	$S3 = 4 \times j$
S5 = n	S5 = n
S6 = b[S4] + S5	S6 = b[S1] + S5





Code Movement

Code before Optimization

for (int j = 0; j < n; j ++)

{

X = Y + Z

 $\mathbf{a}[\mathbf{j}] = 6 \times \mathbf{j}$

Code after Optimization

$$x = v + z$$

for (int
$$j = 0$$
; $j < n$; $j ++$)

{

$$a[j] = 6xj$$
;

}





Dead Code Elimination

• Eliminates the dead code

Code before Optimization

Code after Optimization

```
i = 0;

if (i == 1)

{

a = x + 5;

}
```





Strength Reduction

- Reduces the strength of expressions
- Replaces expensive operators with cheaper one

Example

- B=A*2 →B=A+A
- Cost of multiplication is higher than the addition