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Reg.No

SNS College of Technology, Coimbatore-35.

(Autonomous)

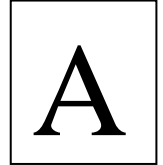
B.E/B.Tech- Internal Assessment -II

Academic Year 2022-2023(ODD)

Fifth Semester

Computer Science and Engineering

19CSB301 – Automata Theory and Compiler Design



Time: 1.5 Hours

Maximum Marks: 50

Part-A (5 x 2 =10)

		CO	Blooms
1.	Name the rules of computing the Firstpos and Lastpos of concatenation and union operation	CO2	Rem
2.	Construct the Syntax tree for $((a+b)^*(a.c)^*)$ and mark the nullable nodes	CO2	Cre
3.	Define LL (1) grammar and its properties	CO3	Rem
4.	Find the first and follow of the CFG given below: S → ABCD A → a ε B → b ε C → c D → d ε	CO3	App
5.	Define Type checking	CO3	Rem

Part-B (2x13+14=40)

6.	a. Construct the ε-NFA to DFA for the given regular expression $(a b)^*abb$	13	CO2	Cre
	or			
	b. Construct DFA from the given regular expression “ $ba(a+b)^*ab$ ” using Direct Method	13	CO2	Cre
7.	a. Check whether the given CFG is LL(1) or not. Check whether the string $id+id*id$ is accepted by this CFG E → TE' E' → +TE' ε T → FT' T' → *FT' ε F → id (E)	13	CO3	Ana

or

- b. Construct the SLR parsing table for the following grammar. check whether the string (a) is accepted or not. $A \rightarrow (A)a$ 13 CO3 Eva
8. a. Construct the DFA from the given Regular Expression $(a|b)^*abb$ using Direct Method 14 CO2 Cre
- or
- b. Construct the canonical parsing table for the grammar given below. the check whether the string "cdcd" is accepted or not. 14 CO3 Eva

$S \rightarrow CC$
 $C \rightarrow cC$
 $C \rightarrow d$

Und-Understanding Rem-Remembering App-Appling
Ana-Analysis Cre-Creating Eva-Evaluating

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Fifth Semester

Computer Science and Engineering

19CSB301 – Automata Theory and Compiler Design

B

Time: 1.5 Hours

Maximum Marks: 50

Part-A (5 x 2 =10)

	CO	Bloom s
1. Draw the ϵ -NFA for the regular expression $(0 1)^*$	CO2	Rem
2. Construct the Syntax tree for $(a b)^*a$ and find the first pos and last pos of each node	CO2	App
3. Eliminate immediate left recursion for the following Grammar E \rightarrow E+T T \rightarrow T*F F \rightarrow (E) id	CO3	App
4. Perform left factoring for the grammar. S \rightarrow iEtS / iEtSeS / a	CO3	APP
5. Mention the two rules for Type Checking	CO3	Rem

Part-B (2x13+14=40)

6. a. Construct the ϵ -NFA to DFA for the given regular expression $(a|b)^*a(a|b)$ using Thompsons Construction 13 CO2 Cre
- or
- b. Convert the Regular Expression $(a|b)^*abb$ to DFA (Direct Method) 13 CO2 Cre

7. a. Construct the predictive parsing table for the following grammar and hence check whether the string (a,a) is accepted or not.

$$S \rightarrow (L)|a$$
$$L \rightarrow L,S|S$$

or

- b. Construct the CLR parsing table for the following grammar. check whether the string (a) is accepted or not.
- CFG $A \rightarrow (A)|a$

8. a. Construct the DFA from the given Regular Expression (a/b)*a (a/b) using Direct Method

or

- b. Construct the SLR parsing table for the following grammar and check whether the string “adad” is accepted or not.

$$S \rightarrow CC$$
$$C \rightarrow aC$$
$$C \rightarrow d$$

**Und-Understanding Rem-Remembering App-Appling
Ana-Analysis Cre-Creating Eva-Evaluating**



Reg.No

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Time: 1.5 Hours

Maximum Marks: 50

Part-A (5 x 2 =10)

		CO	Blooms								
1.	<p>Give the rules of computing the Firstpos and Lastpos of concatenation and union operation</p> <table border="1"> <tr> <td></td> <td><i>nullable(c₁)</i> or <i>nullable(c₂)</i></td> <td><i>firstpos(c₁) ∪ firstpos(c₂)</i></td> <td><i>lastpos(c₁) ∪ lastpos(c₂)</i></td> </tr> <tr> <td></td> <td><i>nullable(c₁)</i> and <i>nullable(c₂)</i></td> <td>if (<i>nullable(c₁)</i>) <i>firstpos(c₁) ∪ firstpos(c₂)</i> else <i>firstpos(c₁)</i></td> <td>if (<i>nullable(c₂)</i>) <i>lastpos(c₁) ∪ lastpos(c₂)</i> else <i>lastpos(c₂)</i></td> </tr> </table>		<i>nullable(c₁)</i> or <i>nullable(c₂)</i>	<i>firstpos(c₁) ∪ firstpos(c₂)</i>	<i>lastpos(c₁) ∪ lastpos(c₂)</i>		<i>nullable(c₁)</i> and <i>nullable(c₂)</i>	if (<i>nullable(c₁)</i>) <i>firstpos(c₁) ∪ firstpos(c₂)</i> else <i>firstpos(c₁)</i>	if (<i>nullable(c₂)</i>) <i>lastpos(c₁) ∪ lastpos(c₂)</i> else <i>lastpos(c₂)</i>	CO2	Rem
	<i>nullable(c₁)</i> or <i>nullable(c₂)</i>	<i>firstpos(c₁) ∪ firstpos(c₂)</i>	<i>lastpos(c₁) ∪ lastpos(c₂)</i>								
	<i>nullable(c₁)</i> and <i>nullable(c₂)</i>	if (<i>nullable(c₁)</i>) <i>firstpos(c₁) ∪ firstpos(c₂)</i> else <i>firstpos(c₁)</i>	if (<i>nullable(c₂)</i>) <i>lastpos(c₁) ∪ lastpos(c₂)</i> else <i>lastpos(c₂)</i>								
2.	<p>Construct the Syntax tree for ((a+b)*+(a.c)*) and mark the nullable nodes</p>	CO2	Und								
3.	<p>Define LL (1) grammar and its properties</p> <p>LL(1), the first L stands for scanning the input from left to right, the second L stands for producing a leftmost derivation, and the 1 stands for using one input symbol of lookahead at each step to make parsing action decision.</p>	CO3	Und								

4.	Find the first and follow of the CFG given below:	CO3	Ana																		
	<table border="1"> <thead> <tr> <th></th> <th>FIRST</th> <th>FOLLOW</th> </tr> </thead> <tbody> <tr> <td>$S \rightarrow ABCD$</td> <td>{a,b,c}</td> <td>{\\$}</td> </tr> <tr> <td>$A \rightarrow a \epsilon$</td> <td>{a, ϵ}</td> <td>{b,c}</td> </tr> <tr> <td>$B \rightarrow b \epsilon$</td> <td>{b, ϵ}</td> <td>{c}</td> </tr> <tr> <td>$C \rightarrow c$</td> <td>{c}</td> <td>{d, \\$}</td> </tr> <tr> <td>$D \rightarrow d \epsilon$</td> <td>{d, ϵ}</td> <td>{S}</td> </tr> </tbody> </table>		FIRST	FOLLOW	$S \rightarrow ABCD$	{a,b,c}	{\\$}	$A \rightarrow a \epsilon$	{a, ϵ }	{b,c}	$B \rightarrow b \epsilon$	{b, ϵ }	{c}	$C \rightarrow c$	{c}	{d, \\$}	$D \rightarrow d \epsilon$	{d, ϵ }	{S}		
	FIRST	FOLLOW																			
$S \rightarrow ABCD$	{a,b,c}	{\\$}																			
$A \rightarrow a \epsilon$	{a, ϵ }	{b,c}																			
$B \rightarrow b \epsilon$	{b, ϵ }	{c}																			
$C \rightarrow c$	{c}	{d, \\$}																			
$D \rightarrow d \epsilon$	{d, ϵ }	{S}																			
5.	Define Type checking Type checking is the process of verifying and enforcing constraints of types in values. A compiler must check that the source program should follow the syntactic and semantic conventions of the source language and it should also check the type rules of the language.	CO3	Und																		

Part-B (2x13+14=40)

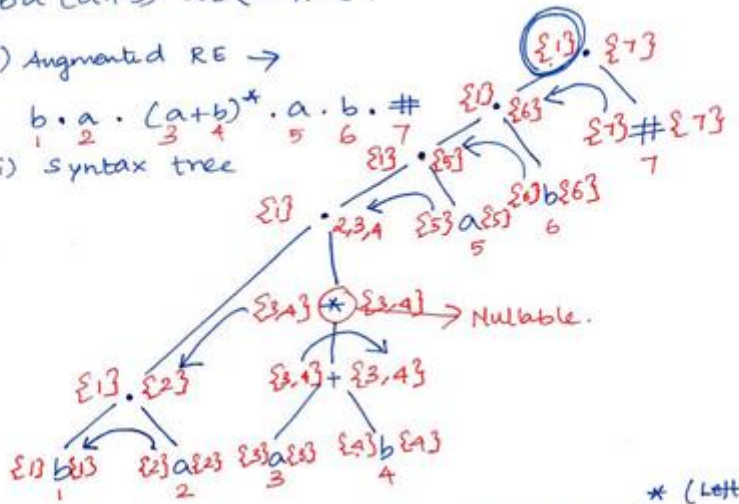
6.	a.	<p>Convert the ϵ-NFA to DFA for the given regular expression $(a b)^*abb$</p> <p><i>f.2021</i> <u>ϵ-NFA to DFA</u> 1. a^+b^+c</p> <p>ϵ-closure(q_0) = {q_0, q_1, q_2} = A</p> <p>A \xrightarrow{a} q_0 = ϵ-closure(q_0) = A</p> <p>A \xrightarrow{b} q_1 = ϵ-closure(q_1) = {q_1, q_2} = B</p> <p>A \xrightarrow{c} q_2 = ϵ-closure(q_2) = {q_2} = C</p> <p>B \xrightarrow{a} ϕ = ϵ-closure(ϕ) = \emptyset (Non-Final state)</p> <p>B \xrightarrow{b} q_1 = ϵ-closure(q_1) = B</p> <p>B \xrightarrow{c} q_2 = ϵ-closure(q_2) = C</p> <p>C \xrightarrow{a} ϕ = \emptyset</p> <p>C \xrightarrow{b} ϕ = \emptyset</p> <p>C \xrightarrow{c} q_2 = C</p> <p>D (Non-Final state) X</p> <table border="1"> <thead> <tr> <th></th> <th>a</th> <th>b</th> <th>c</th> </tr> </thead> <tbody> <tr> <th>A</th> <td>A</td> <td>B</td> <td>C</td> </tr> <tr> <th>B</th> <td>\emptyset</td> <td>B</td> <td>C</td> </tr> <tr> <th>*C</th> <td>\emptyset</td> <td>\emptyset</td> <td>C</td> </tr> <tr> <th>\emptyset</th> <td>\emptyset</td> <td>\emptyset</td> <td>\emptyset</td> </tr> </tbody> </table> <p>DFA</p>		a	b	c	A	A	B	C	B	\emptyset	B	C	*C	\emptyset	\emptyset	C	\emptyset	\emptyset	\emptyset	\emptyset	13	CO2	Und
	a	b	c																						
A	A	B	C																						
B	\emptyset	B	C																						
*C	\emptyset	\emptyset	C																						
\emptyset	\emptyset	\emptyset	\emptyset																						
		or																							
	b.	Convert the given regular expression "ba(a+b)*ab" to DFA (Direct Method)	5	CO2	App																				

1. $ba(a+b)^*ab \leftarrow R.E.$

i) Augmented RE \rightarrow

$b \cdot a \cdot (a+b)^* \cdot a \cdot b \cdot \#$
 1 2 3 4 5 6 7

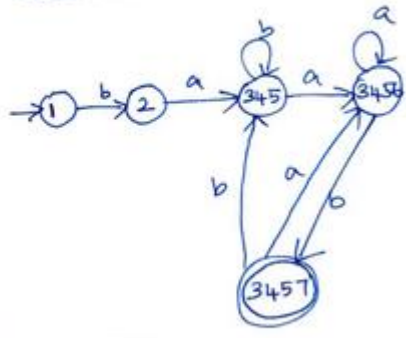
ii) syntax tree



iii) Nullable, Firstpos, Lastpos, Followup
 $\epsilon, *$ operand, operator

* (Left to Right)
 Followup (*, .)
 . (Right to left)

		Followup
b	1	2
a	2	3, 4, 5
a	3	3, 4, 5
b	4	3, 4, 5
a	5	6
b	6	7
#	7	-



DFA
 $ba(a+b)^*ab.$

9 CO2 Und

7. a. Check whether the given CFG is LL(1) or not. Check whether the string $id+id*id$ is accepted by this CFG
 $E \rightarrow TE'$
 $E' \rightarrow +TE' | \epsilon$
 $T \rightarrow FT'$
 $T' \rightarrow *FT' | \epsilon$
 $F \rightarrow id | (E)$

13 CO3 App

LL(1) parsing Table.

① Example 1. (LSE1) $id + id * id \$$

	id	+	*	()	\$
E	$E \rightarrow TE'$			$E \rightarrow TE'$		
E'		$E' \rightarrow +TE'$			$E' \rightarrow \epsilon$	$E' \rightarrow \epsilon$
T	$T \rightarrow FT'$			$T \rightarrow FT'$		
T'		$T' \rightarrow \epsilon$	$T' \rightarrow *FT'$		$T' \rightarrow \epsilon$	$T' \rightarrow \epsilon$
F	$F \rightarrow id$			$F \rightarrow (E)$		

It is an LL(1)

STACK	INPUT	OUTPUT
\$E	id + id * id \$	
\$E'T	id + id * id \$	$E \rightarrow TE'$
\$E'T'F	id + id * id \$	$T \rightarrow FT'$
\$E'T'id	id + id * id \$	$F \rightarrow id$
\$E'T'	+ id * id \$	pop id
\$E'	+ id * id \$	$T' \rightarrow \epsilon$
\$E'T+	+ id * id \$	$E' \rightarrow +TE'$
\$E'T	id * id \$	pop +
\$E'T'F	id * id \$	$T \rightarrow FT'$
\$E'T'id	* id \$	$F \rightarrow id$, pop id
\$E'T'F*	* id \$	$T' \rightarrow *FT'$
\$E'T'id	id \$	pop *, $F \rightarrow id$
\$E'	\$	pop id, $T' \rightarrow \epsilon$
	\$	$E \rightarrow \epsilon$

OR

b. Construct the SLR parsing table for the following grammar. check whether the string (a) is accepted or not
 $A \rightarrow (A)|a$

13

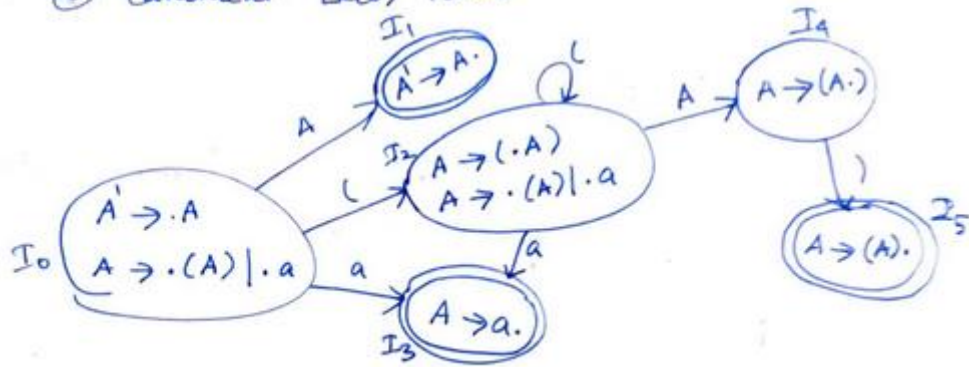
CO3

Und

① Augmented Grammar

$$\begin{aligned} \bar{A} &\rightarrow A \\ A &\rightarrow (A) | a \end{aligned}$$

② Canonical LR(0) items



SLR Table

Final Items - Reduce operation (Follow of ~~first~~ of LHS)
 $A \rightarrow (A) | a$, FIRST $\{ (, a \}$ | FOLLOW $\{), \$ \}$

ITEM	ACTION				GOTO
	a	()	\$	
I ₀	S ₃	S ₂			1
I ₁				Accept	
I ₂	S ₃	S ₂			4
I ₃			r ₂	r ₂	
I ₄			S ₅		
I ₅			r ₁	r ₁	

8. a. Convert the Regular Expression $(a|b)^*abb$ to DFA using Direct Method

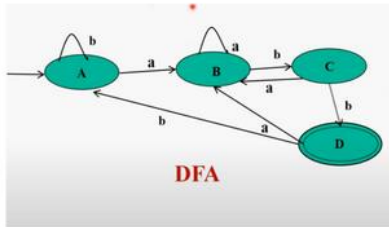
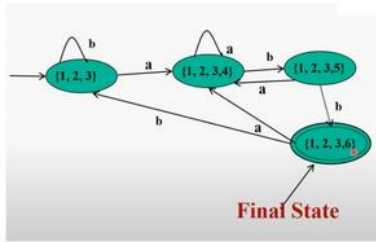
14

CO2

Ana

DFA Construction

	Node	followpos
a	1	1,2,3
b	2	1,2,3
a	3	4
b	4	5
b	5	6
#	6	-



OR

b. Construct the canonical parsing table for the grammar given below. the check whether the string “cdcd “ is accepted or not.

$S \rightarrow CC$

$C \rightarrow cC$

$C \rightarrow d$

STEP 3:
CLR parse Table

	ACTION			GOTO	
	a	b	\$	A	S
I_0	S_3	S_4		2	1
I_1			Accept		
I_2	S_6	S_7		5	
I_3	S_3	S_4		9	
I_4	r_3	r_3			
I_5			r_1		
I_6	S_6	S_7		8	
I_7			r_3		
I_8			r_2		
I_9	r_2	r_2			

14

CO3

Eva

Reg.No

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Answer Key

B

Time: 1.5 Hours

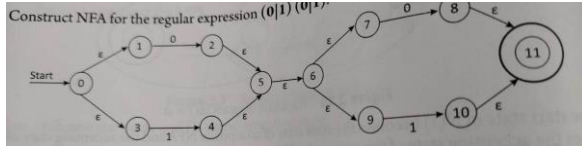
Maximum Marks: 50

Part-A (5 x 2 =10)

CO Blooms

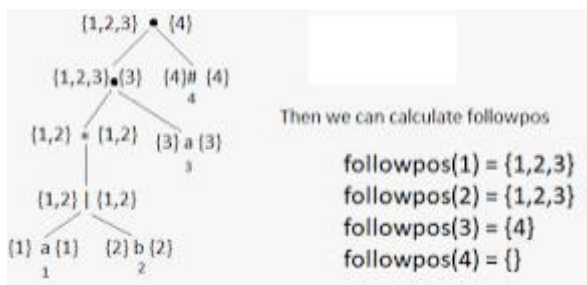
1. Draw the ϵ -NFA for the regular expression $(0|1)(0|1)$

CO Rem 2



2. Construct the Syntax tree for $(a|b)^*a$ and find the first pos and last pos of each node

CO App 2



3. Eliminate immediate left recursion for the following Grammar

CO App 3

$E \rightarrow E+T$

$T \rightarrow T * F$

$F \rightarrow (E) | id$

Eliminating the immediate left recursion from the productions for E and then for T, we obtain

$E \rightarrow TE'$

$E' \rightarrow + TE' | \epsilon$

$T \rightarrow FT'$

$T' \rightarrow * FT' | \epsilon$

$F \rightarrow (E) | id.$

4. Perform left factoring for the grammar.
 $S \rightarrow iEtS / iEtSeS / a$

CO APP
3

5. Mention the two rules for Type Checking

CO Rem
3

- Type Synthesis

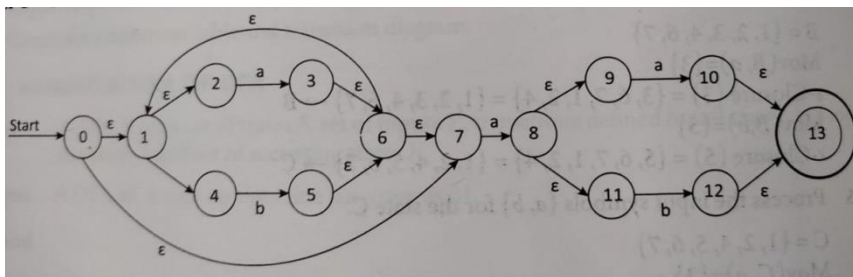
- Builds the type of an expression from the types of its subexpressions
- Requires names to be declared before usage

- Type inference

- determines the type of a construct from the way it is used

Part-B (2x13+14=40)

6. a. Construct the ϵ -NFA to DFA for the given regular expression $(a|b)^*a(a|b)$ using Thompson's Construction 13 CO2



States	Input	
	a	b
A	B	C
B	D	E
C	B	C
D	D	E
E	B	C

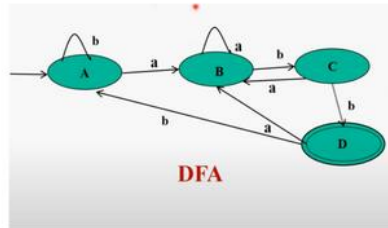
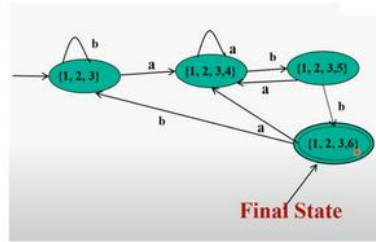
or

- b. Convert the Regular Expression $(a|b)^*abb$ to DFA (Direct Method)

13 CO2

DFA Construction

	Node	followpos
a	1	1,2,3
b	2	1,2,3
a	3	4
b	4	5
b	5	6
#	6	-



7. a. Construct the predictive parsing table for the following grammar and hence check whether the string (a,a) is accepted or not. 13 CO3

$$S \rightarrow (L)a$$

$$L \rightarrow L,S|S$$

Remove left recursion $L \rightarrow SL'$

$$L' \rightarrow \epsilon, SL' / \epsilon$$

The grammar

$$S \rightarrow (L)$$

$$S \rightarrow a$$

$$L \rightarrow SL'$$

$$L' \rightarrow \epsilon, SL' / \epsilon$$

Find FIRST for all the non-terminals.

$$\text{Non-terminals} = \{S, L\}$$

$$\text{Terminals} = \{ (,), a \}$$

$$\text{FIRST}(S) = \{ (, a \}$$

$$\text{FIRST}(L) = \{ \text{FIRST}(S) \} = \{ (, a \}$$

$$\text{FIRST}(L') = \{ \epsilon \}$$

Find FOLLOW for all the non-terminals.

$$\begin{aligned} \text{FOLLOW}(S) &= \{ S \text{ followed by non-terminal } L', \$ \} \\ &= \{ \text{FIRST}(L') = \{ \epsilon \}, \text{ eliminate } \epsilon, \text{ find FOLLOW}(L) \} \\ &= \{), \$ \} \end{aligned}$$

$$\text{FOLLOW}(L) = \{ \}$$

$$\begin{aligned} \text{FOLLOW}(L') &= \{ L' \text{ does not followed by terminal and non-terminal} \} \\ &= \{ \text{FOLLOW}(L) \} \\ &= \{ \} \end{aligned}$$

	()	a	,	\$
S	$S \rightarrow (L)$		$S \rightarrow a$		
L	$L \rightarrow SL'$		$L \rightarrow SL'$		
L'		$L' \rightarrow \epsilon$		$L' \rightarrow ,SL'$	

The above string is accepted

or

- b. Construct the CLR parsing table for the following grammar. check whether the string (a) is accepted or not. 13 CO3

CFG $A \rightarrow (A)a$

Step 1: Write the augmented Grammar

Step 2: Compute Closure function

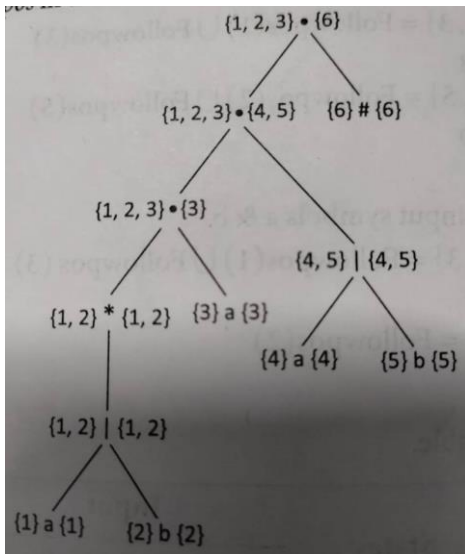
Step 3: Find LR(0) items

Step 4: Find Goto function

Step 5: CLR Parsing Table construction

Step 6: Input string Parsing

8. a. Construct the DFA from the given Regular Expression $(a/b)^*a(a/b)$ using Direct Method 14 CO2



Nodes	followpos
1	{1, 2, 3}
2	{1, 2, 3}
3	{4, 5}
4	{6}
5	{6}
6	—

or

- b. Construct the SLR parsing table for the following grammar and check whether the string “adad” is accepted or not. 14 CO3

$S \rightarrow CC$

$C \rightarrow aC$

$C \rightarrow d$

Step 1 : Find the augmented grammar G

Step 2: Find LR(0) Items

Step 3: Construct Parsing Table

Step 4: Parse the string

States	Action			goto	
	c	d	\$	S	C
0	s3	s4		1	2
1			acc		
2	s3	s4			5
3	s3	s4			6
4	r3	r3	r3		
5	r1	r1	r1		
6	r2	r2	r2		

Stack	Input string	Action
0	cdcd \$	Shift 3
0c3	dcd \$	Shift 4
0c3d4	cd\$	Reduce r3 C → d
0c3C6	cd\$	Reduce r2 C → cC
0C2	cd\$	Shift 3
0C2c3	d\$	Shift 4
0C2c3d4	\$	Reduce r3 C → d
0C2c3C6	\$	Reduce r1 S → CC
0C2c3C6	\$	Reduce r2 C → cC
0C2C5	\$	Reduce r1 S → CC
0S1	\$	Accept

Und-Understanding Rem-Remembering App-Aplying
Ana-Analysis Cre-Creating Eva-Evaluating

Prepared By

Verified By

HoD