

SNS COLLEGE OF TECHNOLOGY



Coimbatore-35
An Autonomous Institution

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DEPARTMENT OF ELECTRONICS & COMMUNICATION ENGINEERING

19ECE351 – IMAGE PROCESSING AND COMPUTER VISION

III B.E. ECE / V SEMESTER

UNIT 3 – IMAGE COMPRESSION AND IMAGE SEGMENTATION

TOPIC - INTRODUCTION





- The most popular technique for removing coding redundancy is due to Huffman (1952)
- When coding the symbols of an information source individually, Huffman coding yields the smallest possible number of code symbols per source symbol
- In terms of the noiseless coding theorem, the resulting code is optimal for a fixed value of n, subject to the constraint that the source symbols be coded one at a time

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The first step in Huffman's approach is

- To create a series of source reductions by ordering the probabilities of the symbols under consideration
- Combining the lowest probability symbols into a single symbol that replaces them in the next source reduction





Huffman Coding

Example: Calculate the Huffman Codes for the set of symbols as shown in table.

Probability 0.4 0.3 0.2 0.1	
)
Solution: $\begin{array}{cccccccccccccccccccccccccccccccccccc$	1 code
Symbol Probability Symbol Probability	
A 0.4 P.O.	
B 0.3 0.3 larg to represent r	wed
$= Z l(\mathcal{R}_{K}) P(\mathcal{R}_{K})$	nessage
$= 1 \times 0.4 + 2 \times 0.3 + 3 \times 0.2 + 9$ $= 1.9 \text{ bits by mbots}$	X0.)





HUFF MANN CODING 1. To Find the average code word Length L = 5 PK [length of mk in bits] 2. To Find the Entropy of the Source H = 5 PK 1092 (PK) 3. code efficiency $\eta = \frac{H \times 100\%}{L}$







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