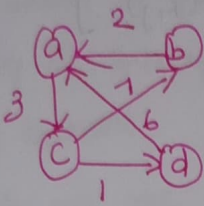


## Floyd's algorithm :

This is for finding the shortest path between every pair of vertices in a given graph.

This is for both undirected and directed graph.  
Also called "All pair shortest path".

Ex:



$$w = \begin{matrix} & a & b & c & d \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 2 & 3 & 1 \\ 2 & 0 & \infty & 6 \\ \infty & 7 & 0 & 1 \\ 6 & \infty & \infty & 0 \end{bmatrix} \end{matrix} \rightarrow \text{weight matrix.}$$

$$D[0] = w \quad D[1] = ? \quad a \quad D[2] = ? \quad a, b \quad D[3] = ? \quad a, b, c$$

$$D[4] = a, b, c, d.$$

↓ intermediate nodes.

$$D[0] = \begin{bmatrix} 0 & 2 & 3 & 1 \\ 2 & 0 & \infty & 6 \\ \infty & 7 & 0 & 1 \\ 6 & \infty & \infty & 0 \end{bmatrix}$$

$$D[1] = \begin{matrix} & a & b & c & d \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} a & 2 & 3 & 1 \\ b & 2 & 0 & 5 \\ c & 2 & 7 & 0 & 1 \\ d & 6 & 2 & 9 & 0 \end{bmatrix} \end{matrix}$$

with intermediate 'a'

$$D[2] = \begin{matrix} & a & b & c & d \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} a & 2 & 3 & 1 \\ b & 2 & 0 & 5 \\ c & 9 & 7 & 0 & 1 \\ d & 6 & 2 & 9 & 0 \end{bmatrix} \end{matrix}$$

with intermediate 'a' and 'b'

$$D[3] = \begin{matrix} & a & b & c & d \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} a & 10 & 3 & 4 \\ b & 2 & 0 & 5 & 6 \\ c & 9 & 7 & 0 & 1 \\ d & 6 & 16 & 9 & 0 \end{bmatrix} \end{matrix}$$

with intermediate 'a', 'b' and 'c'

$$D[4] = \begin{matrix} & a & b & c & d \\ a & \underline{0} & \underline{10} & \underline{3} & \underline{4} \\ b & \underline{2} & \underline{0} & \underline{5} & \underline{6} \\ c & \underline{7} & \underline{1} & \underline{0} & \underline{1} \\ d & \underline{6} & \underline{16} & \underline{9} & \underline{0} \end{matrix}$$

$$D^k(i,j) = \text{Min} [ D^{k-1}(i,j), D^{k-1}(i,k) + D^{k-1}(k,j) ]$$

Algo:

```

p ← w
for k ← 1 to n do
{
  for i ← 1 to n do
  {
    for j ← 1 to n do

```

$$D[i,j] = \text{min} \{ D[i,j], D[i,k] + D[k,j] \}$$

}  
 }  
 }  
 return D.

Analysis:

$$\sum_{k=1}^n \sum_{j=1}^n \sum_{i=1}^n 1$$

$$= \sum_{k=1}^n n^2$$

$$c(n) = n^3$$

$$\theta(n^3)$$