

# Algorithm for backtracking recursive

{ for (each  $a_k$  that belongs to  $T(a_1, a_2, \dots, a_{k-1})$ ) do

{ if ( $B_k(a_1, a_2, \dots, a_k) = \text{true}$ ) then // feasible sequence

{ if ( $(a_1, a_2, \dots, a_k)$  is path to answer node then  
print ( $a[1], a[2], \dots, a[k]$ ))

if ( $k < n$ ) then

Backtrack ( $k+1$ ); // find next set

}

non-recursive

{  $k := 1$ ;

while ( $k \neq 0$ ) do

{ if (any  $a[k]$  that belongs to  $T(a[1], a[2], \dots, a[k-1])$   
remains untried)

AND ( $B_k(a[1], a[2], \dots, a[k])$  is true) then

{ if ( $(a[1], a[2], \dots, a[k])$  is a path to answer node) then

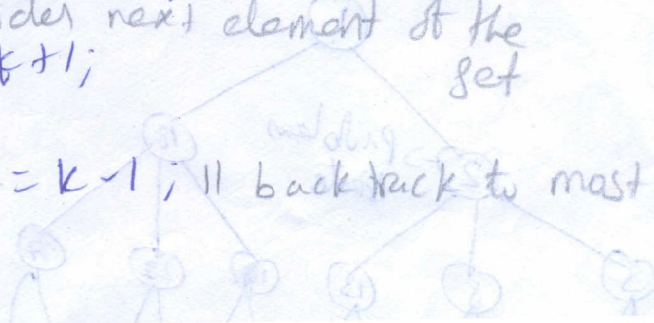
write ( $a[1], a[2], \dots, a[k]$ ); // solution printed

// consider next element of the set  
 $k := k+1$ ;

else

$k := k-1$ ; // backtrack to most recent value

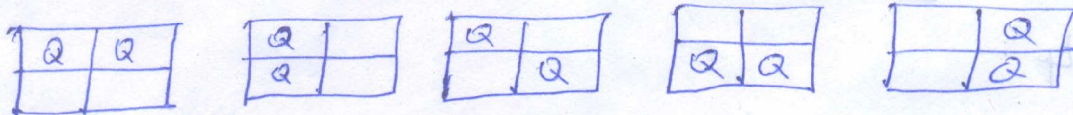
}



Consider  $n \times n$  chessboard on which we have to place  $n$  queens so that no two queens attack each other by being in the same row or in the same column or on the same diagonal.

Examples-

i) Consider  $2 \times 2$  board



\* 2-Queen's problem is not solvable,

\* Because 2-Queen's can be placed on  $2 \times 2$  chessboard

ii) Consider  $4 \times 4$  board

\* Let us take 4-Queen's and  $4 \times 4$  chessboard.

Step 1:

→ now we start with empty chessboard

→ place queen 1 in the 1<sup>st</sup> possible solution of its row.

→ i.e., on 1<sup>st</sup> row & 1<sup>st</sup> column.



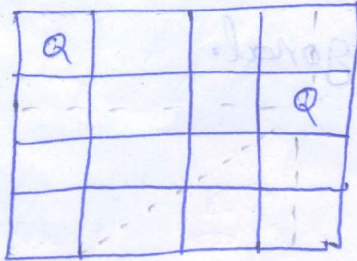
Step 2:

→ Then place queen 2 after trying various unsuccessful

place -  $(1, 2)$ ,  $(2, 1)$ ,  $(2, 2)$  & at  $(2, 3)$  i.e., 2<sup>nd</sup> row & 3<sup>rd</sup> column.



Then place 3<sup>rd</sup> queen cannot be placed on next column, as there is no acceptable position for queens, hence algorithm backtracks & places 2<sup>nd</sup> queen at (2,4) position.



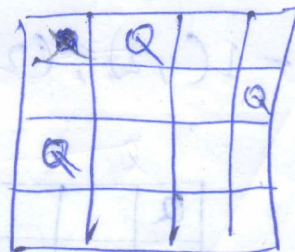
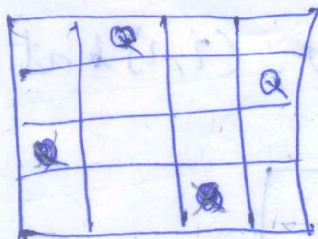
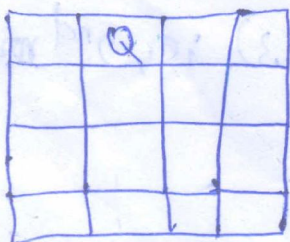
Step: 4

The place 3<sup>rd</sup> queen at (3,2) but it is again another dead end as next queen (4<sup>th</sup> queen) cannot be placed at permissible position.



Step: 5

hence we need to backtrack all the way upto queen 1 & move it to (1,2). place queen 1 at (1,2), queen 2 at (2,4), queen 3 at (3,1) & queen 4 at (4,3)



Algorithm

N queen (k, n)

```

{
  for i = 1 to n do
  {
    if (place (row, column, i)) then // This function checks if
    {
      board [row] [column] → no conflicts so place queen
      if (row = n) then → dead end
      Print board (n) → print the board configuration
      else → try next queen with next position
      queen (row + 1, n)
    }
  }
}

```

place (row, column)

```

{
  for i ← 1 to row - 1 do
  {
    if (board [i] = column) then
    {
      return 0 // Same column by 2 queens
    }
    else if (abs (board [i] - column) = abs (i - row)) then
    {
      return 0 // This formula gives 2 queens then
    }
    → no conflicts hence queen can be placed.
    return 1
  }
}

```