

$i=0$	0	0	0	0	0	0
$i=1$	0	10	12	12	12	12
$i=2$	0	10	22	22	22	22
$i=3$	0					
$i=4$	0					

$$v(1,1) = \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \}$$

$$= \max \{ v(0,1), 12 + v[0, 1-2] \}$$

$$= \max \{ v(0,1), 12 + v[0, -1] \}$$

$$= \max \{ 0, 12 + v[0] \}$$

$$= 12$$

$$v(1,2) = \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \}$$

$$= \max \{ v(0,2), 12 + v[0, 0] \}$$

$$= \max \{ 0, 12 + 0 \}$$

$$= 12$$

$$v(1,3) = \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \}$$

$$= \max \{ v(0,3), 12 + v(0,1) \}$$

$$= \max \{ 0, 12 + 0 \}$$

$$= 12$$

$$v(1,4) = \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \}$$

$$= \max \{ v(0,4), 12 + v(0,2) \}$$

$$= \max \{ 0, 12 + 0 \}$$

$$= 12$$

Dynamic programming & memory optimization

Given n items of known weights w_1, \dots, w_n & values v_1, \dots, v_n & knapsack capacity W , find most valuable subset of items that fill into knapsack.

DP algorithm, derive recurrence relation item $i, 1 \leq i$

weights w_1, \dots, w_i value v_1, \dots, v_i & knapsack capacity $1 \leq j \leq W$

let $v(i, j)$ value of optimal solution

1) among subsets that do not include i th item, value of an optimal subset is by definition $v(i-1, j)$.

2) Among subsets that do include i th item, an opt subset is made up of this item & optimal subset $i-1$ items fit into knapsack capacity $j-w_i$. value of optimal subset is $v_i + v(i-1, j-w_i)$

$$v(i, j) = \begin{cases} \max \{ v(i-1, j), v_i + v(i-1, j-w_i) \} & \text{if } j-w_i \geq 0 \\ v(i-1, j) & \text{if } j-w_i < 0 \end{cases}$$

	0	$j-w_i$	j	w	
0	0	0	0	0	
$i-1$	0	$v(i-1, j-w_i)$	$v(i-1, j)$		
i	0		$v(i, j)$		
n	0			goal.	

	v_i	w_i
1	15	1
2	10	5
3	9	3
4	5	4

Table for filling knapsack

1 2	22	20			
	=3				
3	5	22	item	weight	value
4	4	27	1	2	\$12
2 3	4	30	2	1	\$10
2 1	3	25	3	3	\$20
1 2 3	6		4	2	\$15

Capacity $W=5$

$i=0$	0	0	0	0	0	0
$i=1$	0	10	12	12	12	12
$i=2$	0	10	22	22	22	22
$i=3$	0					
$i=4$	0					

$$\begin{aligned}
 v(1,1) &= \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \} \\
 &= \max \{ v(0,1), 12 + v[0, -1] \} \\
 &= \max \{ 0, 12 + v[0] \} \\
 &= 12
 \end{aligned}$$

$$\begin{aligned}
 v(1,2) &= \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \} \\
 &= \max \{ v(0,2), 12 + v[0,0] \} \\
 &= \max \{ 0, 12 + 0 \} \\
 &= 12
 \end{aligned}$$

$$\begin{aligned}
 v(1,3) &= \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \} \\
 &= \max \{ v(0,3), 12 + v(0,1) \} \\
 &= \max \{ 0, 12 + 0 \} \\
 &= 12
 \end{aligned}$$

$$\begin{aligned}
 v(1,4) &= \max \{ v(i-1, j), w_i + v[i-1, j-w_i] \} \\
 &= \max \{ v(0,4), 12 + v(0,2) \}
 \end{aligned}$$