



# Knapsack using Dynamic programming



# Knapsack problem



## Given:

- n items
- Known weight-  $w_1, w_2, \dots, w_n$
- values=-  $v_1, v_2, \dots, v_n$
- Knapsack capacity  $W$

## Find:

Most valuable subset of items to fill knapsack

## Formula:

$V(i, j)$

Item  $i$ ,  $1 \leq i \leq n$   $i-1$  to  $n$

Weight  $w_1, w_2, \dots, w_i$

Value  $v_1, v_2, \dots, v_i$

Knapsack capacity  $j$   $1 \leq j \leq w$



# Condition



- Subset- ith item doesnot include  $v[i-1,j]$
- Subset- ith item include  $v_i+v[i-1,j-w_i]$

$$V[i,j]= \max\{v[i-1,j], v_i+v[i-1,j-w_i]\}$$



# Formula

$$V[i,j]=\max\{v[l,j-1],v_i+v[i-1,j-w_i]\}$$



	0	J-wi	J	W
0	0	0	0	0
i-1	0	V[i-1,j-wi]	V[i-1,j]	
i	0		V[i,j]	
N	0			Goal



# Algorithm

Knapsack(i,j)

Description : implement the knapsack problem

I/P : non negative integer I & j

O/P : optimal feasible subset

if  $v(i,j) < 0$

if  $j < w[i]$

value  $\leftarrow$  knapsack(i-1,j)

else

value  $\leftarrow$  knapsack(i-1,j), value(i)+ knapsack(i-1,j-wi)

$V(i,j) \leftarrow$  value

Return  $v(i,j)$



# Problem



- Capacity  $W=5$

Items	Weight	Value
1	2	\$12
2	1	\$10
3	3	\$20
4	2	\$15

# Answer

0	0	0	0	0
0	0	12	12	12
0	10	12	22	22
0	10	12	22	30
0	10	15	25	30