

Introduction

- **Introduction**

- What Operating Systems Do
- Computer-System Architecture
- Operating-System Structure
- Operating-System Operations
- Operating-System Services
- User Operating System Interface
- System Calls
- Types of System Calls
- System Programs
- Operating System Structure
- System Boot

- **Process Concept**

- Process Scheduling
- Operations on Processes
- Interprocess Communication

What is an Operating System?

- A program that acts as an intermediary between a **user of a computer** and the **computer hardware**
- **Operating system goals:**
 - Execute user programs and make solving user problems easier
 - Make the computer system convenient to use
 - Use the computer hardware in an efficient manner

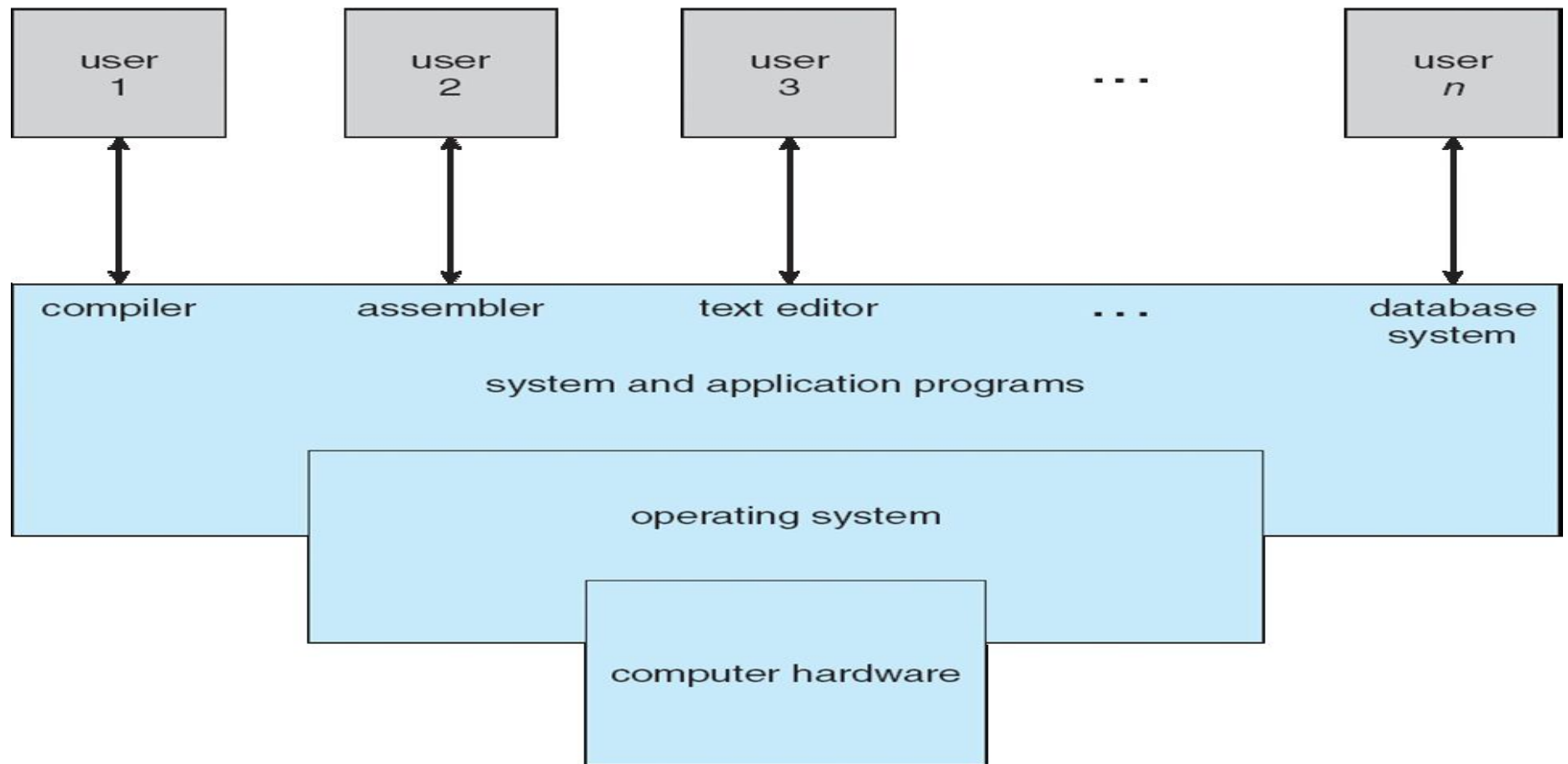


Computer System Structure

- Computer system can be divided into **four components**:
 - **Hardware** – provides basic computing resources
 - CPU, memory, I/O devices
 - **Operating system**
 - Controls and coordinates use of hardware among various applications and users
 - **Application programs** – define the ways in which the system resources are used to solve the computing problems of the users
 - Word processors, compilers, web browsers, database systems, video games
 - **Users**
 - People, machines, other computers



Four Components of a Computer System



WHAT OPERATING SYSTEMS DO

Users want convenience, **ease of use** and **good performance**

- Don't care about **resource utilization**

But shared computer such as **mainframe** or **minicomputer** must keep all users happy

Users of dedicate systems such as **workstations** have dedicated resources but frequently use shared resources from **servers**

Handheld computers are resource poor, optimized for usability and battery life

Some computers have little or no user interface, such as embedded computers in devices and automobiles

OPERATING SYSTEM DEFINITION

- OS is a **resource allocator**
 - Manages all resources
 - Decides between conflicting requests for efficient and fair resource use
- OS is a **control program**
 - Controls execution of programs to prevent errors and improper use of the computer

“The one program running at all times on the computer” is the **kernel**.

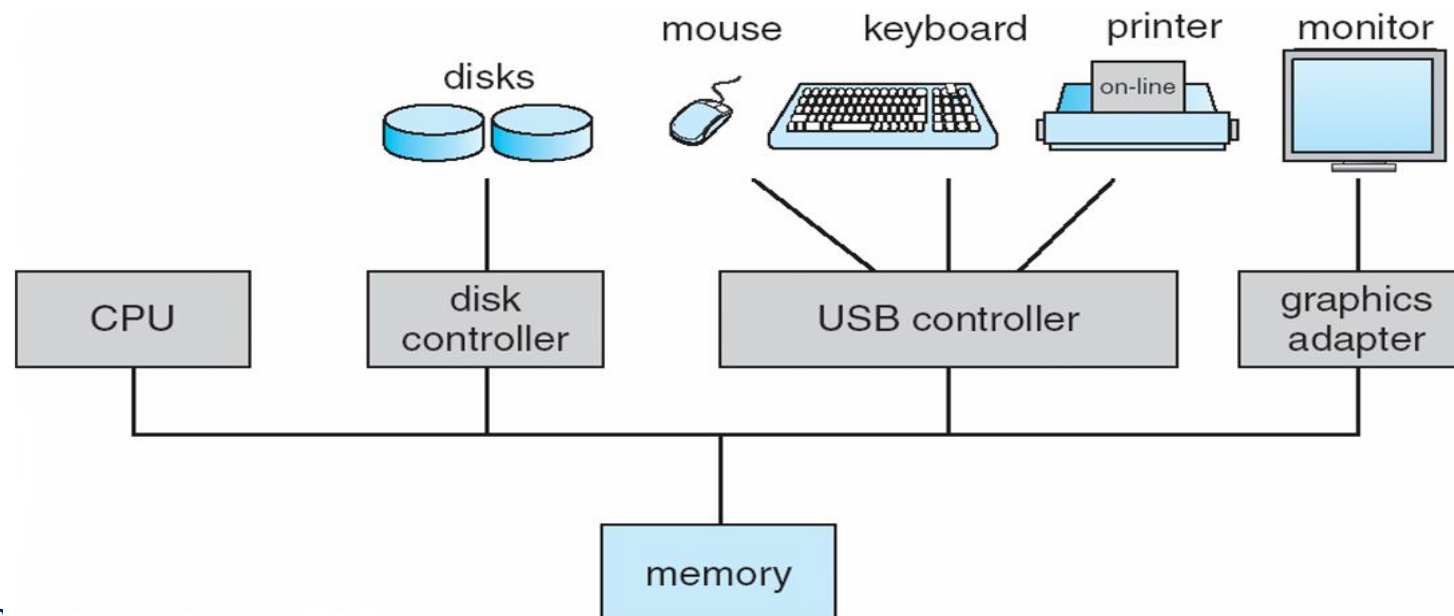


- **bootstrap program** is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, generally known as **firmware**
 - Initializes all aspects of system
 - Loads operating system kernel and starts execution



COMPUTER SYSTEM ORGANIZATION

- Computer-system operation
 - One or more CPUs, device controllers connect through common bus providing access to shared memory
 - Concurrent execution of CPUs and devices competing for memory cycles



COMPUTER-SYSTEM OPERATION

- I/O devices and the CPU can execute concurrently
- Each device controller is in charge of a particular device type
- Each device controller has a local buffer
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller
- Device controller informs CPU that it has finished its operation by causing an interrupt



COMMON FUNCTIONS OF INTERRUPTS

- Interrupt transfers control to the interrupt service routine generally, through the **interrupt vector**, which contains the addresses of all the service routines
- Interrupt architecture must save the address of the interrupted instruction
- A **trap** or **exception** is a software-generated interrupt caused either by an error or a user request
- An operating system is **interrupt driven**



INTERRUPT HANDLING

- The operating system preserves the state of the CPU by storing registers and the program counter
- Determines which type of interrupt has occurred:
 - **polling**
 - **vectored** interrupt system
- Separate segments of code determine what action should be taken for each type of interrupt

- After I/O starts, control returns to user program only upon I/O completion
 - Wait instruction idles the CPU until the next interrupt
 - Wait loop (contention for memory access)
 - At most one I/O request is outstanding at a time, no simultaneous I/O processing

- After I/O starts, control returns to user program without waiting for I/O completion
 - **System call** – request to the OS to allow user to wait for I/O completion
 - **Device-status table** contains entry for each I/O device indicating its type, address, and state
 - OS indexes into I/O device table to determine device status and to modify table entry to include interrupt

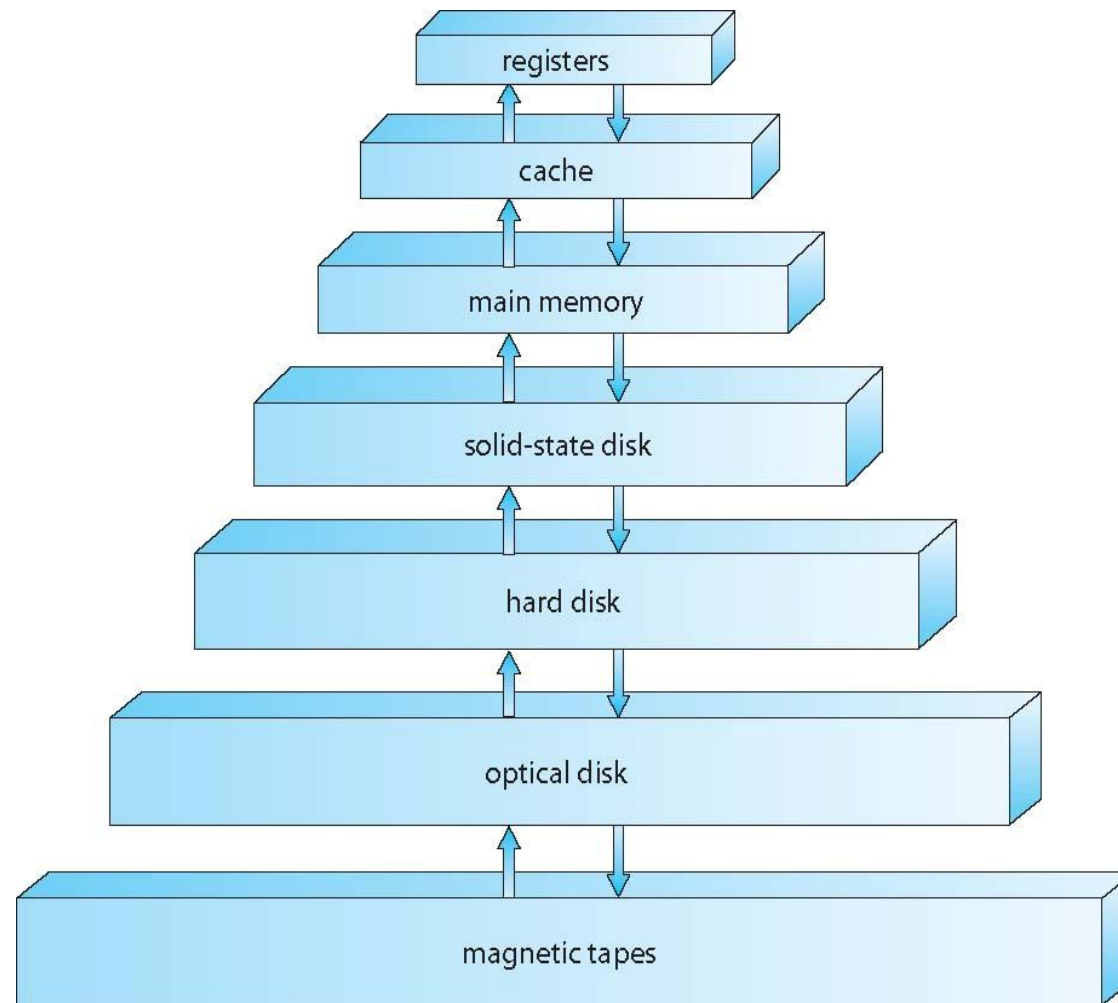


- Main memory – only large storage media that the CPU can access directly
 - **Random access**
 - Typically **volatile**
- Secondary storage – extension of main memory that provides large **nonvolatile** storage capacity
- Hard disks – rigid metal or glass platters covered with magnetic recording material
 - Disk surface is logically divided into **tracks**, which are subdivided into **sectors**
 - The **disk controller** determines the logical interaction between the device and the computer
- **Solid-state disks** – faster than hard disks, nonvolatile
 - Various technologies
 - Becoming more popular



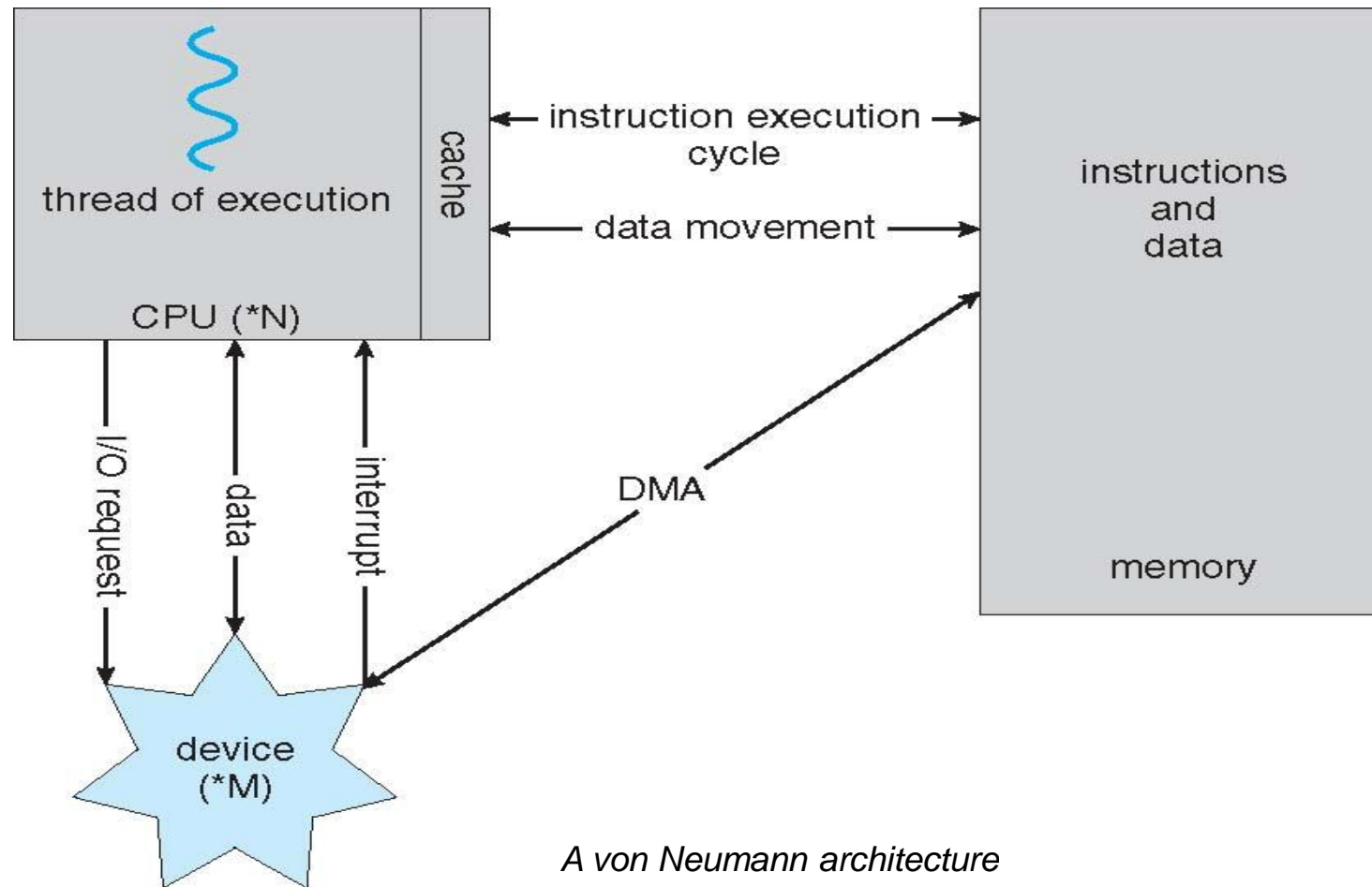
- Storage systems organized in hierarchy
 - Speed
 - Cost
 - Volatility
- **Caching** – copying information into faster storage system; main memory can be viewed as a cache for secondary storage
- **Device Driver** for each device controller to manage I/O
 - Provides uniform interface between controller and kernel

STORAGE-DEVICE HIERARCHY





HOW A MODERN COMPUTER WORKS



A von Neumann architecture