



# **SNS COLLEGE OF ENGINEERING**

**Kurumbapalayam(Po), Coimbatore – 641 107**

**Accredited by NAAC-UGC with 'A' Grade**

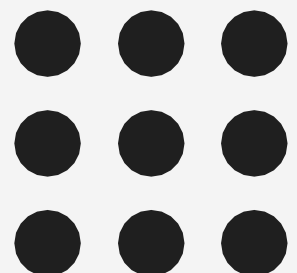
**Approved by AICTE, Recognized by UGC & Affiliated to Anna University, Chennai**

## **Department of Artificial Intelligence and Data Science**

**Course Name – Introduction to Artificial  
Intelligence**

**II Year / III Semester**

**Unit 2 Logic Programming**





# Logic programming

Predicate: It tells something about the subject. sentence involving the predicates that describe the property of objects are denoted by  $P(x)$  where  
P – The Predicate  
x – It is a variable denoting any object.

Characteristics of Predicate Logic:

1. Logical inferencing is allowed.
2. More accurate KR of facts of the real world.
3. Program designing is its application area.
4. Better theoretical foundation.
5. A predicate with no variable is called a Ground Atom.



## Constants

- Names of specific objects – E.g. doreen, gord, william, 32
- Functions
- Map objects to objects – E.g. father(doreen), age(gord), max(23,44)
- Variables
- For statements about unidentified objects or general statements – E.g. x,y,z,...

## Terms represent objects

- The set of terms is inductively defined by the following rules:
  - Constants: Any constant is a term
  - Variables: Any variable is a term
  - Functions: Any Expressions  $(t_1, \dots, t_n)$  of arguments (where each argument is a term and  $f$  is a function of arity  $n$ ) is a term
- Terms without variables are called ground terms
- Examples:
  - $c$
  - $f(c)$
  - $g(x,x)$
  - $g(f(c), g(x,x))$



## Rules of Predicate calculus symbols:

1. Set of letters (uppercase or lowercase) is allowed.
2. Set of digits (0-9) is allowed.
3. Underscore (  ) is allowed.
4. Blanks and non-alphanumeric characters can't be used.
5. Special characters like \$, \*, #, and / are not allowed.

## Symbols:

1. Predicate Symbols: It denotes relations or functional mapping from the elements of a domain to the values true or false.
2. Function Symbols: It denotes relations defined on a domain.
3. Variable Symbols: Lowercase unsubscribed or subscribed letters like x, y, z, t, u, v, etc. It can assume different values over a given domain.

4. Constants: It is fixed value term. They are individual symbols which are names of objects like 1, 2, 3, etc.

5. Quantifiers: There are two types of Quantifiers:

i. Existential Quantifier ( $\exists$ ): It means for some x or there is no x.

ii. Universal Quantifier ( $\forall$ ): It means for all x.

6. Logical Operators: There are five type of logical operators such as – not( $\sim$ ), and ( $\wedge$ ), or( $\vee$ ), implication, equivalence.

Two quantifiers: Universal ( $\forall$ ) and Existential ( $\exists$ )

- Allow us to express properties of collections of objects instead of enumerating objects by name
  - Apply to sentence containing variable
- Universal  $\forall$ : true for all substitutions for the variable
  - “for all”:  $\forall \langle \text{variables} \rangle \langle \text{sentence} \rangle$
- Existential  $\exists$ : true for at least one substitution for the variable
  - “there exists”:  $\exists \langle \text{variables} \rangle \langle \text{sentence} \rangle$
- Examples:
  - $\exists x: \text{Mother}(\text{art}) = x$
  - $\forall x \forall y: \text{Mother}(x) = \text{Mother}(y) \rightarrow \text{Sibling}(x,y)$
  - $\exists y \exists x: \text{Mother}(y) = x$

The set of formulas is inductively defined by the following rules:

1. Preciate symbols: If  $P$  is an  $n$ -ary predicate symbol and  $t_1, \dots, t_n$  are terms then  $P(t_1, \dots, t_n)$  is a formula.
2. Negation: If  $\phi$  is a formula, then  $\neg\phi$  is a formula
3. Binary connectives: If  $\phi$  and  $\psi$  are formulas, then  $(\phi \rightarrow \psi)$  is a formula. Same for other binary logical connectives.
4. Quantifiers: If  $\phi$  is a formula and  $x$  is a variable, then  $\forall x\phi$  and  $\exists x\phi$  are formulas.
  - Atomic formulas are formulas obtained only using the first rule

Any occurrence of a variable in a formula not in the scope of a quantifier is said to be a free occurrence

- Otherwise it is called a bound occurrence
- Thus, if  $x$  is a free variable in  $\phi$ , it is bound in  $\forall x\phi$  and  $\exists x\phi$
- A formula with no free variables is called a closed formula
- Example:  $x$  and  $y$  are bound variables,  $z$  is a free variable  
 $\forall x\forall y(P(f(x)) \rightarrow \neg(P(x)) \rightarrow Q(f(y), x, z))$

$S := \langle \text{Sentence} \rangle \langle \text{Sentence} \rangle :=$   
 $\langle \text{AtomicSentence} \rangle \mid \langle \text{Sentence} \rangle \langle \text{Connective} \rangle \langle \text{Sentence} \rangle \mid \langle \text{Quantifier} \rangle \langle \text{Variable} \rangle, \dots \langle \text{Sentence} \rangle$   
 $\mid \neg \langle \text{Sentence} \rangle \mid ( \langle \text{Sentence} \rangle )$   
 $\langle \text{AtomicSentence} \rangle := \langle \text{Predicate} \rangle ( \langle \text{Term} \rangle, \dots ) \langle \text{Term} \rangle := \langle \text{Function} \rangle ( \langle \text{Term} \rangle, \dots )$   
 $\mid \langle \text{Constant} \rangle$   
 $\mid \langle \text{Variable} \rangle \langle \text{Connective} \rangle := \wedge \mid \vee \mid \rightarrow \mid \leftrightarrow$   
 $\langle \text{Quantifier} \rangle := \mid$   
 $\langle \text{Constant} \rangle := "c" \mid "x1" \mid "john" \mid \dots$   
 $\langle \text{Variable} \rangle := "a" \mid "x" \mid "s" \mid \dots$   
 $\langle \text{Predicate} \rangle := "before" \mid "hasColor" \mid "raining" \mid \dots \langle \text{Function} \rangle := "mother" \mid "leftLegOf" \mid \dots$