

SNS COLLEGE OF ENGINEERING

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DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

19CS502- Automata Theory and Compiler Design

III YEAR / V SEMESTER

UNIT – 3 : Semantics and Context Sensitive Features

Topic : Syntax Directed Translation





- Associate information with the programming language constructs by attaching attributes to grammar symbols.
- Values of these attributes are evaluated by the semantic rules associated with the production rules.
- An attribute may hold almost any thing.
 - \succ a string, a number, a memory location, a complex record.

Grammar \neq **Semantic Rules** \equiv **SDT**









There are two ways to represent the semantic rules associated with grammar symbol

Syntax Directed Definitions (SDD) □ Syntax Directed Translation Schemes(SDTs)

Syntax-Directed Definitions:

- ✓ Give high-level specifications for translations
- \checkmark Hide many implementation details such as order of evaluation of semantic actions.
- \checkmark We associate a production rule with a set of semantic actions, and we do not say when they will be evaluated.

Syntax Directed Translation Schemes:

- \checkmark Indicate the order of evaluation of semantic actions associated with a production rule.
- \checkmark In other words, translation schemes give a little bit information about implementation details.







Eg : SDD vs SDT scheme

SDT Scheme		SDD
$E \rightarrow E + T$	{print '+'}	$E \rightarrow E + T$
$E \rightarrow E - T$	{print '-'}	$E \rightarrow E - T$
$E \rightarrow T$		$E \rightarrow T$
$T \rightarrow 0$	{print '0' }	$T \rightarrow 0$
$T \rightarrow 1$	{print '1' }	$T \rightarrow 1$
• • •		• • •
$T \rightarrow 9$	{print '9' }	$T \rightarrow 9$





T.code ='9'



Conceptually with both the syntax directed translation and translation scheme we

- > Parse the input token stream
- ➢ Build the parse tree
- \blacktriangleright Traverse the tree to evaluate the semantic rules at the parse tree nodes.

Input string \longrightarrow parse tree \longrightarrow dependency graph \longrightarrow evaluation order for semantic rules

Conceptual view of syntax directed translation







- A syntax-directed definition is a generalization of a context-free grammar in which:
 - \succ Each grammar symbol is associated with a set of attributes.
 - > This set of attributes for a grammar symbol is partitioned into two subsets called **Synthesized** and
 - **Inherited** attributes of that grammar symbol.

 \succ Each production rule is associated with a set of semantic rules.

- The value of an attribute at a parse tree node is defined by the semantic rule associated with a production at that node. The value of a **Synthesized attribute** at a node is computed from the values of attributes at the children in that node of the parse tree
- •
- The value of an **Inherited attribute** at a node is computed from the values of attributes at the siblings and parent of that node of the parse tree







Examples:

- Synthesized attribute : $E \rightarrow E1 + E2$ {E.val = E1.val + E2.val}
- Inherited attribute $:A \rightarrow XYZ$ $\{$ Y.val = 2 * A.val $\}$
- Semantic rules set up dependencies between attributes which can be represented by a dependency graph.
- This dependency graph determines the evaluation order of these semantic rules.
- Evaluation of a semantic rule defines the value of an attribute. But a semantic rule may also have some side effects such as printing a value.





Synthesized attribute



Inherited attribute



Syntax-Directed Definitions Example

Production $L \rightarrow E$ $E \rightarrow E1 + T$ $E \rightarrow T$ $T \rightarrow T1 * F$ $T \rightarrow F$ $F \rightarrow (E)$ $F \rightarrow digit$

Semantic Rules print(E.val) E.val = E1.val + T.valE.val = T.val

$$T.val = T1.val * F.val$$

$$T.val = F.val$$

$$F.val = E.val$$

Symbols E, T, and F are associated with a synthesized attribute val. The token digit has a synthesized attribute lexval (it is assumed that it is evaluated by the lexical analyzer).

Terminals are assumed to have synthesized attributes only. Values for attributes of terminals are usually supplied by the lexical analyzer.

The start symbol does not have any inherited attribute unless otherwise stated.







Annotated Parse Tree

Input: 5+3*4

Production $L \rightarrow E \$$ $E \rightarrow E1 + T$ $E \rightarrow T$ $T \rightarrow T1 * F$ $T \rightarrow F$ $F \rightarrow (E)$ $F \rightarrow digit$

Semantic Rules

print(E.val) E.val = E1.val + T.val E.val = T.val T.val = T1.val * F.val T.val = F.val F.val = E.val F.val = digit.lexval













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SDD – Inherited Attributes

Production	Semantic Rules
$D \rightarrow T L$	L.in = T.type
$T \rightarrow int$	T.type = integer
$T \rightarrow real$	T.type = real
$L \rightarrow L_1 id$	$L_1.in = L.in$, addtype(id .entry,L.in)
$L \rightarrow id$	addtype(id.entry,L.in)

- 1. Symbol T is associated with a Synthesized attribute *type*.
- 2. Symbol L is associated with an inherited attribute *in*.







Annotated parse tree







annotated parse tree



Dependency Graph







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SDD of Inherited attribute

Production	Semantic Rules	
$T \rightarrow F T'$	T'.inh = F.val T.val = T'.syn	F.val = 3
$T' \rightarrow *F T'_1$	$T'_{1}.inh = T'.inh \ge F.val$ $T'.syn = T'_{1}.syn$	
$T' \rightarrow \epsilon$	T'.syn = T'.inh	digit./exval = 3
$F \rightarrow \text{digit}$	F.val = digit.lexval	







Dependency Graph









Annotated Parse tree of (3+4)*(5+6)\$







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