

UNIT IV

MEMORY SYSTEM

Basic concepts of Semiconductor RAMs - ROMs – Speed, Size and Cost – Cache memories – Performance consideration – Virtual memory – **Memory**

Management requirements – Secondary storage - Case Study: Memory

Organization in Multiprocessors





Recap the previous Class



Memory management

- **Operating system** is concerned with transferring programs and data between secondary storage and main memory.
- Operating system **needs memory routines** in addition to the other routines.
- Operating system routines are assembled into a virtual address space called **system space**.
- System space is separate from the space in which user application programs reside. (**This is user space.**)
- Virtual address space is divided into **one system space + several user spaces**.



Memory management (contd..)

- MMU uses the **contents of the page table base register** to determine the address of the page table to be used in the translation.
 - Changing the contents of the page table base register can enable us to use a different page table, and switch from one space to another.
- At any given time, the page table base register can point to one page table.
 - Thus, only one page table can be used in the translation process at a given time.
 - Pages belonging to only one space are accessible at any given time.



Memory management (contd..)

- When multiple, independent user programs coexist in the main memory, how to ensure that one program does not modify/destroy the contents of the other?
- Processor usually has two states of operation:
 - Supervisor state & User state.
- **Supervisor state:**
 - Operating system routines are executed.



Memory management (contd..)

■ User state:

- User programs are executed.
- Certain privileged instructions **cannot be executed** in user state.
- These privileged instructions include the ones which change page table base register.
- Prevents one user from accessing the space of other users.

Magnetic Hard Disks

Magnetic disks constitute a traditional method for **non-volatile storage** of information using **magnetic technology**.

- Broadly three types of devices appeared:
 - a) **Floppy disk** : made of bendable plastic
 - b) **Magnetic drum** : made of solid metal
 - c) **Hard disk** : made of metal or glass
- All of these rely on a rotating platter (metal or glass or plastic) coated with **a thin magnetic material**, and **use a moveable read/write head** to read and write data from / to the disk.
 - Data stored as tiny magnets.

Magnetic Hard Disks



Magnetic drum
(62.5 Kbytes)



8" floppy disk
(360 Kbytes)



3.5" floppy disk
(1.2 Mbytes)



3.5" magnetic disk
(1 Tbytes)



1.8" solid-state disk
(512 Gbytes)

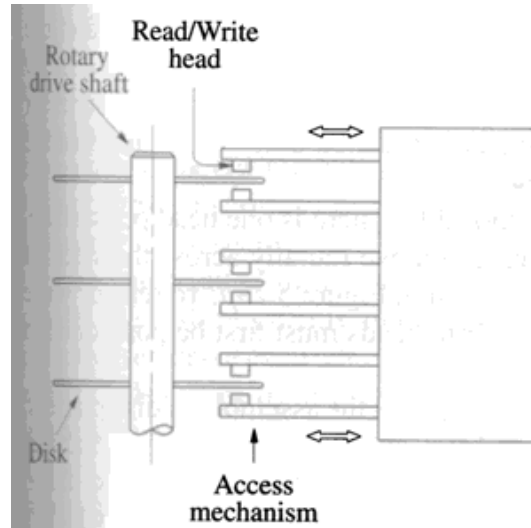
Magnetic Hard Disks

Since the platters in a hard disk are made of rigid metal or glass, they provide **several advantages** over floppy disks:

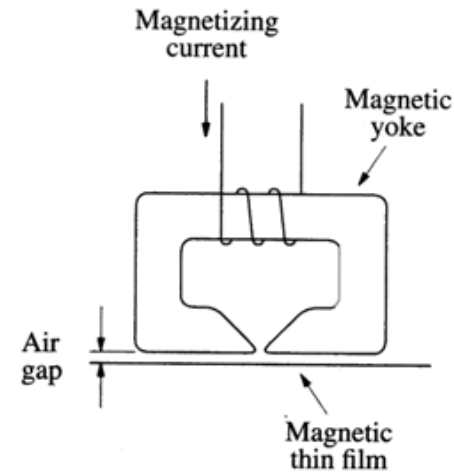
- They can be **larger**.
- Can have **higher density** since they can be controlled more precisely.
- Has a **higher data rate** because it spins faster
- **No physical contact** with read/write head as it spins faster.
 - The read/write head floats on a cushion of air (few microns separation).
 - Requires dustless environment.
 - Results in higher reliability
- More than one platters can be incorporated in the same unit.



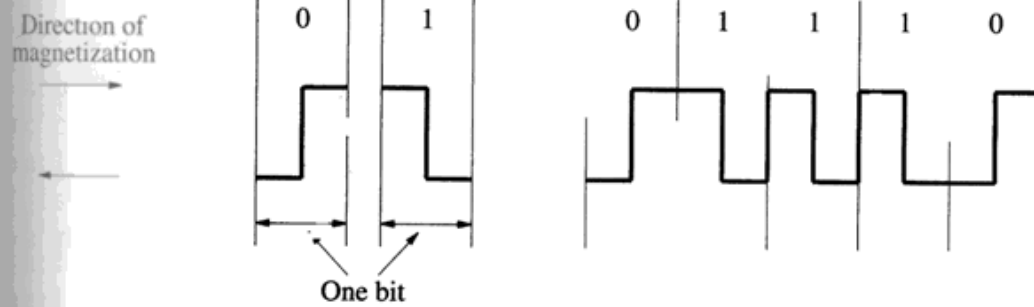
Magnetic Hard Disks



(a) Mechanical structure



(b) Read/Write head detail



(c) Bit representation by phase encoding

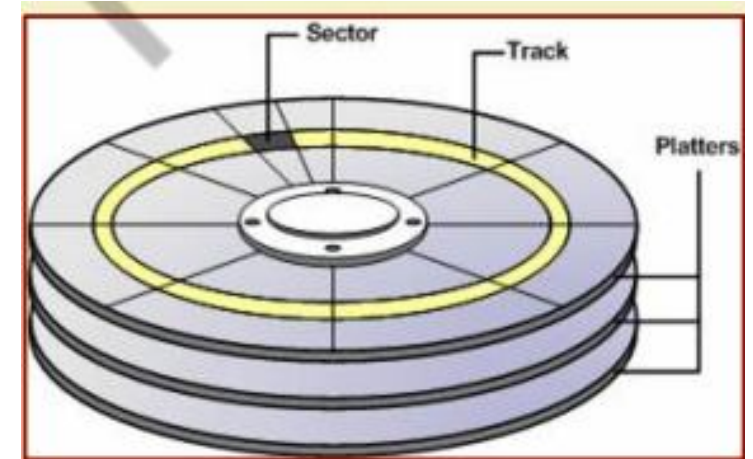
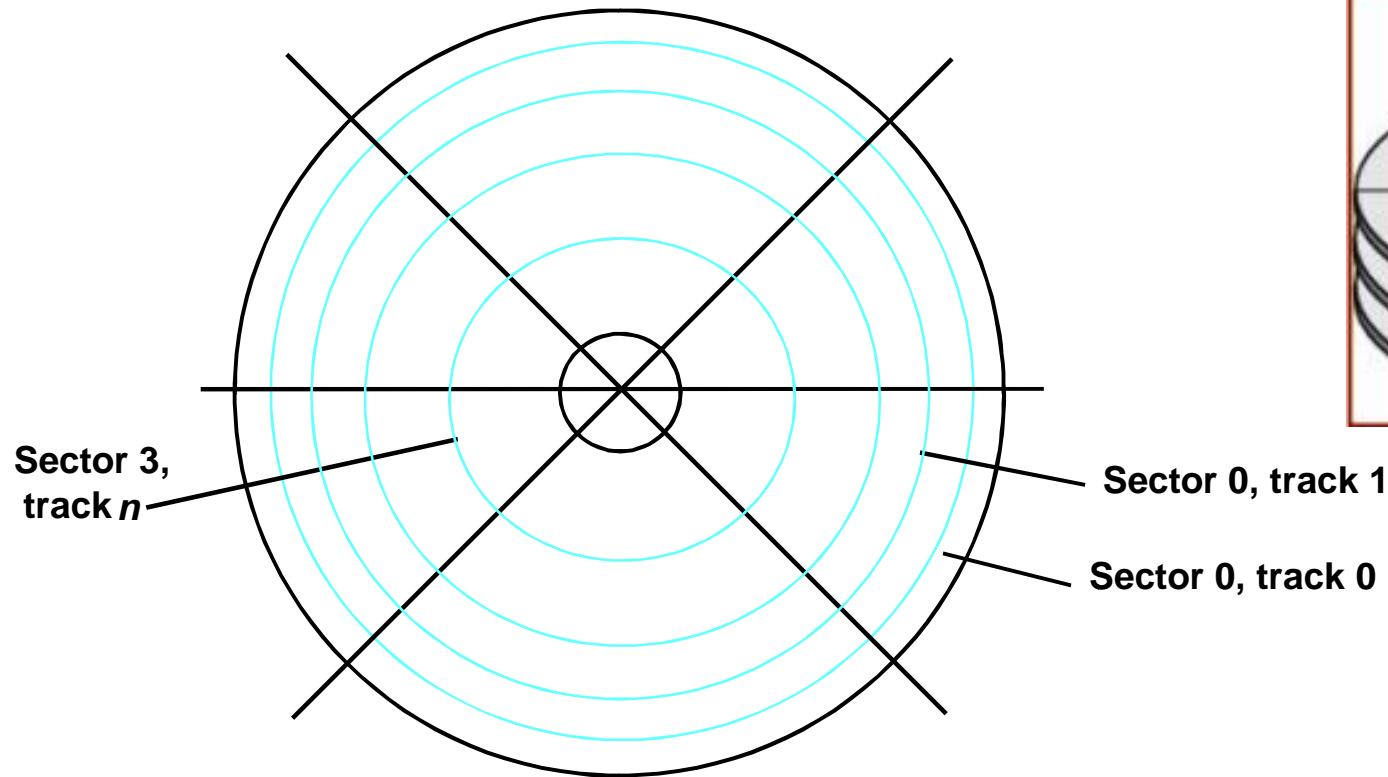
Organization of Data on a Disk

The hard disks consists of a **collection of platters** (typically, 1 to 5), which are connected together and can spin in unison.

- Each platter **has two recording surfaces**, and comes in various sizes (1 – 8 inches).
- The stack of platter typically rotates at a **speed of 5400 to 7200 rpm**.
- Each disk surface is divided into concentric circles called **tracks**.
 - The number of tracks per surface **can vary from 1000 to 5000**.
- Each track is divided into a **number of sectors (64 – 200 sectors/track)**.
 - Typical sector size: **512 – 2048 bytes**.
 - Sector is the **smallest unit** that can be read or written.
- The disk heads for all the surfaces are connected and move together.
- All the tracks under the heads at a given time on all surfaces is called a **cylinder**.



Organization of Data on a Disk



Disk Access Time

There are **three components** to the access time in hard disk:

a) Seek time:

- The time required to **move the head to the desired track**.
- Average seek times are in the range **8 – 20 msec**.
- Actual average can be **25 – 30% less than** this number, since accesses to disks are often localized.

Disk Access Time

b) Rotational delay:

- Once the head is on the correct track, we must wait for the desired sector to rotate under the head.
- The average delay or latency is the **time for half the rotation**.

Examples:

For 3600 rpm, average rotational delay = $0.5 \text{ rotation} / 3600 \text{ rpm} = \mathbf{8.30 \text{ msec}}$

For 5400 rpm, average rotational delay = $0.5 \text{ rotation} / 5400 \text{ rpm} = \mathbf{5.53 \text{ msec}}$

For 7200 rpm, average rotational delay = $0.5 \text{ rotation} / 7200 \text{ rpm} = \mathbf{4.15 \text{ msec}}$



Disk Access Time

c) Transfer time:

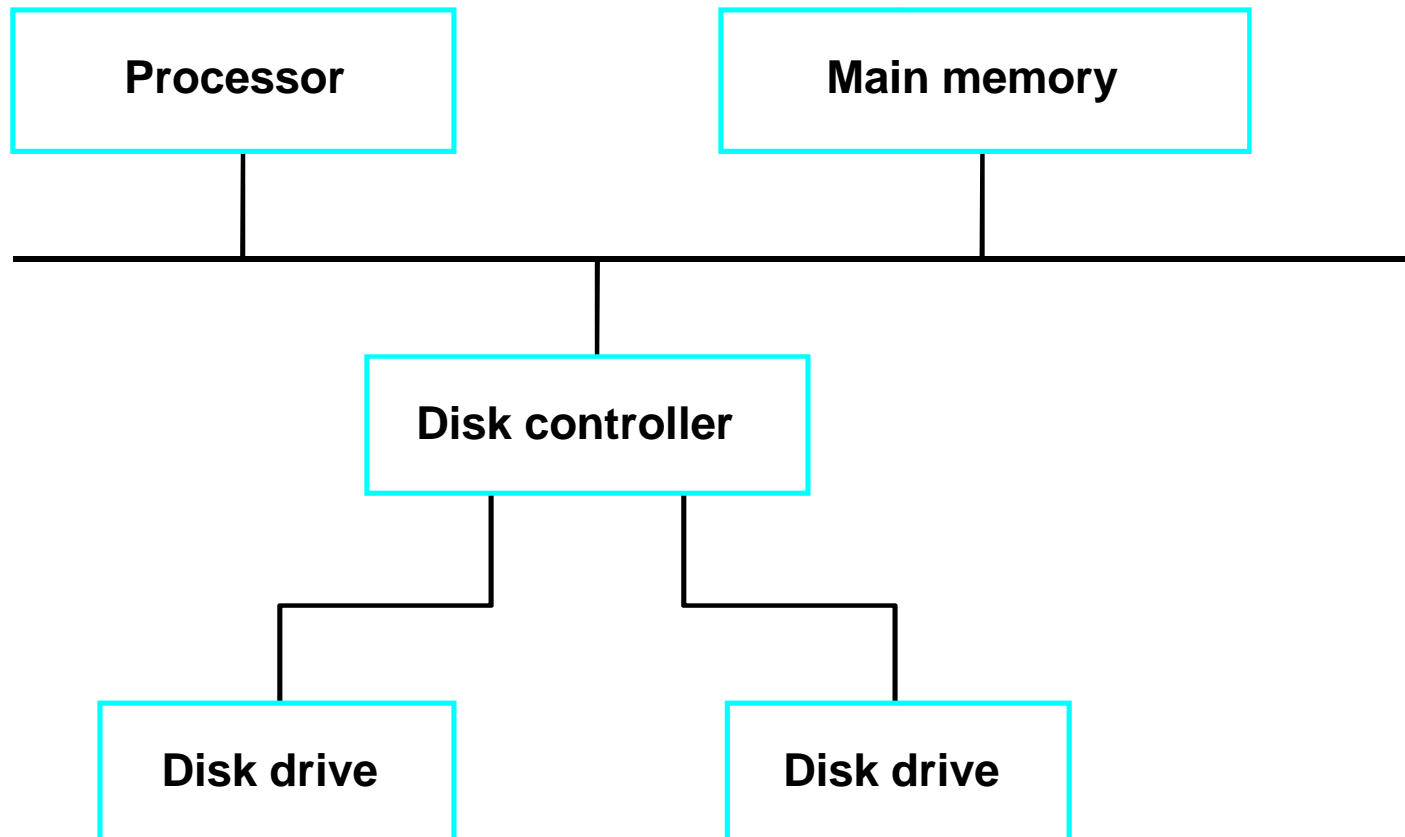
- The total time to transfer a block of data (typically, a sector).
- Transfer rates are typically **15 MB/sec** or more.
- Transfer time **depends on:**
 - Sector size
 - Rotation speed of the disk
 - Recording density on the tracks

Example 1

- Consider a disk with sector size 512 bytes, 2000 tracks per surface, 64 sectors per track, three double-sided platters, and average seek time of 10 msec.
 - a) What is the capacity of the disk?
 - Bytes/track = $512 \times 64 = 32\text{K}$
 - Bytes/surface = $32\text{K} \times 2000 = 64,000\text{K}$
 - Bytes/disk = $64,000\text{K} \times 3 \times 2 = 384,000\text{K}$
 - b) If the disk platters rotate at 7200 rpm, and one track of data can be transferred per revolution, what is the transfer rate?
 - Transfer rate = Capacity of a track / average rotational delay
 - $= 32\text{K} / 4.15 \text{ msec} = 7,711 \text{ Kbytes/sec}$



Disk Controller





Disk Controller

- Seek
- Read
- Write
- Error checking

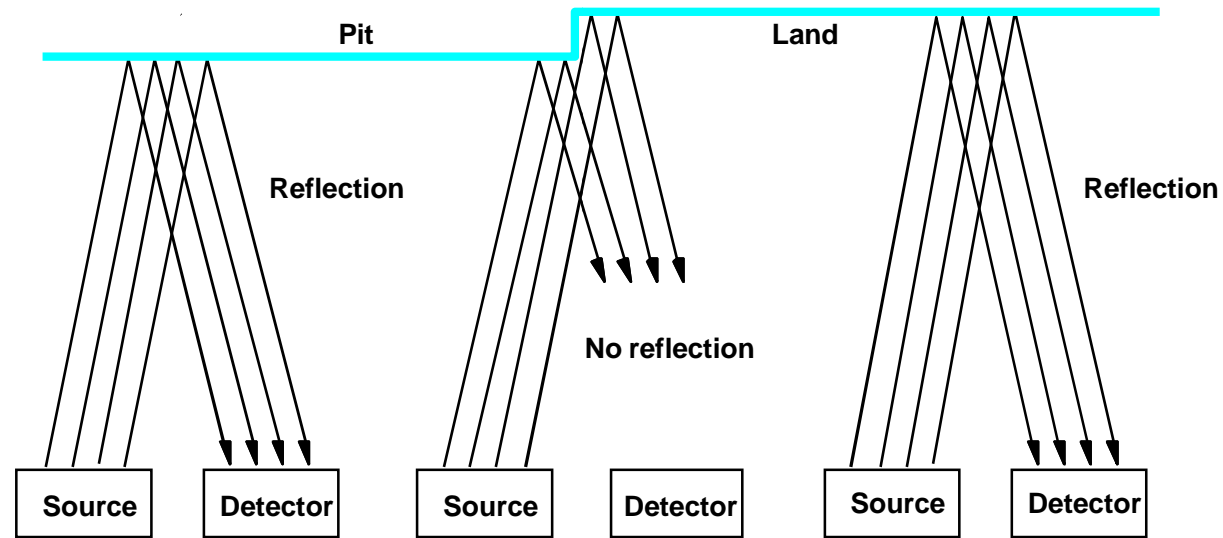
RAID Disk Arrays

- Redundant Array of Inexpensive Disks
- Using multiple disks makes **it cheaper for huge storage**, and also possible to **improve the reliability** of the overall system.
- RAID0 – data striping
- RAID1 – identical copies of data on two disks
- RAID2, 3, 4 – increased reliability
- RAID5 – parity-based error-recovery

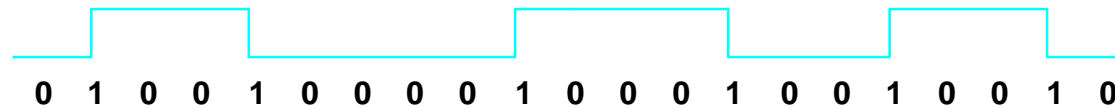


Optical Disks

(a) Cross-section



(b) Transition from pit to land



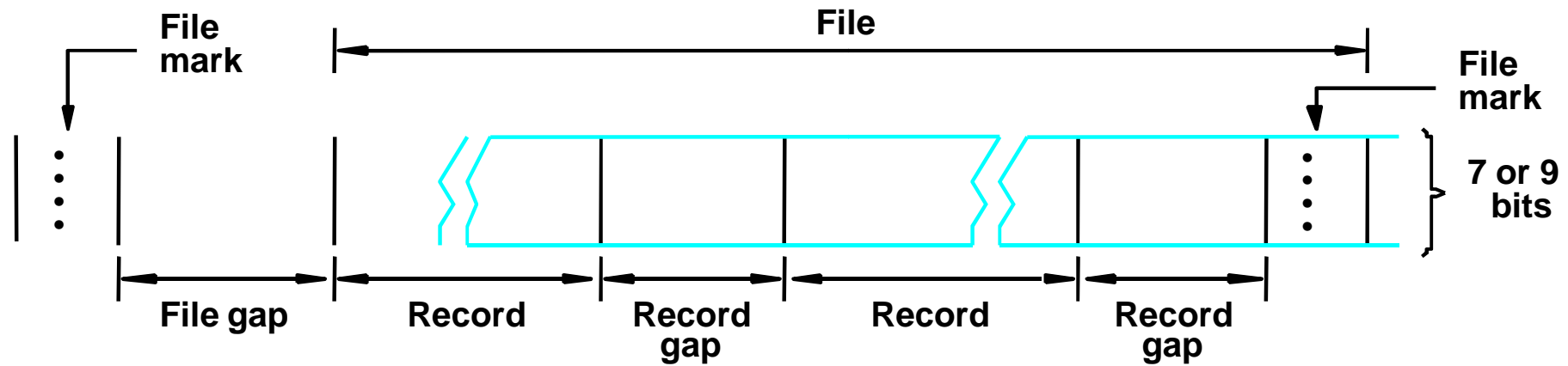
(c) Stored binary pattern

Optical Disks

- CD-ROM
- CD-Recordable (CD-R)
- CD-ReWritable (CD-RW)
- DVD
- DVD-RAM



Magnetic Tape Systems



Organization of data on magnetic tape.



TEXT BOOK

Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", McGraw-Hill, 6th Edition 2012.

REFERENCES

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THANK YOU