



TOPIC:11.- Proof methods and strategy.

## The Theory of Inference for Predicate calculus

### Universal Specification (US)

If a statement of the form  $(\forall x)[A(x)]$  is assumed to be true, then the universal quantifier can be dropped to obtain  $A(y)$  is true for any arbitrary object 'y' in the universe.

### Existential Specification (ES)

From  $\exists x(A(x))$  one can conclude  $A(y)$ , provided that y is not free in any given premise and also not free in any prior step of the derivation.

### Universal Generalization (UG)

From  $A(x)$  one can conclude

$$A(y) \Rightarrow (\forall x)(A(x))$$

### Existential Generalization (EG)

$$A(y) \Rightarrow (\exists x)(A(x))$$



1. Express the statement "Every student in this class has completed Assignment-I" as a quantifiers.

Let  $C(x)$  :  $x$  is in this class

$A(x)$  :  $x$  has completed Assignment-I

For all  $x$ , if  $x$  is in this class, then  $x$  has completed Assignment-I.

∴ Its symbolic form  $(\forall x) (C(x) \rightarrow A(x))$

2.

Write each of the following in symbolic form.

All men are good (b) No men are good

Some men are good (e) Some men are not good

Let  $M(x)$  :  $x$  is a man

$G(x)$  :  $x$  is good

(a) For all  $x$ ,  $x$  is a man, then  $x$  is good.

∴  $(\forall x) (M(x) \rightarrow G(x))$



(b) For all  $x$ , if  $x$  is a man, then  $x$  is not good.  
 $\therefore (\forall x) (M(x) \rightarrow \neg G(x))$

(c) There exists an  $x$ ,  $x$  is a man and  $x$  is good.  
 $\therefore (\exists x) (M(x) \wedge G(x))$

(d) There exists an  $x$ ,  $x$  is a man and  $x$  is not good.  
 $(\exists x) (M(x) \wedge \neg G(x))$

3.

Show that  $(\forall x)(P(x) \rightarrow Q(x)) \wedge (\forall x)(Q(x) \rightarrow R(x))$   
 $\Rightarrow (\forall x)(P(x) \rightarrow R(x))$ .

$\{1\}$	1) $(\forall x)(P(x) \rightarrow Q(x))$	Rule P
$\{1\}$	2) $P(y) \rightarrow Q(y)$	Rule US
$\{3\}$	3) $(\forall x)(Q(x) \rightarrow R(x))$	Rule P
$\{3\}$	4) $Q(y) \rightarrow R(y)$	Rule US
$\{1, 3\}$	5) $P(y) \rightarrow R(y)$	Rule T ( $P \rightarrow Q, Q \rightarrow R$ $\Rightarrow P \rightarrow R$ )
$\{1, 3\}$	6) $(\forall x)(P(x) \rightarrow R(x))$	Rule UG



4)

Show that  $(\forall x)(P(x) \vee Q(x)) \Rightarrow (\forall x) P(x) \vee (\exists x) Q(x)$   
We use indirect method, by assuming

$\neg [(\forall x) P(x) \vee (\exists x)(Q(x))]$  as an additional premise

{1}	1) $\neg [(\forall x) P(x) \vee (\exists x) Q(x)]$	Rule P
{1}	2) $(\exists x) \neg P(x) \wedge (\forall x) \neg Q(x)$	Rule T (DeMorgan's)
{1}	3) $(\exists x) \neg P(x)$	Rule T ( $P \wedge Q \Rightarrow P$ )
{1}	4) $(\forall x) \neg Q(x)$	Rule T ( $P \wedge Q \Rightarrow Q$ )
{1}	5) $\neg P(y)$	Rule ES
{1}	6) $\neg Q(y)$	Rule US
{1}	7) $\neg P(y) \wedge \neg Q(y)$	Rule T ( $P, Q \Rightarrow P \wedge Q$ )
{1}	8) $\neg (P(y) \vee Q(y))$	Rule T (DeMorgan's)
{9}	9) $(\forall x)(P(x) \vee Q(x))$	Rule P





{9}	10) $P(y) \vee Q(y)$	Rule U5
{1,9}	11) $[P(y) \vee Q(y)]^{\wedge}$ $\neg [P(y) \vee Q(y)]$	Rule T ( $P, Q \Rightarrow P \wedge Q$ )
{1,9}	12) F	

5) Verify the validity of the following argument. Every living thing is a plant or animal. John's gold fish is alive and it is not a plant. All animals have hearts.

Therefore John's gold fish has a heart.

Let  $L(x)$  :  $x$  is a living thing  
 $P(x)$  :  $x$  is a plant  
 $A(x)$  :  $x$  is an animal  
 $H(x)$  :  $x$  has a heart

Then, the given premises are

$$(1) (\forall x) [L(x) \rightarrow (P(x) \vee A(x))]$$

$$(2) L(j) \wedge \neg P(j)$$

$$(3) (\forall x) [A(x) \rightarrow H(x)]$$

Conclusion is  $H(j)$



$\{1\}$	1) $(\forall x) [L(x) \rightarrow (P(x) \vee A(x))]$	Rule P
$\{1\}$	2) $L(j) \rightarrow P(j) \vee A(j)$	Rule US
$\{3\}$	3) $L(j) \wedge \neg P(j)$	Rule P
$\{3\}$	4) $L(j), \neg P(j)$	Rule T ( $P \wedge a \Rightarrow P, a$ )
$\{1, 3\}$	5) $P(j) \vee A(j)$	Rule T ( $P, P \rightarrow a \Rightarrow a$ )
$\{1, 3\}$	6) $\neg P(j) \rightarrow A(j)$	Rule T ( $P \rightarrow a \Rightarrow \neg P \vee a$ )
$\{7\}$	7) $(\forall x) (A(x) \rightarrow H(x))$	Rule P
$\{7\}$	8) $A(j) \rightarrow H(j)$	Rule US
$\{1, 3, 7\}$	9) $\neg P(j) \rightarrow A(j)$	Rule T ( $P \rightarrow Q, a \rightarrow R \Rightarrow P \rightarrow R$ )
$\{1, 3, 7\}$	10) $H(j)$	Rule T ( $P, P \rightarrow a \Rightarrow a$ )