

# UNIT II

## ARITHMETIC OPERATIONS

Addition and subtraction of signed numbers – **Design of fast adders** –  
Multiplication of positive numbers - Signed operand multiplication- fast  
multiplication – Integer division – Floating point numbers and operations





# Recap the previous Class



# Binary Adders

**Note:**

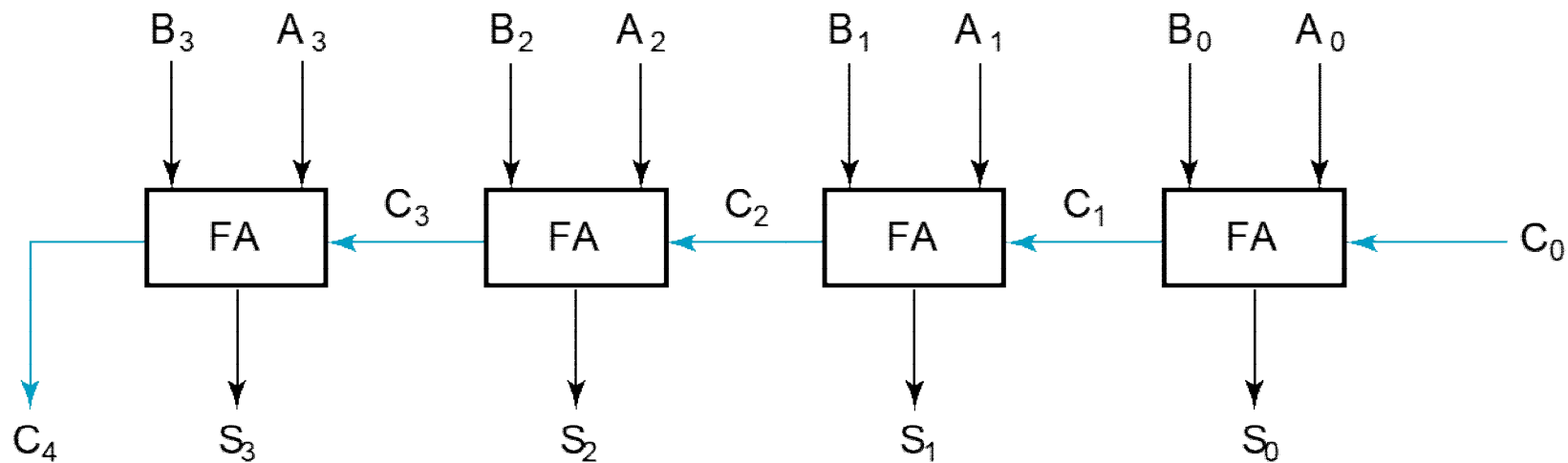
carry out of cell  $i$  becomes carry in of cell  $i + 1$

Description	Subscript 3 2 1 0	Name
Carry In	0 1 1 0	$C_i$
Augend	1 0 1 1	$A_i$
Addend	<u>0 0 1 1</u>	$B_i$
Sum	1 1 1 0	$S_i$
Carry out	0 0 1 1	$C_{i+1}$



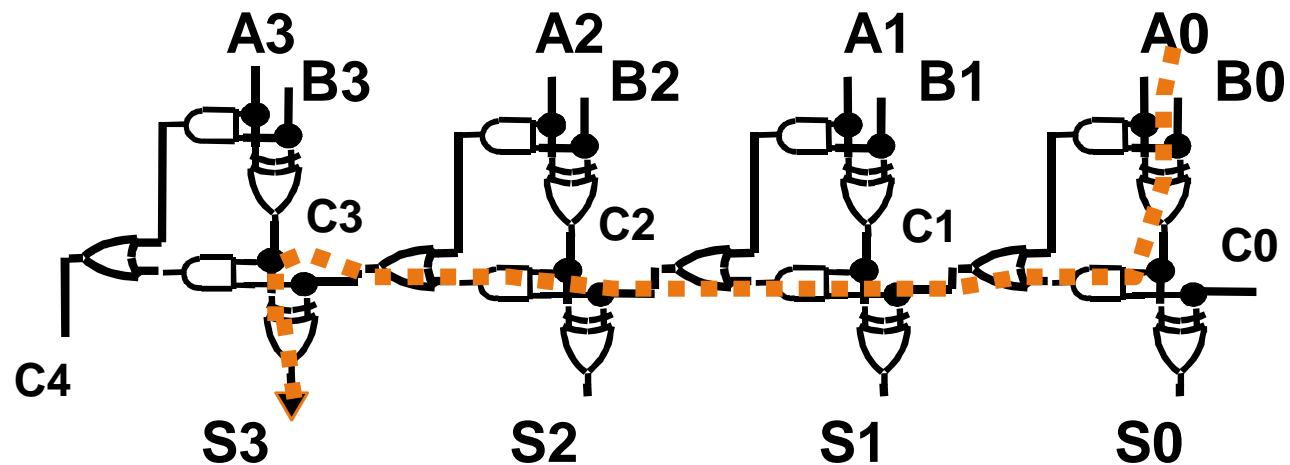
# 4 bit Ripple carry Adder

- A four-bit Ripple Carry Adder made from four 1-bit Full Adders



# Carry Propagation & Delay

- One problem with the addition of binary numbers is **the length of time to propagate the ripple carry from the least significant bit to the most significant bit.**
- The gate-level propagation path for a 4-bit ripple carry adder of the last example:







# Carry Lookahead Adder

$$S_i = x_i \oplus y_i \oplus c_i$$

$$C_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

Factorizing

$$C_{i+1} = x_i y_i + (x_i + y_i) c_i$$

We can write

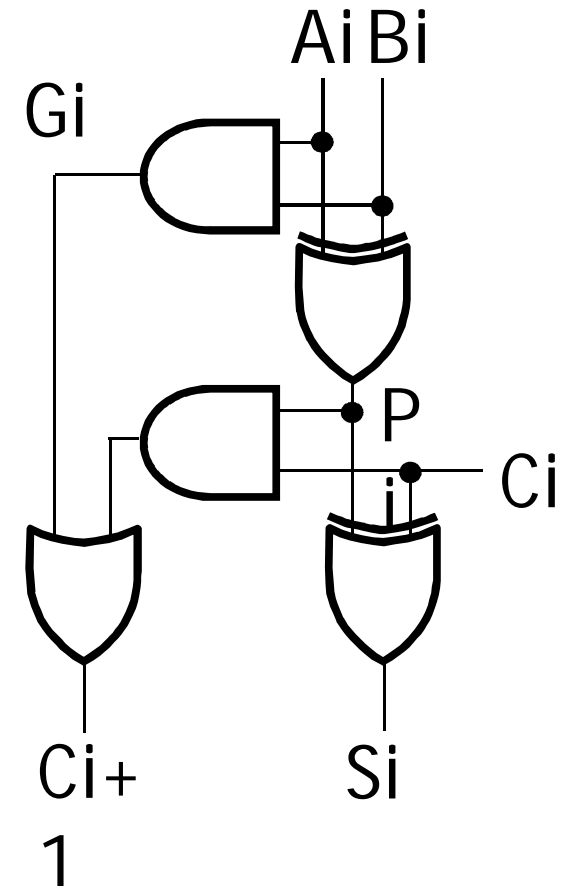
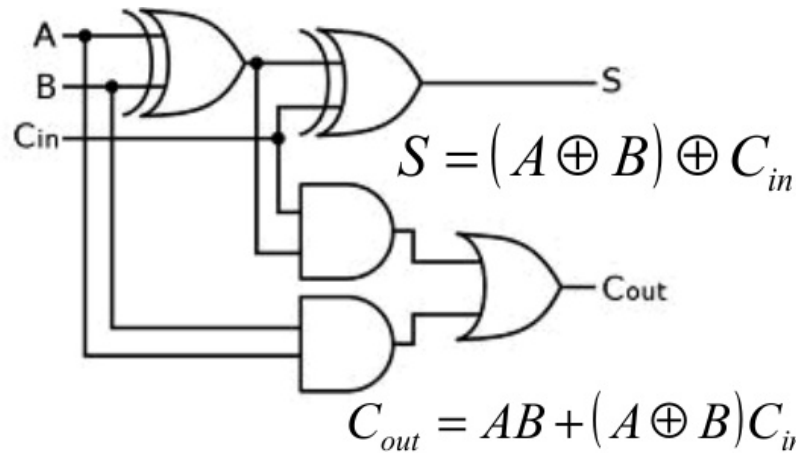
$$C_{i+1} = G_i + P_i c_i$$

where

$$G_i = x_i y_i$$

$$P_i = x_i + y_i$$

These two signal conditions are called *generate*, denoted as  $G_i$ , and *propagate*, denoted as  $P_i$  respectively



# Carry Lookahead Adder

- In the ripple carry adder:
  - $G_i$ ,  $P_i$ , and  $S_i$  are local to each cell of the adder
  - $C_i$  is also local each cell
- In the carry lookahead adder, in order to reduce the length of the carry chain,  $C_i$  is changed to a more global function spanning multiple cells
- Defining the equations for the Full Adder in term of the  $P_i$  and  $G_i$ :

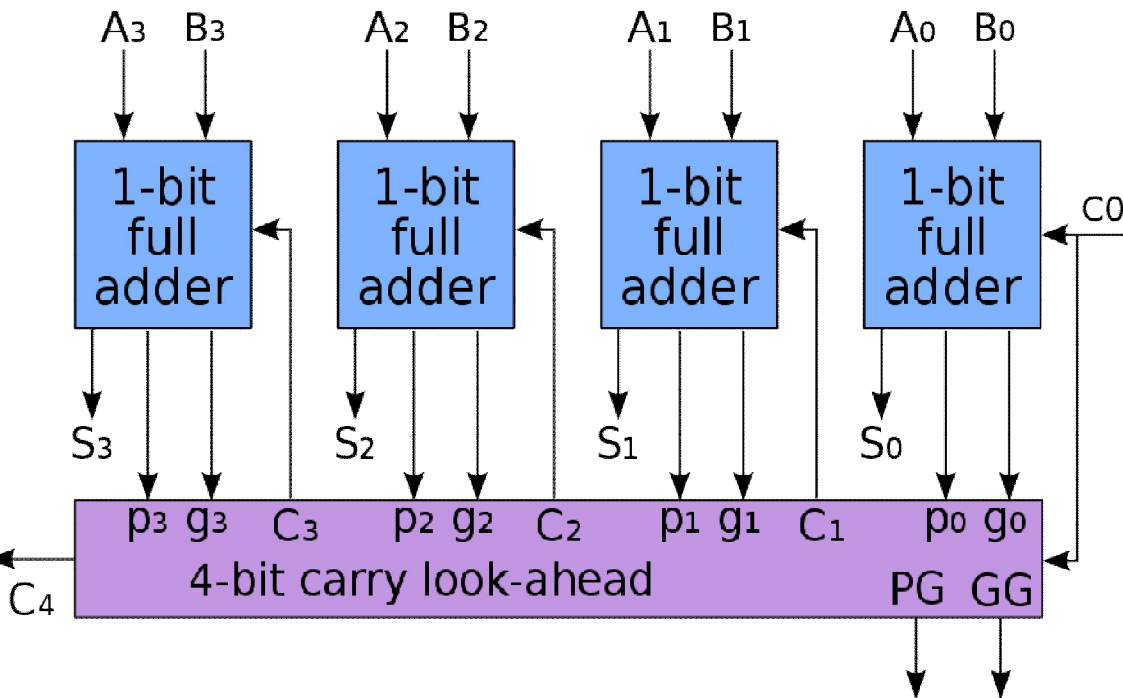
$$P_i = A_i \oplus B_i$$

$$G_i = A_i B_i$$

$$S_i = P_i \oplus C_i$$

$$C_{i+1} = G_i + P_i C_i$$

# Carry Lookahead Adder



$$C_1 = G_0 + P_0 C_0$$

$$C_2 = G_1 + P_1 C_1 = G_1 + P_1(G_0 + P_0 C_0) = G_1 + P_1 G_0 + P_1 P_0 C_0$$

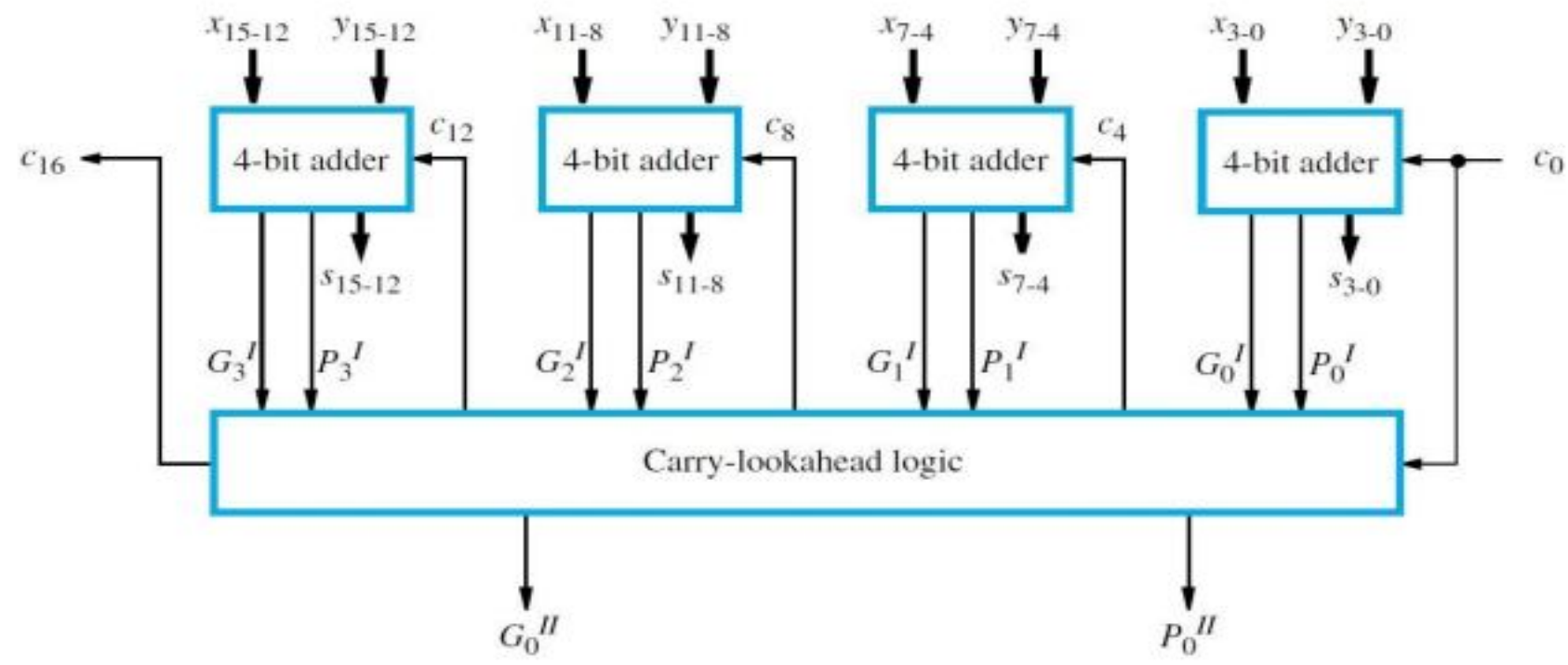
$$C_3 = G_2 + P_2 C_2 = G_2 + P_2(G_1 + P_1 G_0 + P_1 P_0 C_0) = G_2 + P_2 G_1 + P_2 P_1 G_0 + P_2 P_1 P_0 C_0$$

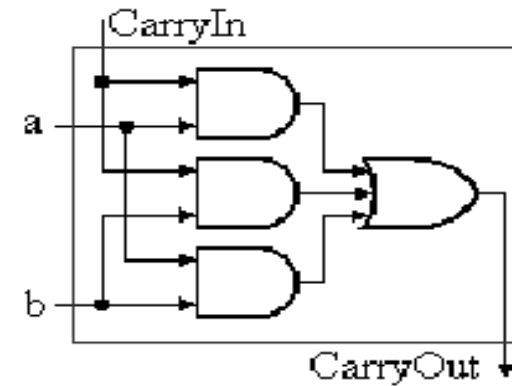
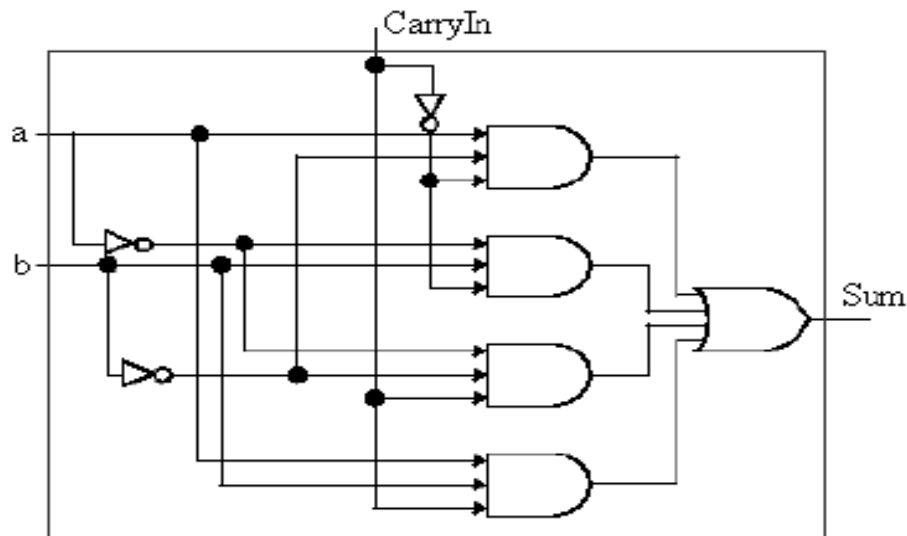
$$C_4 = G_3 + P_3 C_3 = G_3 + P_3 G_2 + P_3 P_2 G_1 + P_3 P_2 P_1 G_0 + P_3 P_2 P_1 P_0 C_0$$

$$C_{i+1} = G_i + P_i G_{i-1} + P_i P_{i-1} C_{i-1}$$



# 16 bit Carry Lookahead Adder





$$\text{Carryout} = (b \cdot \text{CarryIn}) + (a \cdot \text{CarryIn}) + (a \cdot b)$$

$$\text{Sum} = (a \cdot b' \cdot \text{CarryIn}') + (a' \cdot b \cdot \text{CarryIn}') + (a' \cdot b' \cdot \text{CarryIn}) + (a \cdot b \cdot \text{CarryIn})$$

## TEXT BOOK

Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", McGraw-Hill, 6th Edition 2012.

## REFERENCES

1. David A. Patterson and John L. Hennessey, "Computer organization and design", MorganKauffman ,Elsevier, 5th edition, 2014.
2. William Stallings, "Computer Organization and Architecture designing for Performance", Pearson Education 8th Edition, 2010
3. John P.Hayes, "Computer Architecture and Organization", McGraw Hill, 3rd Edition, 2002
4. M. Morris R. Mano "Computer System Architecture" 3rd Edition 2007
5. David A. Patterson "Computer Architecture: A Quantitative Approach", Morgan Kaufmann; 5th edition 2011

# THANK YOU