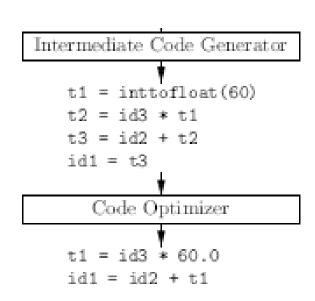




# Unit V Intermediate Code Optimization

- Program transformation technique
- Improves code consume less resources
- Transforms the code to make it more efficient
- Output is not changed
- Intermediate code ②
   optimization ② code
   generation is made easier





# **Code Optimization**



### Optimization

- Machine Independent Optimization
- takes in the intermediate code and transforms a part of the code that does not involve any CPU registers
- Example:

```
do
{
   item = 10;
   value = value + item;
} while(value<100);</pre>
```

This code involves repeated assignment of the identifier item, which if we put this way:

```
Item = 10;
do
{
    value = value + item;
} while(value<100);</pre>
```

- Machine Dependent optimization
  - Target code
    - Rearrangement of machine instructions to improve the efficiency of the code
    - Divide the code into basic blocks





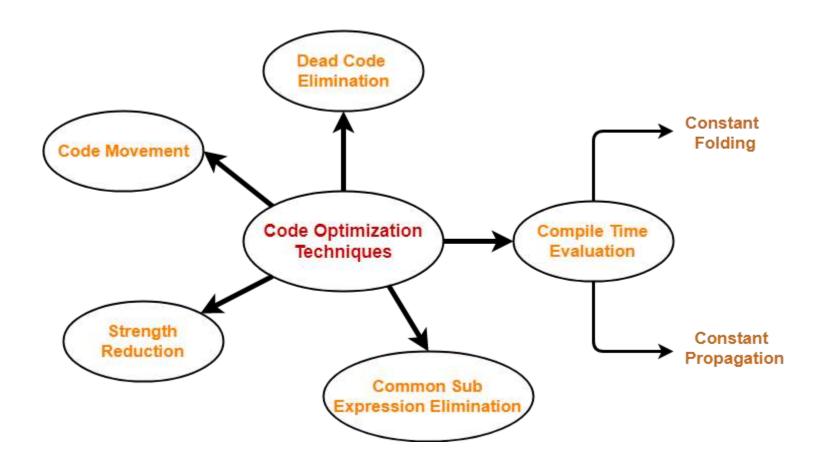
# Peephole Optimization

- Optimization eliminates the redundant instruction from a small area of code
- Set of code peephole / window
- Goals:
  - Improves performance
  - Reduce memory footprint
  - Reduce code size





# Principle sources of optimization





# **Compile Time Evaluation**



### Constant Folding

- Folding the constants
- The expressions that contain the operands having constant values at compile time are evaluated.
- Example:
- return (3+5); ② return 8;
- Cir=(22/7)\*diameter ② cir = 3.14\*diameter

### Constant Propagation

 If some variable has been assigned some constant value, then it replaces that variable with its constant value in the further program during compilation.

### - Example:

- radius =10,pi=3.14
- area=pi\*radius\*radius; 2 area=3.14\*10\*10;





# **Common Sub Expression**

Code before Optimization	Code after Optimization
$S1 = 4 \times i$	$S1 = 4 \times i$
$S2 = a[S1]$ $S3 = 4 \times j$	S2 = a[S1]
S4 = 4 x i // Red und ant Expression	S3 = 4 x j
S5 = n	S5 = n
S6 = b[S4] + S5	S6 = b[S1] + S5





# **Code Movement**

## Code before Optimization

# for (int j = 0; j < n; j ++)

{

X = Y + Z

 $\mathbf{a}[\mathbf{j}] = 6 \times \mathbf{j}$ 

## Code after Optimization

$$x = y + z$$

for (int 
$$j = 0$$
;  $j < n$ ;  $j ++$ )

{

$$a[j] = 6 \times j$$
;

}





# **Dead Code Elimination**

• Eliminates the dead code

### Code before Optimization

### Code after Optimization

```
i = 0;
if (i == 1)
{
a = x + 5;
}
```





# Strength Reduction

- Reduces the strength of expressions
- Replaces expensive operators with cheaper one
- Example
  - B = A \* 2 ? B = A + A
  - Cost of multiplication is higher than the addition