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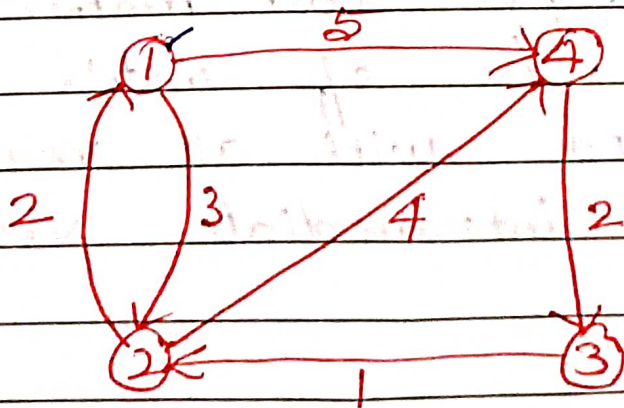
All-pairs shortest path algorithm - Floyd's Marshall algorithm.

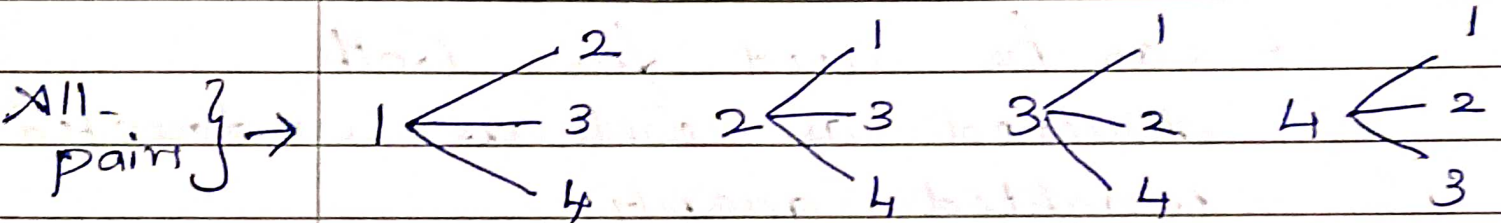
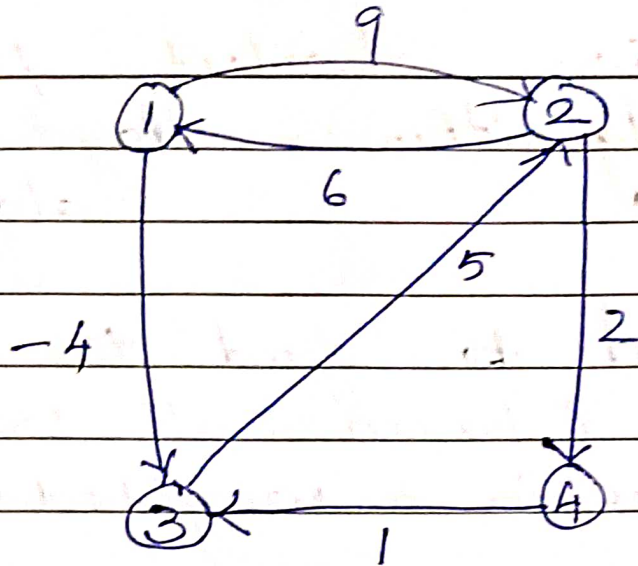
- used to find the shortest path between all pair of vertices in a weighted graph.

- can be used for both directed as well as undirected weighted graph.

- it does not work for graphs with negative cycles (where the sum of edges is negative).

Consider the following graph





The alg. generate a matrix which represent the minimum distance from any node to all other nodes in the graph.

Initially the o/p matrix is same as the give cost matrix of graph. Later, the o/p matrix is updated with all vertices k as the intermediate vertex.

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Algorithm (Procedure)

Begin

for $k = 0$ to n , do

for $i = 0$ to n , do

for $j = 0$ to n , do

if $\text{cost}[i, k] + \text{cost}[k, j]$

$< \text{cost}[i, j]$, then

$\text{cost}[i, j] = \text{cost}[i, k] +$

$\text{cost}[k, j]$

done for

done for

done for

display the current cost
matrix.

End.

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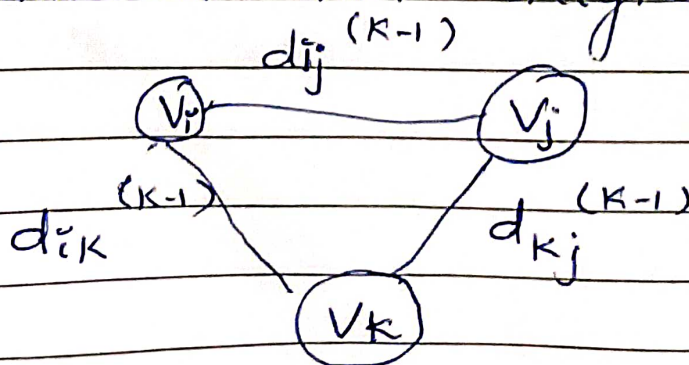
- Floyd's algorithm computes the distance matrix of a weighted graph with 'n' vertices through a series of $n \times n$ matrices

$$D^{(0)} \dots D^{(k-1)}, D^{(k)} \dots D^{(n)}$$

- Compute all the elements of each matrix $D^{(k)}$ from its immediate predecessor $D^{(k-1)}$ in series.

- Let $d_{ij}^{(k)}$ be the element in the i^{th} row and j^{th} column of matrix $D^{(k)}$.

- d_{ij} is equal to the length of the shortest path among all paths from the i^{th} vertex V_i to all the j^{th} vertex V_j with their intermediate vertices numbered not higher than K_i .



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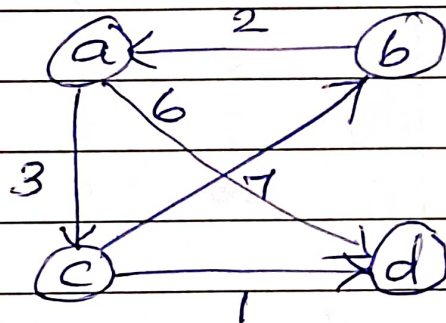
The formula for calculating $d_{ij}^{(k)}$

$$d_{ij}^{(k)} = \min \left\{ d_{ij}^{(k-1)}, d_{ik}^{(k-1)} + d_{kj}^{(k-1)} \right\}$$

for $k=1$

$$d_{ij}^{(0)} = w_{ij}$$

Consider the following example.


$$D^{(0)} =$$

| | a | b | c | d |
|---|----------|----------|----------|----------|
| a | 0 | ∞ | 3 | ∞ |
| b | 2 | 0 | ∞ | ∞ |
| c | ∞ | 7 | 0 | 1 |
| d | 6 | ∞ | ∞ | 0 |

| | | | | | |
|-------------|----------|----------|-----|----------|---|
| $D^{(1)} =$ | | a | b | c | d |
| a | 0 | ∞ | 3 | ∞ | |
| b | 2 | 0 | (5) | ∞ | |
| c | ∞ | 7 | 0 | 1 | |
| d | 6 | ∞ | (9) | 0 | |

a - known node
 b - c through a
 d - c through a

| | | | | | |
|-------------|-----|----------|---|----------|---|
| $D^{(2)} =$ | | a | b | c | d |
| a | 0 | ∞ | 3 | ∞ | |
| b | 2 | 0 | 5 | ∞ | |
| c | (9) | 7 | 0 | 1 | |
| d | 6 | ∞ | 9 | 0 | |

a, b are known node
 c - a through b

| | | | | | |
|-------------|---|------|---|-----|---|
| $D^{(3)} =$ | | a | b | c | d |
| a | 0 | (10) | 3 | (4) | |
| b | 2 | 0 | 5 | (6) | |
| c | 9 | 7 | 0 | 1 | |
| d | 6 | (16) | 9 | 0 | |

a, b, c are known node

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$$D^{(4)} = \begin{array}{c|cccc} & a & b & c & d \\ \hline a & 0 & 10 & 3 & 4 \\ b & 2 & 0 & 5 & 6 \\ c & \textcircled{7} & 7 & 0 & 1 \\ d & 6 & 16 & 9 & 0 \end{array}$$

a, b, c, d are known nodes