

# **SNS COLLEGE OF TECHNOLOGY An Autonomous Institution Coimbatore-35**

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# **DEPARTMENT OF ELECTRONICS & COMMUNICATION ENGINEERING 19ECB301-ANALOG AND DIGITAL COMMUNICATION**

III YEAR / V SEMESTER

# **UNIT 5 – INFORMATION THEORY AND ERROR CONTROL CODING**

**TOPIC** – MUTUAL INFORMATION-CHANNEL CAPACITY

MUTUAL INFORMATION/19ECB301 – ANALOG AND DIGITAL COMMUNICATION/S.KAVIPRIYA/ECE/SNSCT

8/11/2023





## **MUTUAL INFORMATION**

- Mutual information measures the amount of information that can be obtained about ► one random variable by observing another.
- It is important in communication where it can be used to maximize the amount of information shared between sent and received signals.
- ▶ The mutual information denoted by I (X, Y) of a channel is defined by:  $I(X; Y) = H(X) - H\left(\frac{X}{V}\right) b/symbol$
- Since H (X) represents the uncertainty about the channel input before the channel output is observed and H (X/Y) represents the uncertainty about the channel input after the channel output is observed, the mutual information I (X; Y) represents the uncertainty about the channel input that is resolved by observing the channel output.







# **MUTUAL INFORMATION**

- Properties of Mutual Information I (X; Y) I(X; Y) = I(Y; X) $I(X;Y) \ge 0$ ▶ I(X; Y) = H(Y) - H(Y/X)I(X; Y) = H(X) + H(Y) - H(X, Y)
- The Entropy corresponding to mutual information [i.e. I (X, Y)] indicates a measure of the information transmitted through a channel. Hence, it is called 'Transferred information'.







# **CHANNEL CAPACITY**

- The channel capacity represents the maximum amount of information that can be transmitted by a channel per second.
- To achieve this rate of transmission, the information has to be processed properly or coded in the most efficient manner.
- Channel Capacity per Symbol Cs: The channel capacity per symbol of a discrete memoryless channel (DMC) is defined as  $C_S = \max_{\{P(x,y)\}} I(X;Y)b/symbol$

Where the maximization is over all possible input probability distributions  $\{P(x_i)\}$  on X.

Channel Capacity per Second C: I f 'r' symbols are being transmitted per second. then the maximum rate od transmission of information per second is 'r Cs', this is the channel capacity per second and is denoted by C (b/s) i.e.  $C = rC_S b/s$ 

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# **CAPACITIES OF SPECIAL CHANNELS**

- **Lossless Channel:** For a lossless channel, H(X/Y) = 0 and I(X; Y) = H(X).
- Thus the mutual information is equal to the input entropy and no source information is lost in transmission.

$$C_S = \max_{\{P(x_i)\}} H(X) = \log_2 m$$

Where m is the number of symbols in X.

- **Deterministic Channel:** For a deterministic channel, H(Y|X) = 0 for all input distributions  $P(x_i)$  and I(X; Y) = H(Y).
- Thus the information transfer is equal to the output entropy. The channel capacity per symbol will be

 $C_S = \max_{\{P(x_i)\}} H(Y) = \log_2 n$ 

where n is the number of symbols in Y.





# **CAPACITIES OF SPECIAL CHANNELS**

**Noiseless Channel:** since a noiseless channel is both lossless and deterministic, we have I (X; Y) = H (X) = H (Y) and the channel capacity per symbol is

 $C_S = \log_2 m = \log_2 n$ 

**Binary Symmetric Channel:** For the BSC, the mutual information is  $I(X;Y) = H(Y) + p \log_2 p + (1-p) \log_2 (1-p)$ 

And the channel capacity per symbol will be  $C_S = 1 + p \log_2 p + (1 - p) \log_2 p$ 

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$$p_2(1-p)$$



# **THANK YOU**

8/11/2023

