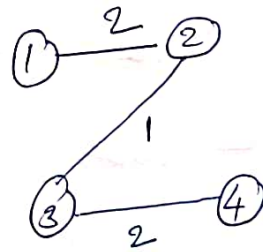
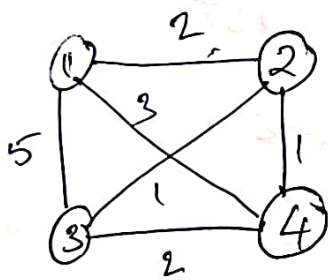


Minimum Spanning Tree

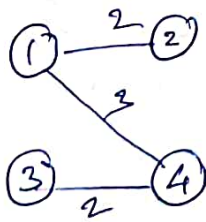
→ A spanning tree of a connected graph is a connected acyclic graph that contains all the vertices of a subgraph

- All vertices
- Connected
- No cycles

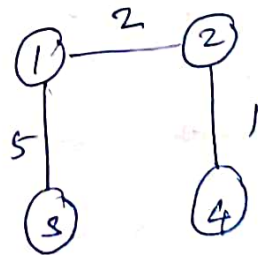
→ Remove edges if it forms cycle



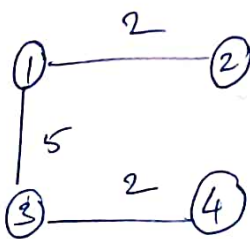
Cost = 5



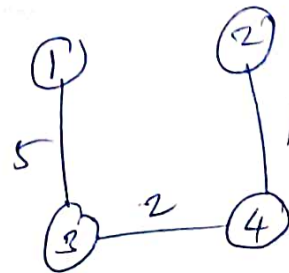
Cost = 7



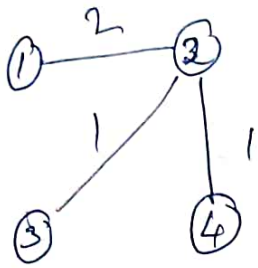
Cost = 8



Cost = 9



Cost = 8



Cost = 4

2 types algorithm to find MST

① Prim's Algorithm

② Kruskal Algorithm

Prim's Algorithm:

→ one way compute MST.

→ uses Greedy Techniques

→ begins with set 'u' initialized to 'v'

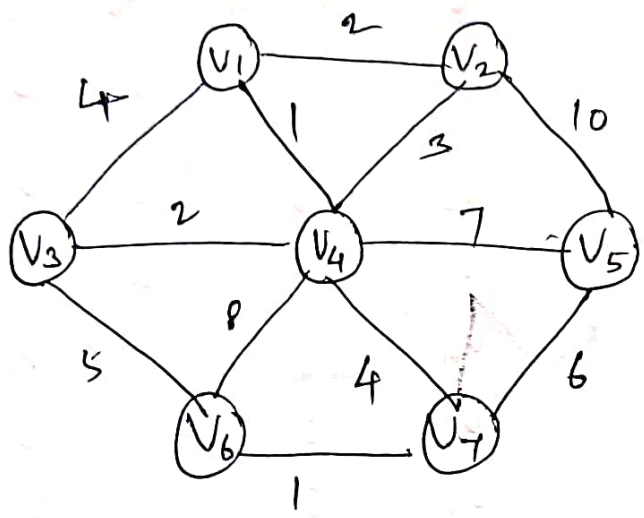
→ choose one edge at a time

→ each state it finds the shortest edge (u, v) will not in MST.

↙ ↘ Cost is low among all edges.
will be in MST

→ $d_w = \text{Min}(d_w, c_{w,v})$

eg:



- ① Choose an arbitrary start vertex
- ② keep including connected edges

Initialize Config

V	known	dv	Pv
V ₁	0	0	0
V ₂	0	∞	0
V ₃	0	∞	0
V ₄	0	∞	0
V ₅	0	∞	0
V ₆	0	∞	0
V ₇	0	∞	0

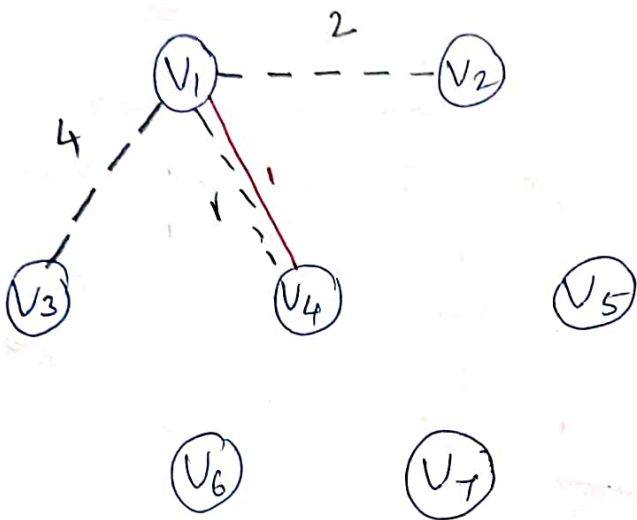
* Consider V₁ as source vertex, and find adjacent vertices of V₁

$$V_1 \rightarrow V_2, V_3, V_4$$

$$V_1 \xrightarrow{2} V_2$$

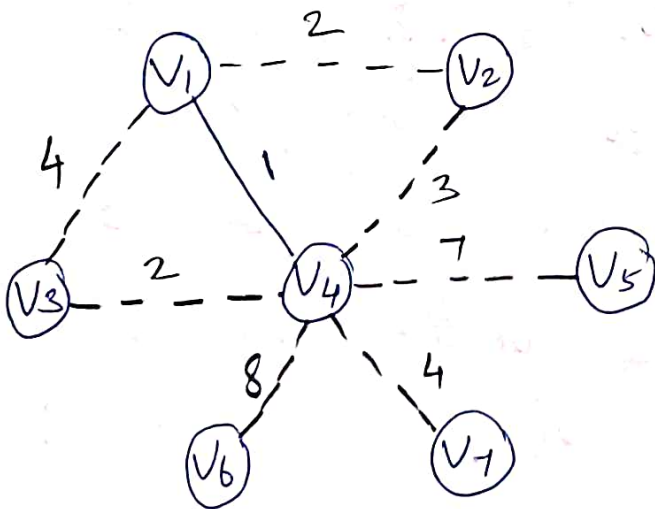
$$V_1 \xrightarrow{4} V_3$$

$$V_1 \xrightarrow{1} V_4 \quad \text{- minimum distance so select edges}$$



V	known	dv	P _v
V ₁	1	0	0
V ₂	0	2	V ₁
V ₃	0	4	V ₁
V ₄	0	1	V ₁
V ₅	0	∞	0
V ₆	0	∞	0
V ₇	0	∞	0

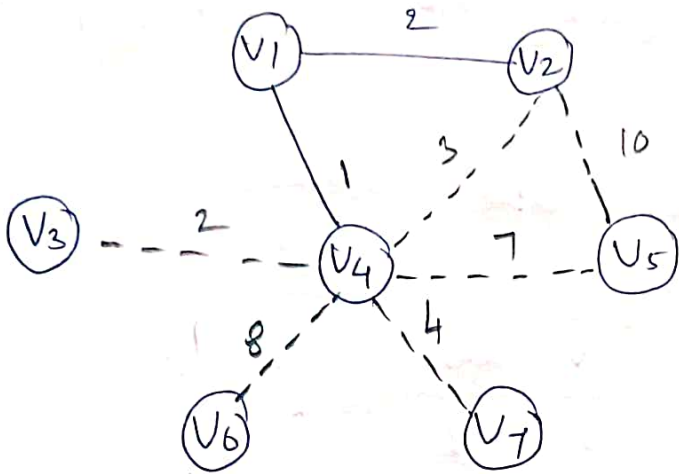
→ In this table minimum dv is 1 i.e V₄ it is unknown vertex, choose V₄ is known vertex find adjacent vertices V₄ → V₂, V₃, V₅, V₆, V₇



V	known	dv	P _v
V ₁	1	0	0
V ₂	0	2	V ₁
V ₃	0	2	V ₄
V ₄	1	1	V ₁
V ₅	0	7	V ₄
V ₆	0	8	V ₄
V ₇	0	4	V ₄

→ Next minimum dv is V₂, V₃ [unknown vertex], choose V₂ as known vertex

$V_2 \rightarrow V_1, V_4, V_5$ [V_1, V_4 already known vertex].



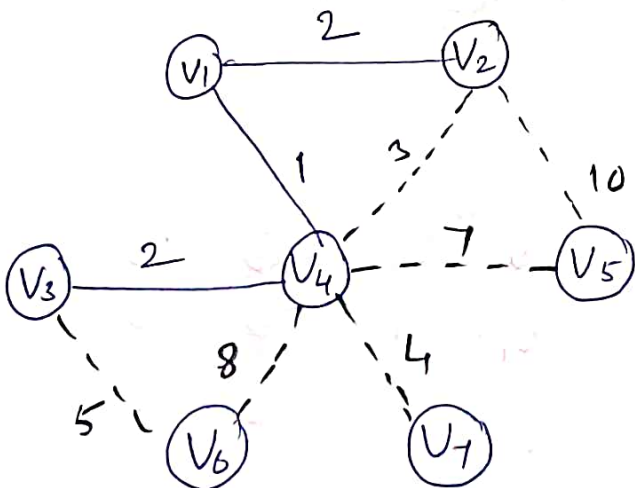
V	known	d_V	P_V
V_1	1	0	0
V_2	1	2	V_1
V_3	0	2	V_4
V_4	1	1	V_1
V_5	0	7	V_4
V_6	0	8	V_4
V_7	0	4	V_4

V_3 as known

~~V_3~~ adjacent of $V_3 \rightarrow$ ^{already known vertex} V_1, V_4, V_6

$$V_4 - V_6 = 8$$

$V_3 \rightarrow V_6 = 5 \rightarrow$ Minimum distance, update table entry



V	known	d_V	P_V
V_1	1	0	0
V_2	1	2	V_1
V_3	1	2	V_4
V_4	1	1	V_1
V_5	0	7	V_4
V_6	0	5	V_4
V_7	0	4	V_4

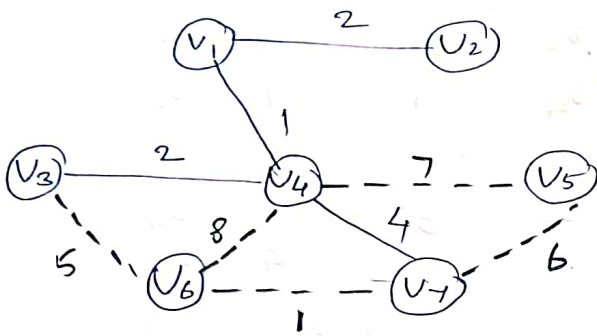
V_7 having minimum distance, so consider V_7 as known vertex

V_7 adjacent vertex $u = V_6, V_4, V_5$
 \rightarrow Already known.

$V_4 - V_5$ is 7 }
 $V_5 - V_7$ is 6 } \rightarrow Min update table

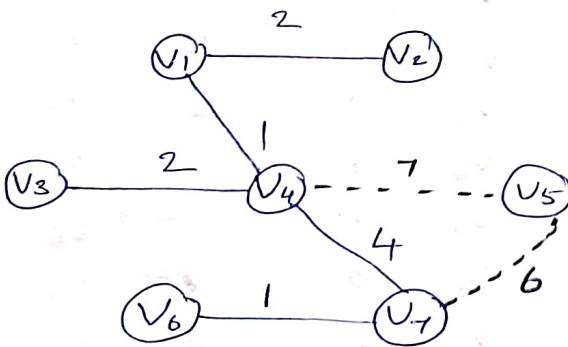
$V_6 - V_3 = 5$

$V_6 - V_7 = 1 \rightarrow$ Min update table



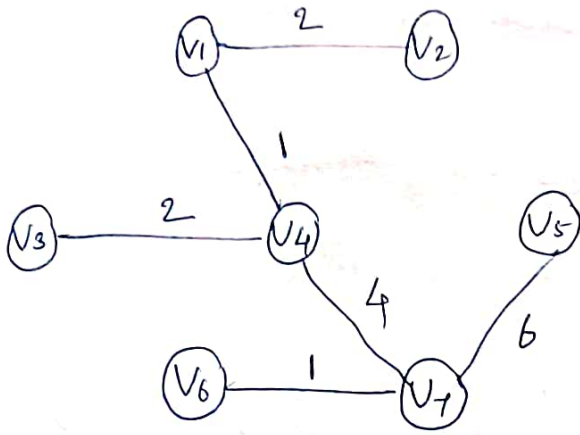
v	known	dv	Pv
V1	1	0	0
V2	1	2	V1
V3	1	2	V4
V4	1	1	V1
V5	0	5	V7
V6	0	1	V7
V7	1	4	V4

$\rightarrow V_6$ as known vertex.
 V_6 adjacent - V_3, V_4, V_7



v	known	dv	Pv
V1	1	0	0
V2	1	2	V1
V3	1	2	V4
V4	1	1	V1
V5	0	6	V7
V6	1	1	V7
V7	1	4	V4

V_5 as known vertex



V	known	dv	Pv
V ₁	1	0	0
V ₂	1	2	V ₁
V ₃	1	2	V ₄
V ₄	1	1	V ₁
V ₅	1	6	V ₇
V ₆	1	1	V ₇
V ₇	1	4	V ₄

Routine:

Void Prims (Table T)

{ int i;

Vertex v, w;

for (i = 0; i < numvertices; i++)

{

 T[i].known = false;

 T[i].dist = Infinity;

 T[i].path = 0;

}

for (i; i)

{ let v be the start vertex
with smallest distance

 T[v].dist = 0;

 T[v].known = True;

 for each w adjacent to v

 if (!T[w].known)

 T[w].dist = Min(T[w].dist, c_{vw});

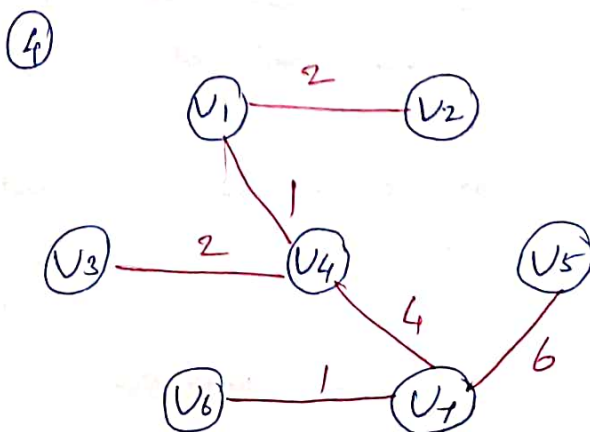
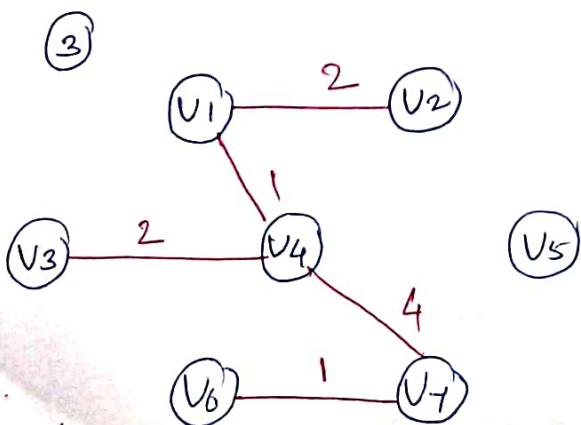
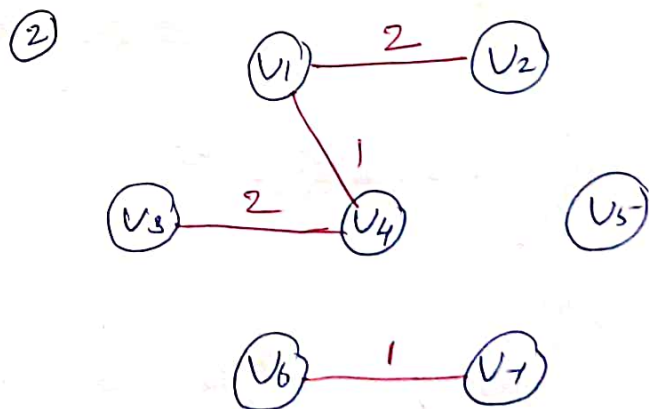
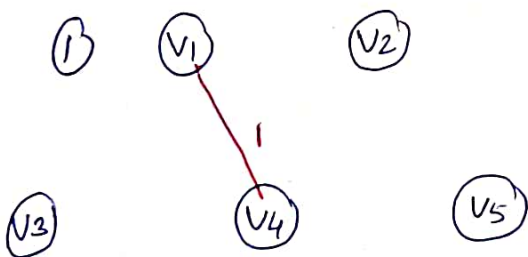
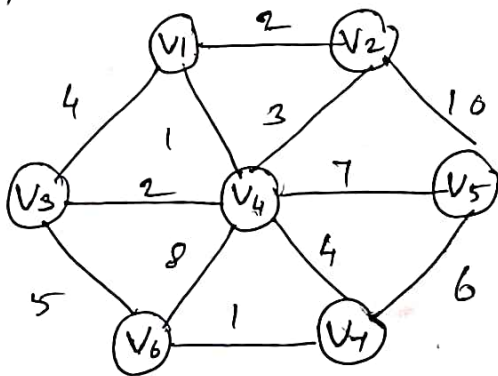
 T[w].path = v;

}

Kruskal Algorithm:

- Compute MST
- select edges in the order of smallest weight
- does not form cycle
- Terminate if all edges are accepted
- Maintain forest collection of trees.
- Two operation
 - ① find — returns the root of the tree
 - ② union — Merge two trees, make root ptr.

eg: 1



Edge	Weight	Action
(v_1, v_4)	1	Accepted
(v_6, v_7)	1	Accepted
(v_1, v_2)	2	Accepted
(v_4, v_3)	2	Accepted
(v_4, v_2)	3	Rejected
(v_4, v_7)	4	Accepted
(v_1, v_2)	4	Rejected
(v_4, v_5)	7	Rejected
(v_4, v_5)	6	Accepted
(v_2, v_5)	10	Rejected

Route:

void kruskal (Graph G)

{

int edges accepted;

Disjset S;

PriorityQueue H;

Vertex u, v;

setType Uset, Uset';

Edge E;

Initialize (S);

Read graph into heap array (G, H);

Build heap (H);

Edges accepted = 0;

while Edges Accepted \neq NumVertex - 1
{

$E = \text{DeleteMin}(H);$

$Uset = \text{find}(U, s);$

$Vset = \text{find}(V, s);$

if $(Uset \neq Vset)$

↳

Edges Accepted $++$;

setUnion $(s, Uset, Vset);$

↳ }
↳ }