

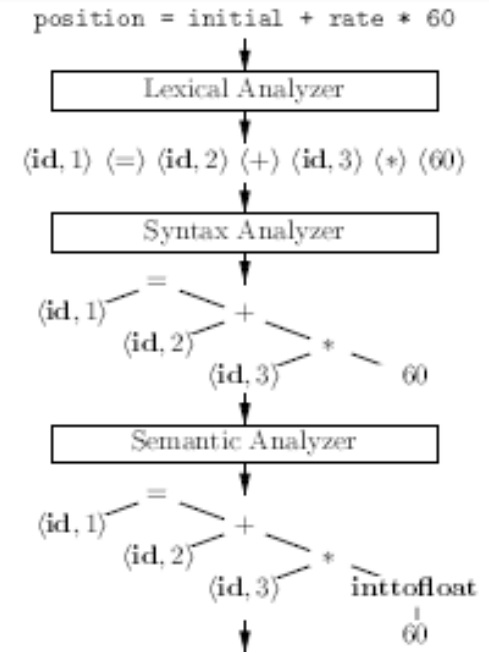


Semantic Analysis

- 3rd phase in compilation
- Parse tree & symbol table
 - Consistency of code
- **Type Checking**
- Semantic consistency
- Semantic Errors
 - Type mismatch
 - Undeclared variables
 - Reserved keywords
 - Accessing out of scope variable
 - Actual and formal parameter mismatch

| | | |
|---|----------|-----|
| 1 | position | ... |
| 2 | initial | ... |
| 3 | rate | ... |
| | | |

SYMBOL TABLE





Type Systems

- Type system
 - Rules that assign types to program constructs (constraints added to check the validity of program)
 - Language specific
- Type Checking (verifying the type)
- Static check (Compilation)
 - Type check, Type conversion, Indexing checks, Function call check, Flow of Control check
- Dynamic Check (Execution)
- Error Recovery
 - Location



Type System

- Type Expression
- Language Construct – part of program formed with lexical tokens. EX: if, else,...python → print (language construct)
- **Type Expression** – Associate each language construct with an expression describing its type
- Type expression → Basic Type & Type name
- **Basic Type** → int, char, boolean, void, float,...
- **Type Name / (Type Constructor)**
 - Array, Function, Pointer
 - Ex: Array → T is a Type Expression
 - → Array(I,T) is a Type Expression
 - → Array([1...10],int)



Type System

- Type Equivalence
 - Name Equivalence
 - Ex: `char a='A'; char b=a;`
 - Structure Equivalence
- Implementation of Type System is called Type Checker



Specification of a Simple Type Checker



- Example:

$P \rightarrow D ; E$

$D \rightarrow D ; D \mid id : T$

$T \rightarrow char \mid integer \mid array [num] \text{ of } T \mid \uparrow T$

$E \rightarrow literal \mid num \mid id \mid E \text{ mod } E \mid E [E] \mid E \uparrow$

- Translation Scheme

$P \rightarrow D ; E$

$D \rightarrow D ; D$

$D \rightarrow id : T \{ \text{addtype} (id.entry, T.type) \}$

$T \rightarrow char \{ T.type := char \}$

$T \rightarrow integer \{ T.type := integer \}$

$T \rightarrow \uparrow T1 \{ T.type := \text{pointer}(T1.type) \}$

$T \rightarrow array [num] \text{ of } T1 \{ T.type := \text{array} (1 \dots num.val, T1.type) \}$



Specification of a Simple Type Checker



- **Type Checking of Expression**

- $E \rightarrow \text{literal} \{ E.\text{type} := \text{char} \}$ $E \rightarrow \text{num} \{ E.\text{type} := \text{integer} \}$
- $E \rightarrow \text{id} \{ E.\text{type} := \text{lookup} (\text{id.entry}) \}$
- $E \rightarrow E1 \text{ mod } E2 \{ E.\text{type} := \text{if } E1.\text{type} = \text{integer} \text{ and } E2.\text{type} = \text{integer} \text{ then integer else type_error} \}$

- **Type checking of statements**

1. Assignment statement: $S \rightarrow \text{id} := E$
 $S \rightarrow \text{id} := E \quad \{ S.\text{type} := \text{if id.type} = E.\text{type} \text{ then void else type_error} \}$

- **Type checking of Functions**

2. Conditional statement: $S \rightarrow \text{if } E \text{ then } S1$
 $S \rightarrow \text{if } E \text{ then } S1 \quad \{ S.\text{type} := \text{if } E.\text{type} = \text{boolean} \text{ then } S1.\text{type} \text{ else type_error} \}$

3. While statement:

$S \rightarrow \text{while } E \text{ do } S1$
 $S \rightarrow \text{while } E \text{ do } S1 \quad \{ S.\text{type} := \text{if } E.\text{type} = \text{boolean} \text{ then } S1.\text{type} \text{ else type_error} \}$