



TOPIC:11.- Proof methods and strategy.

The Theory of Inference for Predicate calculus

Universal Specification (US)

If a statement of the form $(\forall x)[A(x)]$ is assumed to be true, then the universal quantifier can be dropped to obtain A(y) is true for any arbitrary object 'y' in the universe.

Existential Specification (ES)

From $J_{2}(A(z))$ one can conclude A(y), provided that y is not free in any given premix and also not free in any prior step of the durivation.

Universal Generalization (UG)

From A (4) one can / condude

$$A(y) \Rightarrow (\forall x)(A(x))$$

Existential Grunnalization (EG) $A(y) \Rightarrow (\exists z)(A(z))$





Express the Statement "Every student in this class as quantifiers."

Let C(x): x is in this class A(x): x has completed Assignment-I

For all x, if x is in this class, then x has completed Assignment-I

ompleted Assignment-L.

The Symbolic form $(\forall x)$ $(C(x) \rightarrow A(x))$

2.

Write each of the following in symbolic form.

All men are good (b) No men are good

Some men are good (e) Some men are not good

Let M(x): x is a man G(x): x is good

For all x, x is a man, then x is good.

($\forall x$) $(M(x) \rightarrow G(x))$





- (c) There exists an z, x is a man and z is : (3z) (M(z) A G(z))
- (d) There exists an x, x is a man and x is not good.

 (72) (11(2) 1 7 G(2))

3.

Show that
$$(\forall x)(P(x) \rightarrow Q(x)) \land (\forall x)(Q(x) \rightarrow R)$$

 $\Rightarrow (\forall x)(P(x) \rightarrow R(x))$.

§13	$ (\forall x) (P(x) \rightarrow Q(x)) $	Rule P
313	2) $P(y) \rightarrow Q(y)$	Rule US
{3}	3) $(\forall x) (Q(x) \rightarrow R(x))$	Rule P
	4) $Q(y) \rightarrow R(y)$	Rule US
{3} {1,3}	$5) P(y) \rightarrow R(y)$	Rule T ($P \rightarrow Q, Q \rightarrow R$) $\Rightarrow P \rightarrow R$)
ξ1,3 ζ	6) (4x) (P(x) → R(x)).	
})	





4)

Show that $(\forall x)(P(x) \lor Q(x)) \Rightarrow (\forall x) P(x) \lor (\exists x)$ We use indirect method, by assuming

P(2) V (3x)(Q(x)) as an additional primise

£13	(x) P(x) P(x) V (∃x) Q(x)	Rule P
\$1 }	2) $(\exists x) \neg P(x) \land$ $(\forall x) \neg Q(x)$	Rule T (Demorgan's)
{1}	3) (3x) ¬p(x)	Rule T (PAQ ⇒ P)
13	A) (+71) -1 Q(7)	Ruli T (PAa ⇒ a)
§13	5) ¬ P(y)	Rule ES
£13	6) 7 Q(4)	Rule US
£13	7) - P(4) 1-Q(4)	Rule T (P,Q ⇒ PAQ)
£13	8) - (P(4) VQ(4))	Rule T (Dumorgan's)
{9 }	9) (+x) (P(x) VQ(x))	Rule P



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593	10) P(4) VQ(4)	Rule US
{1,9}	11) [P(y) VQ(y)]^ -[P(y) VQ(y)]	Rule T (P, a → PAQ)
£1,97	12) F	

bring thing is a plant or animal John's gold fish is alive and it is not a plant. All animals have hearts. Therefore John's goldfish has a heart.

Let
$$L(x)$$
: x is a living thing $P(x)$: x is a plant $A(x)$: x is an animal has a heart

H(x): x has a heart

Then, the given premises are

M, the given
$$\Gamma$$

(1) $(\forall x) \left[L(x) \rightarrow (P(x) \cdot VA(x)) \right]$

(2) L(j)
$$\Lambda \neg P(j)$$

(3)
$$(\forall x) [A(x) \rightarrow H(x)]$$

conclusion is H(i)





	/ 1	and the second s
{1}	1) $(\forall x)$ $[L(x) \rightarrow (P(x) \lor A(x))]$	x)) Rule P
{ 1 }	2) $L(j) \rightarrow P(j) \vee A(j)$) Rule US
{3}	3) L(j) A¬P(j)	Rule P
{3}	4) L(j), ¬P(j)	Rule T (PAa ⇒P.a)
{1,3}	5) P(j) V [A(j)	Rule T(P, P→a ⇒a)
<i>§</i> 1,3 <i>§</i>	6) $\neg P(j) \rightarrow A(j)$	Rule T(P>a ⇒¬PVa)
ξ ηζ	7) $(\forall x) (A(x) \rightarrow H(x))$	Rule P
ξη	8) $A(j) \rightarrow H(j)$	Rule US
1,3,7}	q) $\neg P(j) \rightarrow \mathbf{h}(j)$	Rule T (P→Q,Q→R ⇒P→R)
1,3,7 }	10) H(j)	Rule $T(P, P \rightarrow a \Rightarrow a)$