



- Modification Anomalies
- Ex: (Empid, Empname, Projectid, Projectname)
- ➢ Update Anomalies: If any changes in Projectname , it changes to empid, empname in the relation.
- ≻ Insertion Anomalies: If Projectid is NULL.
- Deletion Anomalies: If delete projectid, project name is deleted, then it has NULL values





Sid	Sname	Credits	DeptName	Building	RoomNo
1	Rahul	5	CSE	B1	101
2	Jiva	8	CSE	B1	101
3	Jenny	9	IT	B2	201
4	Palani	9	IT	B2	201
5	Ashok	7	Civil	B3	110
6	Aakash	7	ECE	B1	115
7	Vanitha	8	Civil	B1	110
8	Tom	7	CSE	B1	101





- > Decomposition
 - \succ the process of breaking down one table into multiple tables.

ID	Name	Age
1	А	20
2	А	21

ID	Name	Name	Age
1	А	А	20
2	А	А	21

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➤ Decomposition – Lossy:

ID	Name		Name	Age
1	А	$\left \times\right $	А	20
2	А	<i>v</i> 4	А	21

ID	Name	Age
1	А	20
1	А	21
2	А	20
2	А	21

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> Decomposition – Problems:

> Some queries more expensive

 \succ loss of information during decomposition





- Decomposition Properties:
 - ➤ Loss-less Join or Non Loss Decomposition When all information found in the original database is preserved after decomposition.
 - Dependency Preservation This is a property in which the constraints on the original table can be maintained by simply enforcing some constraints on each of the smaller relations.





> a process of organizing the data in database to avoid

data redundancy

≻Forms:

≻ First normal form(1NF)

≻ Second normal form(2NF)

≻ Third normal form(3NF)

► Boyce & Codd normal form (BCNF)





- > First normal form(1NF) :
 - ➤ an attribute (column) of a table cannot hold multiple values. It should hold only atomic values.

emp_id	emp_name	emp_address	emp_mobile
101	Herschel	New Delhi	8912312390
102	100	Variation	8812121212
102 J	JOII	Kanpur	9900012222
103	Ron	Chennai	7778881212
104	Lester	Bangalore	9990000123
			8123450987





 \geq each attribute of a table must have atomic (single) values.

emp_id	emp_name	emp_address	emp_mobile
101	Herschel	New Delhi	8912312390
102	Jon	Kanpur	8812121212
102	Jon	Kanpur	9900012222
103	Ron	Chennai	7778881212
104	Lester	Bangalore	9990000123
104	Lester	Bangalore	8123450987

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- Prime Attribute:
 - ➢ Member of some candidate Key.
 - ➢ Example: RegisterNo.
- > Non-prime Attribute:
 - ➢ not a member of any candidate key





- Second normal form(1NF) : A table is said to be in
 2NF if both the following conditions hold:
 - ≻ Table is in 1NF (First normal form)
 - ➢ All the non-key columns are dependent on the table's primary key. (OR)
 - Every attribute is in an candidate Key,
 - > Every attribute depend fully on every candidate key.





- ➤ (<u>id</u>, <u>projid</u>, hrs, name, projname)
- > (id, projid) -> hrs, name, projname
- ➢ id → name
- > projid -> projname
- ≻Name and projname partial dependency
- > where name and projname provides redundancy.
- > Name modified all relation's names are modified.





- (id, projid, hrs, name, projname)
- (id,projid) -> hrs, name, projname
- ➢ id → name
- > projid -> projname
- ►2NF: (<u>id,projid,</u>hrs), (<u>id,</u>name),

(projid, projname)





NORMALIZATION – 2NF

	teacher_id	subject	teacher_age
Δ	111	Maths	38
	111	Physics	38
	222	Biology	38
	333	Physics	40
	333	Chemistr y	40

Candidate Keys: {teacher_id, subject}

Non prime attribute: teacher_age





NORMALIZATION – 2NF

Ν	teacher_id	subject	teacher_age
	111	Maths	38
	111	Physics	38
	222	Biology	38
	333	Physics	40
	333	Chemistry	40

The table is in first normal form and all the columns depend on the table's primary key.

		teacher_id	subject
teacher id	teacher age	111	Maths
		111	Physics
111	38	222	Biology
222	38	222	Diology
333	40	333	Physics
	DBMS/	R.KAMALAKKANNAAJAP/CSE-IOT	Chemistry





- A table design is said to be in 3NF if both the following conditions hold:
 - ≻ Table must be in 2NF
 - Transitive functional dependency of non-prime attribute on any super key should be removed.

OR

Every non prime attribute

- ➢ fully FD on every key, and
- ➢ non transitively FD on every key





- ➤ (id, name, projid, projname)
- ➤ id -> name, projid
- projid -> projname
- > From projname, can find projid, from projid find id.
- (id, name, projid) and (projid, projname)





- (id, dist, taluk, area, price, rate)
- ➢ id dist, taluk, area, price, rate
- > dist, taluk -> id, area, price, rate
- ➤ dist -> rate
- ➤ area -> price





- Second Normal Form
- ➤(id, dist, taluk, area, price, rate)
- ➢ id dist, taluk, area, price, rate
- > dist, taluk -> id, area, price, rate
- ➤ dist -> rate
- ➤ area -> price
- ➢ price not fully FD





- Second Normal Form
- \succ (<u>id</u>, dist, taluk, area, price)
- ➢ id dist, taluk, area, price, rate
- > dist, taluk -> id, area, price, rate
- > area -> price
- \succ (<u>dist</u>, rate)
- dist -> rate





- Third Normal Form
- \succ (<u>id</u>, dist, taluk, area, price)
- id -> area -> price -> it is transitive
- ➤ (id, dist, taluk, area)
- ➤ (<u>area</u>, price)
- \succ (<u>dist</u>, price) It is 3NF







Thank You.....

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