



SNS COLLEGE OF ENGINEERING
(Autonomous)
DEPARTMENT OF ELECTRONICS AND COMMUNICATIONS ENGINEERING



Quine-McCluskey (Tabular) Minimization



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Two step process utilizing tabular listings to:

- Identify prime implicants (implicant tables)
- Identify minimal PI set (cover tables)

All work is done in tabular form

- Number of variables is not a limitation
- Basis for many computer implementations
- Don't cares are easily handled

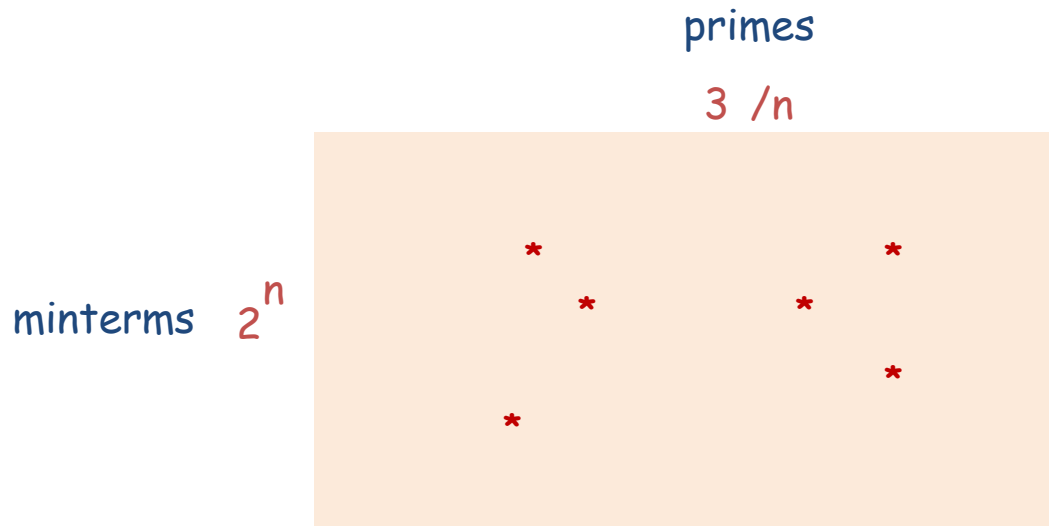
Proper organization and term identification are key factors for correct results



Difficulty

Note: Can be

- 2^n minterms
- $\sim 3^n/n$ primes



- ✓ Thus $O(2^n)$ rows and $O(3^n/n)$ columns and minimum covering problem is NP-complete.

Example

$$F = \bar{w}\bar{x}\bar{y}\bar{z} + \bar{w}x\bar{y}z + \bar{w}xy\bar{z} + wx\bar{y}z$$

$$d = x\bar{y}\bar{z} + w\bar{y}\bar{z} + xyz + \bar{x}y\bar{z} + wy\bar{z} + \bar{w}\bar{x}\bar{y}z$$

Karnaugh map

$\bar{w}\bar{y}$	w	x	yz	
				\bar{z}
				x

	\bar{z}	x	$\bar{w}\bar{y}$
$\bar{w}\bar{x}\bar{y}\bar{z}$	1	0	1
$\bar{w}x\bar{y}z$	0	1	1
$\bar{w}xy\bar{z}$	1	1	0
$wx\bar{y}z$	0	1	0

$$\sum_c = \bar{z} + x + \bar{w}\bar{y}$$

Covering Table

Solution: $\sum_m = \bar{z} + x$
 (also $x + \bar{w}\bar{y}$)



Covering Table

Minterms of f

	\bar{z}	x	$\bar{w}\bar{y}$	Primes of f+d
$\bar{w}\bar{x}\bar{y}\bar{z}$	1	0	1	
$\bar{w}x\bar{y}z$	0	1	1	
$\bar{w}xy\bar{z}$	1	1	0	
$\bar{w}\bar{x}\bar{y}z$	0	1	0	← Row singleton (essential minterm)

↑
Essential prime

- ✓ **Definition:** An essential prime is any prime that **uniquely** covers a minterm of f.



Quine-McCluskey Minimization (cont.)



- ✓ Terms are initially listed **one per line in groups**
 - Each group contains terms with the **same number of true and complemented variables**
 - Terms are listed in numerical order within group
- ✓ Terms and implicants are identified using one of three common notations
 - full variable form
 - cellular form
 - 1,0,- form



Example of Different Notations



$$F(A, B, C, D) = \sum_m (4, 5, 6, 8, 10, 13)$$

	Full variable	Cellular	1,0,-
1	$\overline{A}\overline{B}\overline{C}\overline{D}$	4	0100
	$\overline{A}\overline{B}C\overline{D}$	8	1000
2	$\overline{A}B\overline{C}D$	5	0101
	$\overline{A}BC\overline{D}$	6	0110
3	$A\overline{B}\overline{C}\overline{D}$	10	1010
	$AB\overline{C}D$	13	1101



Notation Forms



- Full variable form:** variables and complements in algebraic form
- hard to identify when adjacency applies
 - very easy to make mistakes
- ✓ **Cellular form:** terms are identified by their decimal index value
- easy to tell when adjacency applies; indexes must differ by **a power of two (one bit)**
- ✓ **1,0,- form:** terms are identified by their binary index value
- **easier to translate to/from full variable form**
 - **easy to identify when adjacency applies, one bit is different**
 - **shows variable(s) dropped when adjacency is used**
- ✓ Different forms may be mixed during the minimization



Implication Table

(1,0,-)



✓ Quine-McCluskey Method

- Tabular method to systematically find all prime implicants

$$f(A,B,C,D) = \sum_m(1,2,5,6,7,9,10) + \sum_d(0,13,15)$$

- Part 1: Find all prime implicants
- Step 1: Fill Column 1 with **onset** and **DC-set** minterm indices.
Group by number of true variables (# of 1's).

NOTE THAT DCs ARE INCLUDED IN THIS STEP!

Implication Table		
Column I		
0000		
0001		
0010		
0101		
0110		
1001		
1010		
0111		
1101		
1111		



Minimization - First Pass (1,0,-)



✓ Quine-McCluskey Method

- Tabular method to systematically find all prime implicants
- $f(A,B,C,D) = \sum_m(1,2,5,6,7,9,10) + \sum_d(0,13,15)$
- Part 1: Find all prime implicants
- Step 2: Apply Adjacency - Compare elements of group with N 1's against those with N+1 1's. One bit difference implies adjacent. Eliminate variable and place in next column.

E.g., 0000 vs. 0100 yields 0-00

0000 vs. 1000 yields -000

When used in a combination, mark with a check. If cannot be combined, mark with a star. These are the prime implicants.

Repeat until nothing left.

Implication Table

Column I		Column II	
0000	0	000-	0,1
		00-0	0,2
0001	1		
0010	2	0-01	1,5
		-001	1,9
0101	5	0-10	2,6
0110	6	-010	2,10
1001	9		
1010	10	01-1	5,7
		-101	5,13
0111	7	011-	6,7
1101	13	1-01	9,13
1111	15	-111	7,15
		11-1	13,15



Minimization - Second



Pass (1,0,-)

- ✓ Quine-McCluskey Method
 - Step 2 cont.: Apply Adjacency - Compare elements of group with N 1's against those with N+1 1's. One bit difference implies adjacent. Eliminate variable and place in next column.
E.g., 0000 vs. 0100 yields 0-00
00-0 vs. 10-0 yields -0-0
 - When used in a combination, mark with a check ✓.
 - If cannot be combined, mark with a star *. **THESE ARE THE PRIME IMPLICANTS.**
Repeat until nothing left.
- ✓ The set of * constitutes the Complete Sum Σ_c

Implication Table		
Column I	Column II	Column III
0000 ✓ 0	000- 0,1 00-0 0,2	--01 1,5,9 -1-1 5,7,11
0001 ✓ 1		
0010 ✓ 2	0-01 1,5 -001 1,9	
0101 ✓ 5	0-10 2,6	
0110 ✓ 6	-010 2,10	
1001 ✓ 9		
1010 ✓ 10	01-1 5,7 -101 5,13	
0111 ✓ 7	011- 6,7	
1101 ✓ 13	1-01 9,13	
1111 ✓ 15	-111 7,15 11-1 13,15	



Prime Implicants



$$f(A,B,C,D) = \Sigma_m (1,2,5,6,7,9,10) + \Sigma_d (0,13,15)$$

		C			
		D	00	01	11
A	B	00	01	11	10
	00	X	1	0	1
	01	0	1	1	1
	11	0	X	X	0
	10	0	1	0	1

Prime Implicants:

$$000 - = \bar{A} \bar{B} \bar{C}$$

$$0 - 10 = \bar{A} C \bar{D}$$

$$011 - = \bar{A} B C$$

$$- - 01 = \bar{C} D$$

$$00 - 0 = \bar{A} \bar{B} \bar{D}$$

$$- 010 = \bar{B} C \bar{D}$$

$$- 1 - 1 = B D$$

Stage 2: find smallest set of prime implicants that cover the active-set
 Note that essential prime implicants must be in the final expression



Coverage Table



rows = prime implicants
 columns = ON-set elements (minterms)

place an "X" if ON-set element is covered by the prime implicant

NOTE: DON'T INCLUDE DCs IN COVERAGE TABLE; THEY DON'T HAVE TO BE MANDATORY COVERED

Coverage Chart
minterms

		1	2	5	6	7	9	10
primes	0,1	000-	X					
	0,2	00-0		X				
	2,6	0-10		X		X		
	2,10	-010		X				X
	6,7	011-				X	X	
	1,5,9,13	--01	X		X			X
	5,7,13,15	-1-1			X		X	



Row and Column Dominance

- ✓ **Definition:** Given two rows i_1 and i_2 , a row i_1 is said to **dominate** i_2 if it has checks in all columns in which i_2 has checks, i.e. it is a superset of i_2

Example:

i_1	x	x	x	x	x
i_2	x	x		x	x

i_1 dominates i_2

- ✓ We can remove row i_2 , because we would never choose i_2 in a minimum cover since it can always be replaced by i_1 (i_2 is anymore a prime implicant).

DOMINATED ROWS (IMPLICANTS) CAN BE ELIMINATED



Row and Column Dominance



- ✓ **Definition:** Given two columns j_1 and j_2 , if the set of primes of column j_2 is contained in the set of primes of column j_1

Example:

j_2 dominates j_1

	j_1	j_2
	x	
	x	x
	x	
	x	x

- ✓ We can remove column j_1 since we have to choose a prime to cover j_2 , any such prime also covers j_1 , that would result covered as well.

DOMINATED COLUMNS (MINTERMS) CAN BE ELIMINATED



Pruning the Covering Table



1. Remove all rows covered by **essential primes** (columns in row **singletons**). Put these primes in the cover G .
 2. **Group identical rows together and remove dominated rows.**
 3. **Remove dominating columns.** For equal columns, keep just one to represent them.
 4. Newly formed row singletons define **n -ary essential primes.**
 5. Go to 1 if covering table decreased.
- ✓ The algorithm may terminate successfully with a set of primes and an empty table.
 - ✓ In case it terminate with a non empty table, the resulting reduced covering table is called the **cyclic core**. This has to be solved. A minimum solution for the cyclic core must be added to the resulting G .



Coverage Table (cont.)



Coverage Chart

		1	2	5	6	7	9	10
0,1	000-	X						
0,2	00-0		X					
2,6	0-10		X		X			
2,10	-010		X					X
6,7	011-				X	X		
1,5,9,13	--01	X		X			X	
5,7,13,15	-1-1			X		X		

		1	2	5	6	7	9	10
0,1	000-	X						
0,2	00-0		X					
2,6	0-10		X		X			
2,10	-010		X					X
6,7	011-				X	X		
1,5,9,13	--01	X		X			X	
5,7,13,15	-1-1			X		X		

If column has a single x, than the implicant associated with the row is **essential**.
It must appear in the minimum cover



Coverage Table (cont.)



		1	2	5	6	7	9	10
0,1	000-	X						
0,2	00-0		X					
2,6	0-10		X		X			
2,10	-010		X					X
6,7	011-				X	X		
1,5,9,13	--01	X		X			X	
5,7,13,15	-1-1			X		X		

Eliminate all columns covered by essential primes

		1	2	5	6	7	9	10
0,1	000-	X						
0,2	00-0		X					
2,6	0-10		X		X			
2,10	-010		X					X
6,7	011-				X	X		
1,5,9,13	--01	X		X			X	
5,7,13,15	-1-1			X		X		

Find minimum set of rows that cover the remaining columns

$$F = \bar{B}C\bar{D} + \bar{A}BC + \bar{C}D$$

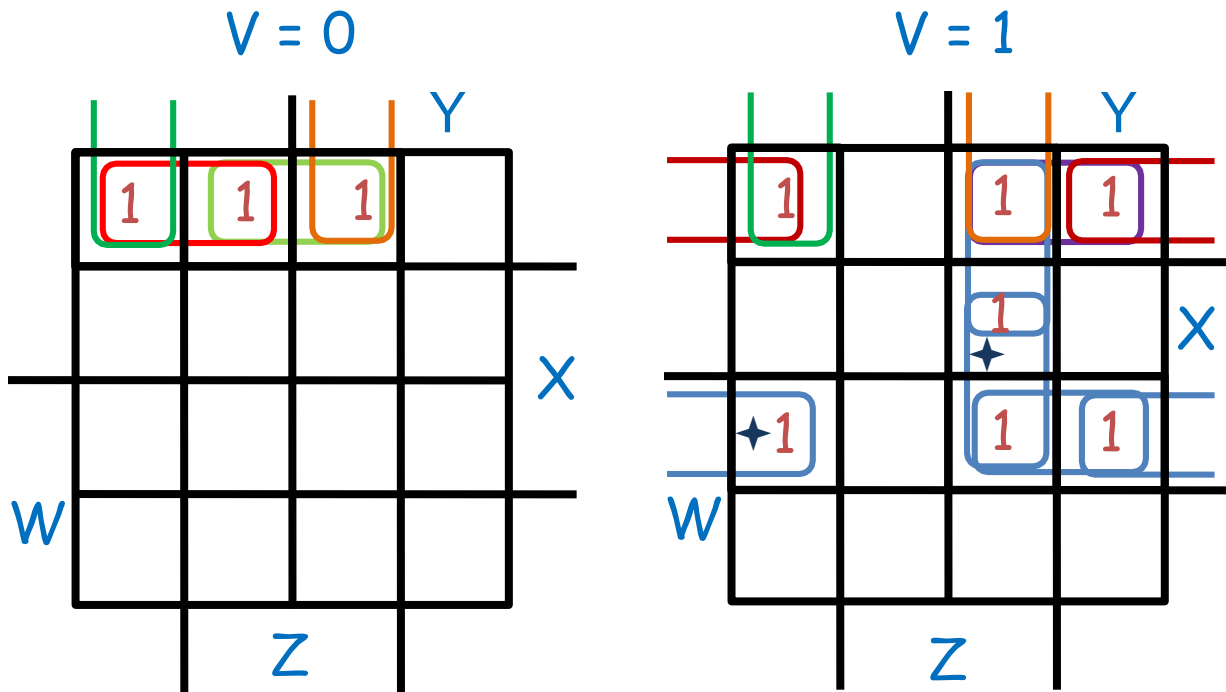


Quine Mc Clunskey: Cyclic

Core example



$$F = \sum_m (0, 1, 3, 16, 18, 19, 23, 28, 30, 31)$$



$$F = v'w'y'x' + v'w'x'z + w'x'y'z + w'x'yz + vw'x'z' + vw'x'y + vwXZ' + vwxy + vw'y z + vx y z$$

A B C D E F G H I J



Implication Table (1,0,-)



✓ Quine-McCluskey Method

- Tabular method to systematically find all prime implicants
- $f(v,w,x,y,z) = \sum m(0,1,3,16,18,19,23,28,30,31)$
- Part 1: Find all prime implicants
- Step 1: Fill Column 1 with active-set and DC-set minterm indices. Group by number of true variables (# of 1's).

Implication Table			
	Column I	Column II	
0	00000	0000-	A: 0
1	00001	-0000	C: 0 1
16	1 0000	000-1	B: 1
3	00 011	100-0	E: 16 18
18	1001 0	-0011	D: 3 19
19	100 11	1001-	F: 18 19
28	1 1 100	10-11	I: 19 23
23	10 111	111-0	G: 28 30
30	111 10	1-111	J: 23 31
31	1111 1	1111-	H: 30 31

$$F = v'w'y'x' + v'w'x'z + w'x'y'z' + w'x'yz + vw'x'z' + vw'x'y + vw'xz + vwxy + vw'yz + vxyz$$

A B C D E F G H I J



Quine Mc Clunskey



$$F = \sum_5(0,1,3,16,18,19,23,28,30,31)$$

	0	1	3	16	18	19	23	28	30	31
A	*	*								
B		*	*							
C	*			*						
D			*			*				
E				*	*					
F					*	*				
G								*	*	
H									*	*
I						*	*			
J							*			*



$G+J+A+D+E;$

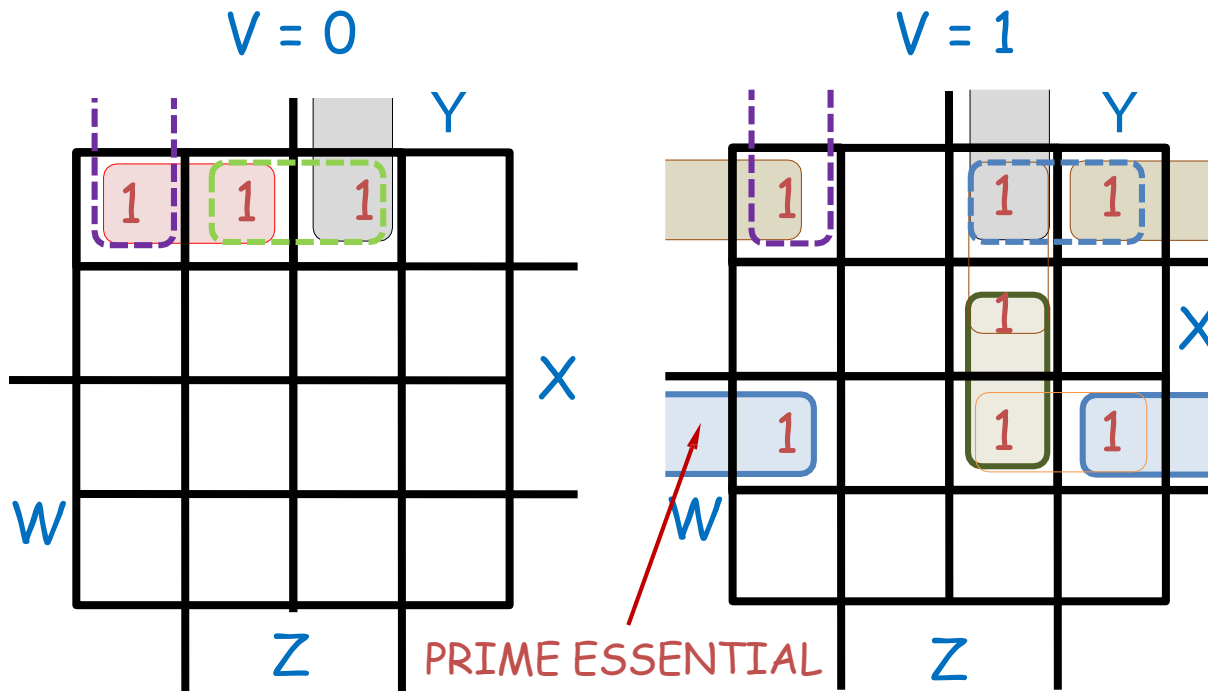


Quine Mc Clunskey: Cyclic

Core example



$$F = \sum_m(0,1,3,16,18,19,23,28,30,31)$$



G: 28 3
 J: 23 3
 A: 0
 D: 3 1
 E: 16 1

$$F = v'w'y'x' + v'w'x'z + w'x'y'z' + w'x'yz + vw'x'z' + vw'x'y + vw'xz' + vwxy + vw'y'z + vxyz$$

A B C D E F G H I J

$$G + J + A + D + E;$$

Generating Primes – multiple outputs

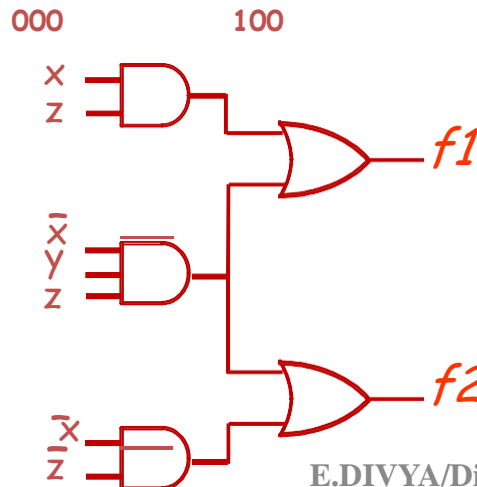
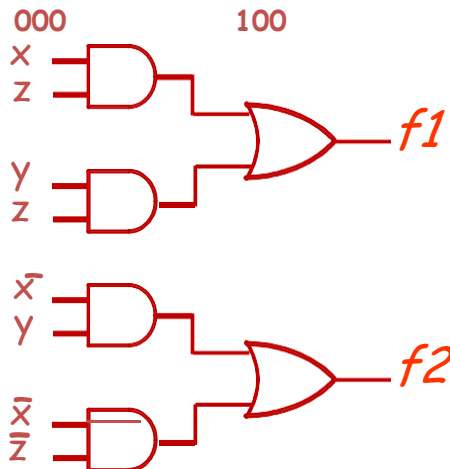
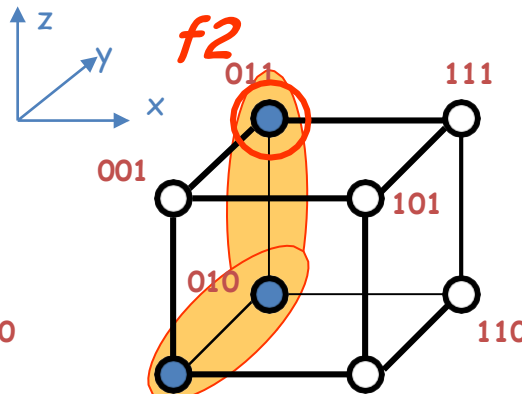
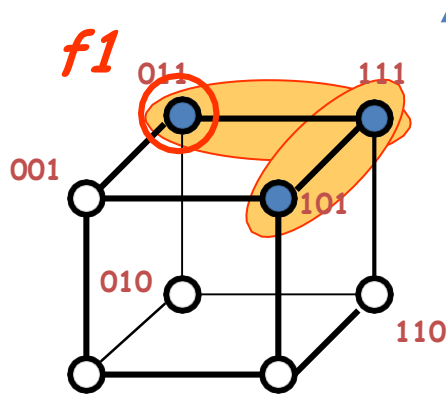
Example: $f_1(x, y, z) = \sum m(3,5,7)$, $f_2(x, y, z) = \sum m(0,2,3)$

x	y	z	
0	1	1	-11
1	0	1	1-1
1	1	1	

$$f_1 = yz + zx$$

x	y	z	
0	0	0	-0-0
0	1	0	01-
0	1	1	

$$f_2 = \bar{x}\bar{z} + \bar{x}y$$



The idea is that we can share terms:

- using separate optimizations 6 gates and $G=12$
- sharing a term 5 gates and $G=11$.

$$f_1 f_2 = (yz + zx)(\bar{x}\bar{z} + \bar{x}y)$$

$$f_1 f_2 = \bar{x}yz$$



Generating Primes – multiple output

- ✓ Theorem: if p_1 is a prime implicant for f_1 , and p_2 is a prime implicant for f_2 , then if $p_1 \cdot p_2 \neq 0$, $p_1 \cdot p_2$ is a prime implicant of $f_1 \cdot f_2$
- ✓ Theorem: if p_3 is a prime implicant for $f_1 \cdot f_2$, then there exist p_1 for f_1 , and p_2 for f_2 , such that $p_3 = p_1 \cdot p_2$
- ✓ We can conclude that all prime implicants of $f_1 \cdot f_2$ are minimal sharable products for f_1 and f_2 ; and that all prime implicants for $f_1 \cdot f_2$ are created by products of prime implicants for f_1 and f_2
- ✓ The way to use this is to make the prime implicants of $f_1 \cdot f_2$ available to the minimizations of f_1 and f_2 by extending the table concept



Generating Primes – multiple outputs



- ✓ Procedure similar to single-output function, except: include also the primes of the products of individual functions

	f_1 minterms	f_2 minterms
Rows for f_1 prime implicants: mark only f_1 columns		
Rows for f_2 prime implicants: mark only f_2 columns		
Rows for $f_1 f_2$ prime implicants: mark both f_1 and f_2 columns		



Minimize multiple-output cover



✓ Example, cont.

$$f_1 \quad m_3 = 011 \quad p_1 = yz$$

$$m_5 = 101 \quad p_2 = xz$$

$$m_7 = 111$$

$$f_2 \quad m_0 = 000 \quad p_3 = \bar{x}y$$

$$m_2 = 010 \quad p_4 = \bar{x}\bar{z}$$

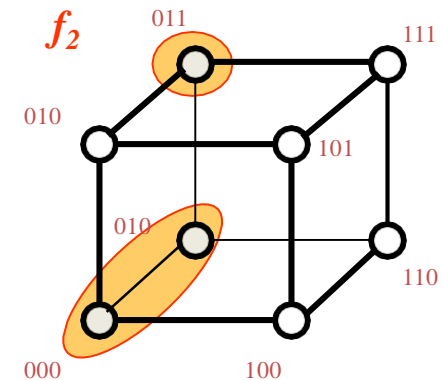
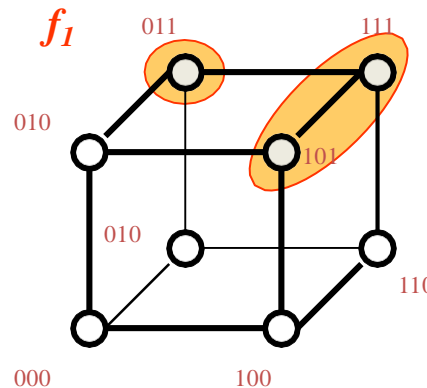
$$m_3 = 011$$

$$f_2 \quad m_3 = 011 \quad p_5 = \bar{x}yz$$

Min cover has 3 primes:
 $F = \{ p_2, p_4, p_5 \}$

	m_3	m_5	m_7	m_0	m_2	m_3		m_3	m_3
p_1	✓		✓				p_1	✓	
p_2		✓	✓				p_3		✓
p_3					✓	✓	p_5	✓	✓
p_4				✓	✓				
p_5	✓								✓

Note that row p_5 dominates rows p_1 and p_3 , removing these rows, the coverage is complete





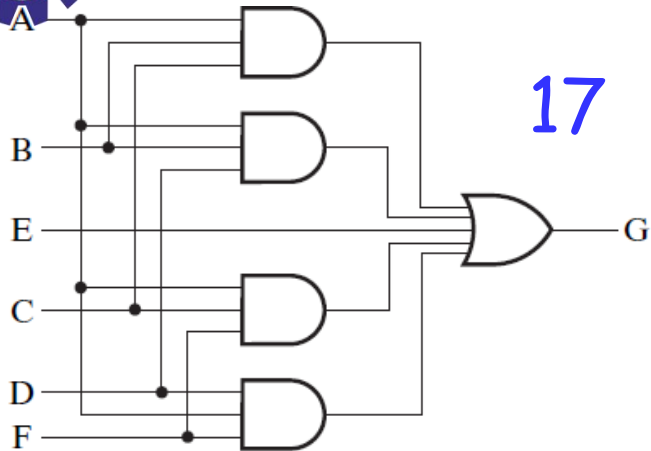
Multiple-Level Optimization



- ✓ Multiple-level circuits: circuits that are not two-levels (with or without input and/or output inverters)
- ✓ **Multiple-level circuits can have reduced gate input cost** compared to two-level (SOP and POS) circuits, obviously **augmenting the execution time**
- ✓ Multiple-level optimization is performed by applying transformations to circuits represented by equations while evaluating cost and execution time

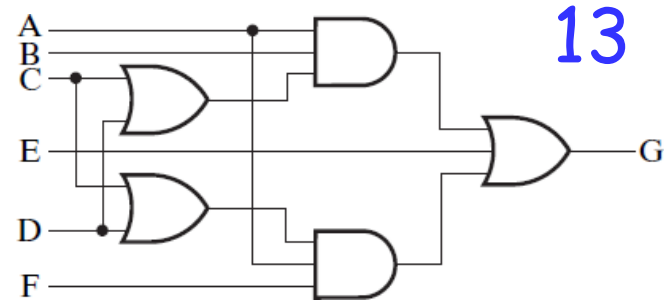


Multi level circuits – circuit analysis

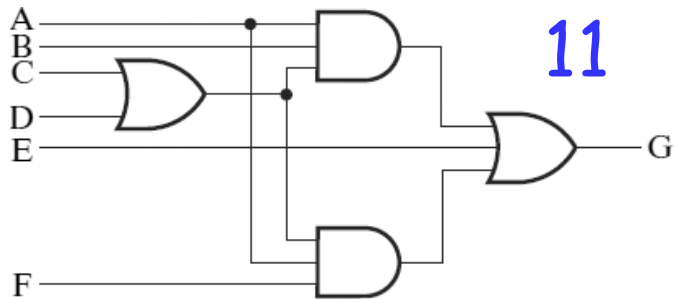


(a)

$$G = ABC + ABD + E + ACF + ADF$$

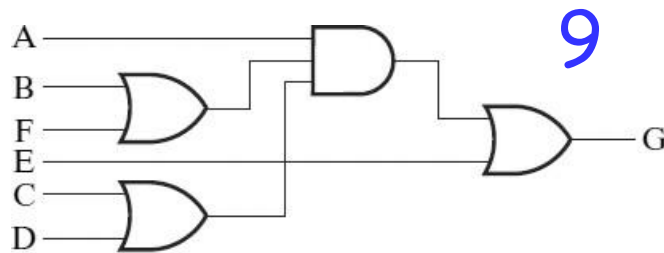


$$G = AB(C+D) + E + AF(C+D)$$



(c)

$$G = (C+D)(AB+AF) + E$$



(d)

$$G = A(C+D)(B+F) + E$$



Basic Transformations



- ✓ **Factoring** - finding a factored form from SOP or POS expression
- ✓ **Elimination of G into F** - expression function F as a function of G and some or all of its original variables
- ✓ **Extraction** - decomposition applied to multiple functions simultaneously



Transformation Examples



✓ Algebraic Factoring

$$F = \bar{A}\bar{C}\bar{D} + \bar{A}B\bar{C} + ABC + AC\bar{D} \quad G = 16$$

- Factoring:

$$F = \bar{A}(\bar{C}\bar{D} + B\bar{C}) + A(BC + C\bar{D}) \quad G = 18$$

- Factoring again:

$$F = \bar{A}\bar{C}(B + \bar{D}) + AC(B + \bar{D}) \quad G = 12$$

- Factoring again:

$$F = (\bar{A}\bar{C} + AC)(B + \bar{D}) \quad G = 10$$



Transformation Examples



Elimination

Beginning with two functions: $X = B + C$ $Y = A + B$

$$Z = \bar{A} X + C Y \qquad G = 10$$

- Eliminating X and Y from Z :

$$Z = \bar{A} (B + C) + C (A + B) \qquad G = 10$$

- "Flattening" (Converting to SOP expression):

$$Z = \bar{A} B + \bar{A} C + AC + BC \qquad G = 12$$

- This has increased the cost, but has provided a new SOP expression for two-level optimization.
- Two-level Optimization

$$Z = \bar{A} B + C \qquad G = 4$$

- Increasing gate input count G temporarily can result in a final solution with a smaller G



Transformation Examples



Extraction

- Beginning with two functions:

$$E = \overline{A}\overline{B}\overline{D} + \overline{A}BD$$

$$H = \overline{B}C\overline{D} + BCD$$

$$G = 16$$

- Finding a common factor and defining it as a function:

$$F = \overline{B}\overline{D} + BD$$

- We perform extraction by expressing E and H as the three functions:

$$E = \overline{A}F, H = CF$$

$$G = 10$$

- The reduced cost G results from the sharing of logic between the two output functions



THANK YOU